Computer Science Program, The University of Texas, Dallas

Client-Server Architecture

Clients and Servers Client/Server with File Servers Client/Server with Database Servers Client/Server Communication Client/Server with Transaction Processing Client/Server Groupware Web Client/Server

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Paradigm Shift: Past, Present and Future

Centralized processing

A host computer (often a mainframe) handls all processing, including input, output, data storage and retrieval

Predominant computing mode in the `70s

Distributed processing

A number of computers (minis, workstations, PCs, ...) handle all processing, They are distributed physically and connected thru a communications network

The computing mode of the present

Cooperative processing

A number of computers (minis, workstations, PCs, ...) handle all processing, They are distributed physically and connected thru a communications network

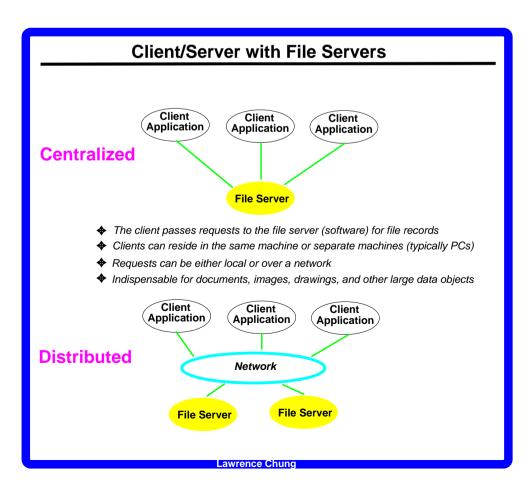
Processing thru sharing of resources, transparently to the users

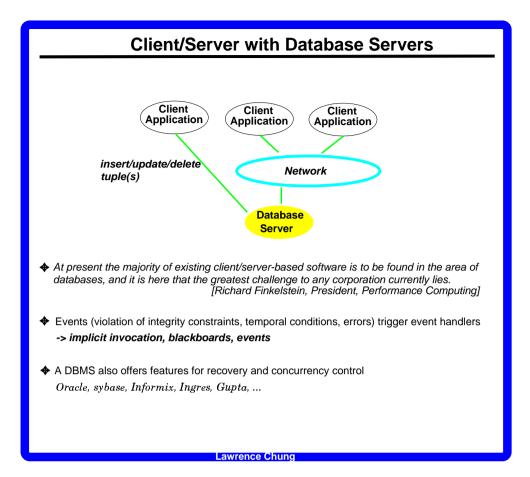
The computing mode of the future

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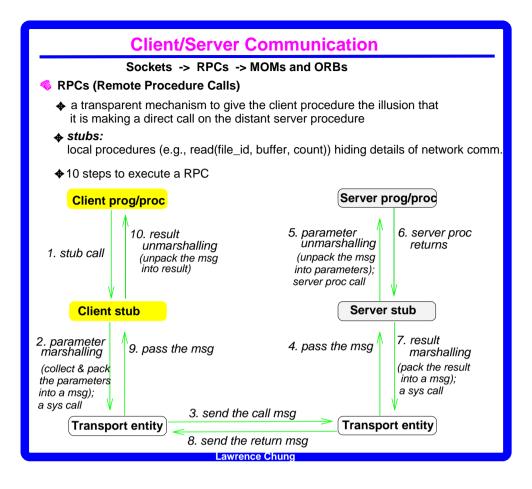
Clients and Servers Star. **Basic Definition** Client: requests for services ♣ Server: provides services ♣ Service: any resource * (e.g., data, type defn, file, control, object, CPU time, display device, etc.) **Typical Properties** S. A service request is about "what" is needed, and it often made abstractly It is up to the server to determine how to get the job done -> the notion of module (cf. MILs, ADTs --- Larch, Z) ☑ The ideal client/server software is independent of hardware or OS platform ☑ The location of clients and servers are usually transparent to the user A client may become a server; a server may become a client A client/server system can be scaled horizontally, i.e., by adding/removing client workstations with only a slight performance impact vertically. i.e., by migrating to a larger and faster server machines or multiservers

