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Abstract Data Types (ADTs)

Why ADTs?

Main Concepts of ADTs

ADTs as Objects

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Why Abstract Data Types (ADTs)?

Why Module Interconnection Languages (MILs)?

"Programming-in-the-small vs. programming-in-the-large"

Programming-in-the-small

Programming-in-the-small

Goal building "programs"

Problem Usually clear, small

building "software"

Usually unclear, large
sw. architecture

Emphasis Detailed design & impl.

Technique Structured programming

"divide & conquer" "separation of concerns"

Notation PLs

(Formal) (OO) Specification Lang.

anpower single person/small number usually single

multi-person multi-version

interface-oriented

hiding the representations

localized change: change in detailed design & impl. does not affect other (client) objects

specifying the representations

conceptual consolution

conquering complexity

decomposing problems into collections of interacting components (-> encapsulation)

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Main Concepts of ADTs

- model system as a collection of ADTs
 - @ algebraic specification
 - model-theoretic (axiomatic) specification -> Z
- Larch is for Algebraic Interface Specification

["Larch: Languages and Tools for Formal Specification", J. V. Guttag, J. J. Horning, S. J. Garland, K. D. Jones, A. Modet & J. M. Wing, Springer-Verlag, 1993]

- **☞ Each ADT includes:**
 - ø data objects
 - c operations on data objects (specified as functions with domain and range)
 - essential properties of operations algebraic equations in FOL with equality (specified in algebraic equations in FOL with equality)
- Two-tiered approach to software development
 - LCL Larch Common Language
 - PL-independent notation for writing interface spec.
 - LSL Larch Shared Language
 - PL-dependent notation for writing interface spec. LSLs exist for SmallTalk, Modula-3, C, C++, CLU, ...
- large library of predefined specs for common data types
- Larch theorem prover for correctness of properties

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Main Concepts of ADTs

```
Example
```

```
/* A trait describes an ADT, i.e., data objects, ops, properties of op
Stack (E, C): trait
                         & links to other ADTs */
```

introduces

/* E is an element (e.g., int, str, bool) of C, a container */

new: -> C

push: C, E -> C

top: C -> E exempting top (new)

pop: C -> C exempting pop (new)

isEmpty: C -> Bool

asserts

C generated by new, push

forall stk: C, e: E

top(push(stk, e)) == e

pop(push(stk, e)) = stk

isEmpty (new)

~ isEmpty (push (stk, e))

implies

LinearContainer (push for insert, top for first, pop for rest)

Main Concepts of ADTs

Questions

```
Stack (E, C): trait

introduces

new: -> C

push: C -> E exempting top (new)
pop: C -> C exempting pop (new)
isEmpty: C -> Bool
asserts

C generated by new, push
forall stk: C, e: E
top ( push ( stk, e ) ) == e
pop ( push ( stk, e ) ) == stk
isEmpty ( new)
isEmpty ( new)
isEmpty ( push ( stk, e ))
implies

LinearContainer (push for insert, top for first, pop for rest)
```

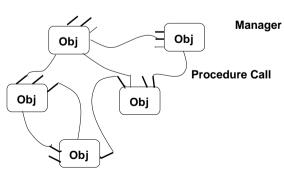
- pop (top (new)) ==
- isEmpty (pop (push (new, 5))) ==
- first (rest (insert (insert (new, 10), 11))) ==

/* Assume new for new */

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ADTs as Objects

- Does Fortran/COBOL/C/Pascal offer subroutines for operations? Does Fortran/COBOL/C/Pascal offer information hiding?
- Objects (also called "managers") are responsible for
 - preserving the integrity of their resources &
 - iding the representations from other objects
 - interactions take place mostly thru function and procedure invocations



- encourages reuse (<- hiding & encapsulation)
- interaction via procedure call
 - a (client) object should know the identity of another (server) object (cf. pipe&filter)
 - change of an object identity necessitates modification of all other objects that explicitly invoke it

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