

Theory of Computational Complexity

by Ding-Zhu Du and Ker-I Ko

Errata

Page 21, line 4: “ $A \in DTIME(c \cdot t(n))$ ” should read “ $L(M) \in DTIME(c \cdot t(n))$ ”

Page 47, line 11 from the bottom: “ EXP ” should read “ $EXPPOLY$ ”

Page 56, line 13: Both occurrences of “ $x[j]$ ” should be changed to “ $x[k]$ ”

Page 67, line 24: Delete “at most $3K_T/2$ and, hence,”

Page 87, line 17: Add the following sentence at the end:

We also let 3-CNF-TAU_k (3-DNF-TAU_k) denote the problem of determining whether a given 3-CNF (3-DNF, respectively) formula F satisfies

$$(\forall \tau_1 : X_1 \rightarrow \{0, 1\})(\exists \tau_2 : X_2 \rightarrow \{0, 1\}) \cdots \\ (Q_{k+1} \tau_k : X_k \rightarrow \{0, 1\}) F|_{\tau_1, \tau_2, \dots, \tau_k} = 1,$$

where $Q_k = \exists$ if k is odd, and $Q_k = \forall$ if k is even.

Page 87, Corollary 3.14: The second half of part (a) should read

“and 3-DNF-TAU_k is \leq_m^P -complete for Π_k^P .”

The second half of part (b) should read

“and 3-CNF-TAU_k is \leq_m^P -complete for Π_k^P .”

Page 110, Exercise 3.13: This exercise should be changed to the following:

3.13 Prove that the game GEOGRAPHY played on undirected graphs is solvable in polynomial time.

Page 114, line 8 from the bottom: “In particular, $|S_{n-1}| = \varphi(n-1) > 0$ ” should be replaced by the following paragraph (in the following, $a|b$ denotes that a divides b):

Since, for all d that do not divide $n-1$, we have $S_d = \emptyset$, we get $\sum_{d|(n-1)} |S_d| = |Z_n^*| = n-1$. In the above, we have proved that if d divides $n-1$, then either $S_d = \emptyset$ or $|S_d| = \varphi(d)$. Thus, by Euler’s Theorem, for each d that divides $n-1$, we must have $|S_d| = \varphi(d)$, for otherwise $\sum_{d|(n-1)} |S_d|$ would be less than $n-1$. It follows that $|S_{n-1}| = \varphi(n-1) > 0$.

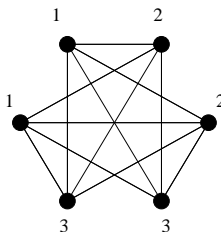
Page 115, line 18: “less than $n-1$ ” should read “less than or equal to $(n-1)/2$ ”

Page 121, line 8 from the bottom: “approximation to f ” should read “approximation to f^{-1} ”

Page 204, line 13: “ $\langle 0, w \rangle \neq B$ ” should read “ $\langle 0, w \rangle \notin B$ ”

Page 204, line 16: “then $\langle \psi, w \rangle \in B$ ” should read “then $\langle \phi, w \rangle \in B$ ”

Page 206, Figure 6.5: The figure at the right-hand side should be changed to



Page 218, line 1: “Each subcircuit C_i ” should read “Each subcircuit C_{i+1} ”

Page 239, line 9: “ $\mathcal{D} = P^{\mathcal{C}}$ ” should read “ $\mathcal{D} = P^{\mathcal{C}}$ for some class \mathcal{C} that contains P as a subclass”

Page 258, line 22: “there is a co-sparse set” should read “there is no co-sparse set”

Page 258, line 23: “is also em no” should read “is also NO”

Page 291, Lemma 8.5: Add two more items to Lemma 8.5:

(c) For all $x, y > 1$ such that $\gcd(x, n) = \gcd(y, n) = 1$, $(\frac{x}{n}) \cdot (\frac{y}{n}) = (\frac{xy}{n})$.

(d) For all odd $n > 2$, $(\frac{2}{n}) = (-1)^{(n^2-1)/8}$.

Page 291, line 14 from the bottom: “Lemma 8.5(b) above shows” should read “Lemma 8.5 above shows”

Page 306, line 12: “ $F|_{x=1} \in \text{SAT}$ ” should read “ $F|_{x_1=1} \in \text{SAT}$ ”

Page 333, line 11: “a cover of H_k ” should read “the path (u, u_1, v_1, v) and a cycle cover of $H_k - \{u, u_1, v_1, v\}$ ”

Page 421: The first half of the proof of Theorem 11.24 has an error and should be changed as follows:

Lines 26 to 27: Replace “We assume” by “We assume that Φ has $k(h+1)^m$ variables and $(h+1)^m$ clauses. In addition, we assume”

Line 27: Replace “where Π ” by “where each f_i , $1 \leq i \leq k$, defines Boolean assignments to $(h+1)^m$ variables, and Π ”

Lines 31 to 33: Replace the two sentences “Then, V_k simulates $V \dots$ uses $f_i(\vec{a})$ for it.” by “Then, V_k simulates V with the following modification:

It uses Π_{k+1} as the additional proof, and replaces $\bar{G}^f(\vec{u}_1, \vec{u}_2, \vec{u}_3, \vec{v})$ by

$$\bar{G}_1^{f_1, \dots, f_k}(\vec{u}_1, \vec{u}_2, \vec{u}_3, \vec{v}) = \prod_{j=1}^3 \prod_{i=1}^k \bar{\theta}_{ij}(\vec{u}_j, \vec{v})(\bar{\varphi}_j(\vec{v}) - f_i(\vec{u}_j)),$$

where $\theta_{ij}(s, t)$ is equal to 1 if the s th variable in the i th block is the j th variable in C_t , and equal to 0 otherwise. Note that the sum check proof system with respect to function $\bar{G}_1^{f_1, \dots, f_k}$ now requires three values of $f_i(s)$ for $1 \leq i \leq k$ and $s \in \mathcal{F}^m$.”

Page 432, line 6 from the bottom: “and F has” should read “and F_x has”

Page 432, line 5 from the bottom: “total of $d_1 n^c$ clauses).” should read “total of $d_1 n^c$ clauses) if $x \notin L$.”

Page 436, line 13: Delete the period at the end of the formula for S' .

Page 465, lines 26 to 28: Change this reference item to

Kozen, D. [1977], Lower bounds for natural proof systems, *Proceedings of the 18th IEEE Symposium on Foundations of Computer Science*, IEEE Computer Society Press, Los Angeles, pp. 254–266.

Updated June 30, 2006