

### Assignment #1:

1. Consider the following problem:  $m =$  the number of machines  $= 1$ ;  $n =$  the number of jobs  $= n$ ;  $r_j =$  "release" date – time at which processing of job  $j$  can begin;  $p_j =$  total processing time for job  $j$ ;  $w_j =$  weight of job  $j$ . [Assume all these are input data of the problem and are positive numbers]. We want a schedule  $S$  that minimizes  $\sum_{j=1}^n w_j F_j(S)$  where  $F_j(S)$  is the "flow" time of job  $j$  in  $S$  and is given by

$$F_j(S) = C_j(S) - r_j$$

where  $C_j(S)$  is the completion time of job  $j$  in  $S$ .

- (a) First consider the case where  $w_j = 1 \forall j$ . If preemptions are allowed, show that SRPT rule is optimal. SRPT rule assigns for processing at all times the job with the *shortest remaining processing time*. [Make sure to argue that inserted idle time is not present in optimal schedules]
  - (b) Now consider the case with general values for  $w_j$ . Devise an algorithm to solve this problem again allowing preemptions. [Make sure to argue that inserted idle time is not present in optimal schedules]
  - (c) Now consider the case where preemption is not allowed. In this part, let  $w_j = 1 \forall j$ . Show that we can no longer assume that optimal solutions have no inserted idle time.
2. Consider the following problem:  $m = 1$ ;  $r_j = 0 \forall j$ ;  $p_j, w_j$  are as in the previous problem. We want a schedule  $S$  that minimizes  $\sum_{j=1}^n w_j [1 - e^{-\alpha C_j(S)}]$  where  $0 < \alpha < 1$ . Show that WDSPT [weighted discounted shortest processing time) rule produces an optimal schedule. This is a schedule that processes jobs in decreasing order of  $\frac{w_j e^{-\alpha p_j}}{1 - e^{-\alpha p_j}}$ . [Make sure you consider inserted idle time and preemption.]