

## Classical MILs

Why MILs?

Main Concepts of MILs

Trend in Industry

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## Why Module Interconnection Languages (MILs)?

*"Programming-in-the-small vs. programming-in-the-large"*

	<i>Programming-in-the-small</i>	<i>Programming-in-the-small</i>
<b>Goal</b>	building "programs"	building "software"
<b>Problem</b>	Usually clear, small	Usually unclear, large
<b>Emphasis</b>	Detailed design & impl.	sw. architecture
<b>Technique</b>	Structured programming	"divide & conquer" "separation of concerns"
<b>Notation</b>	PLs	(Formal) (OO) Specification Lang.
<b>Manpower</b>	single person/small number	multi-person
<b>Version</b>	usually single	multi-version

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## Why Module Interconnection Languages (MILs)?

*Programming-in-the-large requires conquering complexity!*

☆ *Division of work*

Only the owner team needs to know how to implement a particular part

☆ *Multi-paradigm implementation*

Different people are good at different PLs  
Different PLs are good for different things  
Different things are developed at different times

☆ *Evolvable software*

Impact of changes should be localized  
change in data structure or algorithm should be hidden  
change in PLs should be hidden (or localized -> wrapper)

☆ *Information protection*

Only on need-to-know basis

☆ *Reuse of components (in the library)*

Reduce development & verification effort

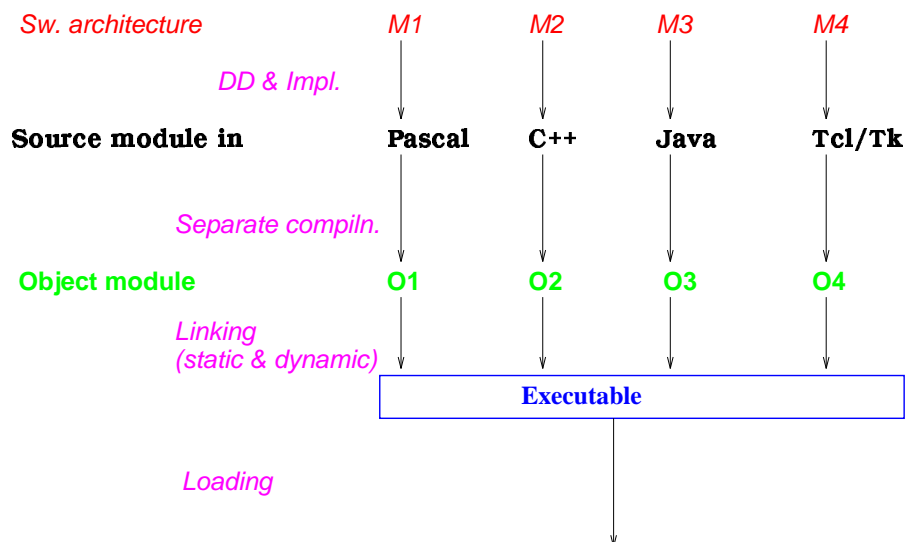
☆ *Separate compilation*

Can't compile 1M LOC for each change

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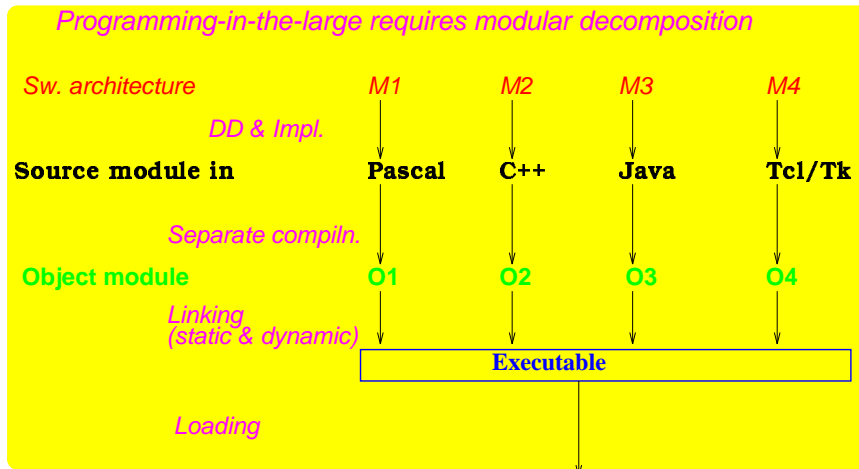
## Why Module Interconnection Languages (MILs)?

*Programming-in-the-large requires modular decomposition*



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## Why Module Interconnection Languages (MILs)?