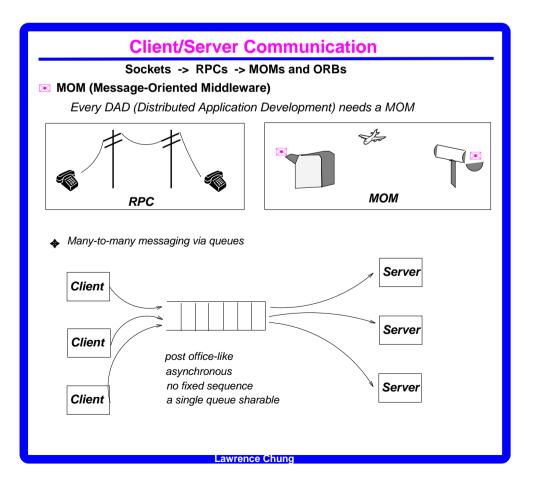
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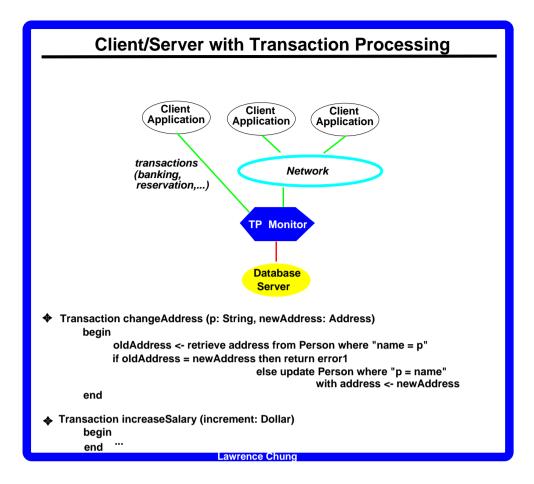




Client/Server Communication		
Sockets -> RPCs -> MOMs and ORBs		
Sockets		
♦ introduced in 1981 as the Unix BSD 4.2 generic interface		
provide Unix-to-Unix comr	nunications over networks	
•	/insock): based on the Unix BSD 4.2 sockets interface	
Net_id.Host_id Internet Address	Port Address	
unique Net id multiple Host id	Ports (entry points to applns on a host)	
	connection-oriented/connectionless, protocol:TCP/IP	
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Client/Server with Transaction Processing

Transactions are a way to make ACID operations a general commodity [Transaction Processing Concepts and Techniques, Jim Gray and Andreas Reuter, 1993]

Atomicity

- a transaction is an indivisible unit of work
- □ an all-or-nothing proposition
- all updates to a database, displays on the clients' screens, message queues
- Le.g., salary increase for all 1 million employees or none

Consistency

- a transaction is an indivisible unit of work
- □ S -> [T | abort] -> S
- □ integrity constraints (e.g., mgr.salaray > salary)

Isolation

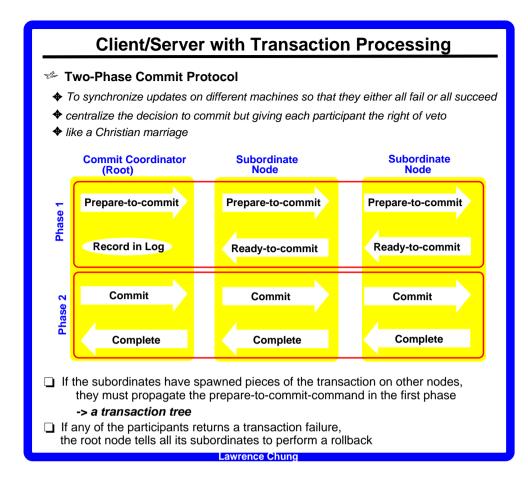
- a transaction's behavior not affected by other transactions running concurrently
- Le.g., reserve a seat
- serialization techniques

🗖 Durability

- persistence
- a transaction's effects are permanent after it commits.

ACID is like motherhood and apple pie. It's necessary- and you can't have too much of it. OK, enough preaching.

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Client/Server Groupware	
✦ Why?	
support for business reengineering: bcfh -> maximize profit	
change the way people communicate with each other helps manage (and track) the product thru its various phases	
computer-supported cooperative work (CSCW)	
collaborative/workgroup computing	
return on investment: 16% to 1666% on a median investment of \$100K [a study of 65 Notes users in 1994]	
Five foundation technologies	
"combining technologies and creating new synergy"	
 Multimedia document management from electronic imaging (scanning, digitization, display, storage and retrieval) to document (component types: text, image, graphics, faxes, mail, voice clips, BBs) Workflow 	
🛿 E-mail	
Conferencing thru microphones, video windows, electronic whiteboards, controlled access to shared document	
Scheduling	
electronic scheduling of meetings, sharing calendars and "to do" lists, triggering events	
Why not just DBMS?	
DBMS: highly structured data Groupware: highly unstructured data document database	
Why not just TP Monitors?	
TP Monitors: transaction processes Groupware: (currently) not transaction-oriented in the ACID sense	
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