### But

- ✂ **static type-checking & consistency checking at an intermediate level of descr.**  
*e.g., M1 uses a variable V in M3*  
*Is V defined in M3?*  
*Is M1 allowed to access V?*
- ✂ **controlling different versions, assembling components for a complete system**

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## Main Concepts of MILs

**MILs provide formal grammar constructs for various module interconnection specifications for assembling a complete software system.**

- ☛ **The first MIL was developed in 1975**  
["Programming-in-the-Large versus Programming-in-the-Small",  
DeRemer & Kron, IEEE TSE 2(1), June, 1976]
- ☛ **Variations among different schemes**  
["Module Interconnection Languages",  
Prieto-Diaz & Neighbors, The Journal of Systems and Software 6, 1986]
- ☛ **Module structure called "System Tree"**
  - ↳ modules that provide/export/synthesize resources and require/import/inherit them
  - ↳ a resource is any entity that can be named in a PL (e.g., variables, constants, procedures, type defs)
  - ↳ interface-oriented, without details of how functions or modules are implemented

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## Main Concepts of MILs

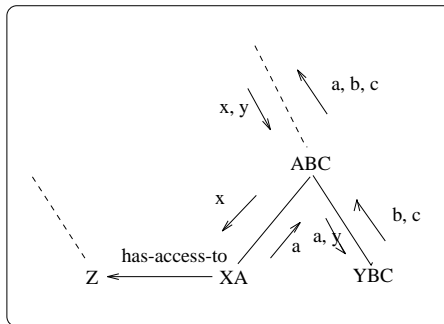
### Example

```

module ABC
  provides: a, b, c /* resources defined in ABS - "statement of origin" */
  requires: x, y /* resources used but not in ABS - "statement of usage" */
  consist-of: function XA, module YBC /* nesting of module */

  function XA
    provides: a
    requires: x
    has-access-to: module Z /* any resource provided by Z */
    real x, integer a
  end XA

  module YBC
    provides: b, c
    requires: a, y
    real y, integer a, b, c
  end YBC
end ABC
  
```



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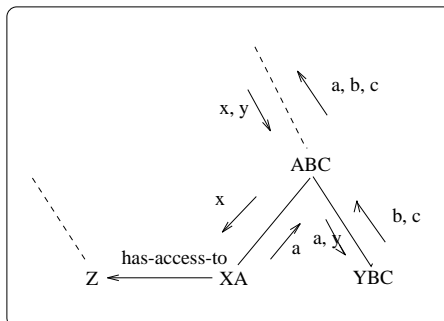
## Main Concepts of MILs

### Some constraints on accessibility

- ✂ has-access-to is not transitive  
e.g.,  $L \leftarrow \text{-----} M \leftarrow \text{-----} N$  does not mean  $L \leftarrow \text{-----} N$
- ✂ inheritance can be all (e.g., XA can access both x, y) or restricted (e.g., x only)

### Questions

- ✂ If  $L \leftarrow \text{-----} M$  and  $M'$  is a child of  $M$ , can  $M'$  access  $L$ ?
- ✂ If  $L \leftarrow \text{-----} M$  and  $L'$  is a child of  $L$ , can  $M$  access  $L'$ ?



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## Main Concepts of MILs

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☛ **systems supporting module interconnection include:**

☞ **Ada (package), Module (module)**

highly modular and provide for version definitions

☞ **Protel (PROcedure Oriented Type Enforcing Language)**

implemented in 1975 by BNR  
 used extensively but mainly by BNR  
 based on compile-link-load paradigm  
 performs type checking across modules

☞ **SCCS (Source Code Control System)**

part of PWB (Programmer's Work Bench) facility  
 by Bell Labs in 1973  
 a file storage system for recording various versions of a text file  
 supports creation of any revision of a source program or text  
 file protection against accidental changes

☞ **SARA (System ARchitect's Apprentice)**

supports a structural multi-level requirements-driven methodology  
 for the design of reliable sw or hw digital systems  
 developed at UCLA in 1976 and under continual development

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## Trend in Industry

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\* **"buy, don't build"**

[Brooks, "No silver bullet: Essence and Accidents of Software Engineering",  
 Computer 20(4), pp. 10-19, Apr. 1987]

**faster (reduced development time)**

**increased reliability**

**increased flexibility**

\* **increasing component size and complexity**

**e.g1., TPS at GTE (1000 small, 50 large)**

**e.g2., F-22 fighter aircraft**

OS      network mgmt system  
 ballistic eqns for free-ball bombs  
 navigation algorithms  
 process scheduler    compilers    dbms  
 aircraft-specific delays  
 UIM

aircraft-specific delays
ballistic eqns for free-ball bombs
process scheduler navigation algorithms UIM
dbms
OS compilers network mgmt system

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## Trend in Industry

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\* ***coordination among components***

*coordination infrastructure & standards for components plug-in*  
e.g., CORBA

\* ***from subroutines to subsystems***

*["Component-based Software Development",  
American Programmer 8(11), Nov. 1995]*

\* ***architecture specification as a deliverable***

*"If a project has not achieved a system architecture, including its rationale,  
the project should not proceed to full-scale system development." [Boehm]*

*increasing importance of sw. architecture and specification*