Lecture 2: OCaml Functions

CS 6371: Advanced Programming Languages January 16, 2014

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\#let solve w x y z =
                                                    Top-level and nested "let" operations are
  let prod a b = a*b in
                                                    different: A top-level "let" defines a global
    (prod w x) + (prod y z);
                                                    variable (usually a function). It must be
solve : int -> int -> int -> int =
                                                    followed by ";;". An inner "let" declares a local
#prod;;
                                                    variable that can be used only within a
Toplevel input:
                                                    subexpression. It must be followed by "in".
>prod;;
>^^^
                                                    When variable names conflict, the innermost
The value identifier prod is unbound.
                                                    declaration "shadows" the others.
#type foo = int;;
                                                    User-defined types can be primitive types (int,
#type foo = bool;;
                                                    bool, string, etc.), or they can be lists, tuples, or
#type foo = string;;
                                                    variants that include any of the above.
#type foo = int list;;
#type foo = int * string;;
#type btree = BLeaf
             | BNode of (int*btree*btree);;
#type ntree = NLeaf
             | NNode of ((int*ntree) list);;
#let identity x = x;
                                                    When possible, OCaml gives functions a
identity : 'a -> 'a = <fun>
                                                    polymorphic type. Polymorphic functions can
#identity 3;;
                                                    be applied to arguments of any type.
- : int = 3
#identity "foo";;
- : string = "foo"
\#let apply f x = (f x);;
                                                    However, there must be some consistent way
apply : ('a -> 'b) -> 'a -> 'b = <fun>
                                                    to instantiate each type variable. Here we see
\#let add (x,y) = x+y;;
                                                    an example where no such instantiation exists
add : int * int -> int = <fun>
#apply add (1,2);;
                                                    and the compiler therefore rejects the code.
- : int = 3
#apply add "foo";;
Toplevel input:
>apply add "foo";;
This expression has type string,
but is used with type int * int.
#let rec map f l =
                                                    Lists can also have polymorphic type.
  (match 1 with
     [] -> []
   | x::t -> (f x)::(map f t));;
map:('a -> 'b) -> 'a list -> 'b list = <fun>
\#let addone n = n+1;;
addone : int -> int = <fun>
#map addone [23;42;64];;
-: int list = [24; 43; 65]
\#map (fun n \rightarrow n+1) [23;42;64];;
                                                    Use "fun" to create anonymous (i.e., unnamed)
-: int list = [24; 43; 65]
                                                    functions. "fun ... -> ..." is the same as if you
\#(fun n -> n+1);;
                                                    typed "let foo ... = ...;;" and then used "foo".
- : int -> int = <fun>
\#(fun n -> n+1) 2;;
- : int = 3
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#let compose f g = (fun x \rightarrow f (g x));
                                                     Using anonymous functions, you can build and
compose : ('a->'b)->('c->'a)->'c->'b = <fun>
                                                     return functions as values at runtime.
\#let cool = (compose (fun n -> n+1)
                        (fun n -> n*2));;
cool : int -> int = <fun>
#cool 1;;
- : int = 3
\#let addx x = (fun y \rightarrow x+y);;
                                                     An anonymous function may refer to variables
addx : int -> int -> int = <fun>
                                                     declared in outer scopes.
#addx 1;;
- : int -> int = <fun>
#(addx 1) 2;;
- : int = 3
#addx 1 2;;
-: int = 3
\#let add = (fun x -> (fun y -> x+y));;
                                                     Actually "let foo x y = ..." is just an abbreviation
add : int -> int -> int = <fun>
                                                     for "let foo = (\text{fun x} \rightarrow (\text{fun y} \rightarrow ...))". If you give
#add 1 2;;
                                                     such a function fewer arguments than it
-: int = 3
#add 1;;
                                                     expects, it yields a function from the remaining
- : int -> int = <fun>
                                                     arguments to the original value. Functions
\#let add x y = x+y;;
                                                     written this way are called "curried functions".
add : int -> int -> int = <fun>
#map (add 3) [1;2;3];;
                                                     Applying fewer arguments is called "partial
-: int list = [4; 5; 6]
                                                     evaluation".
#let rec map2 f 11 12 =
                                                     Binary operators can be used in prefix rather
  (match (11,12) with
                                                     than infix syntax by enclosing the operator in
     ([],x) \rightarrow x
                                                     parentheses. This allows you to pass a binary
     (x,[]) \rightarrow x
   | (h1::t1, h2::t2) ->
                                                     operator as a function argument.
         (f h1 h2)::(map2 f t1 t2));;
map2 : ('a -> 'a -> 'a) -> 'a list -> 'a
list -> 'a list = <fun>
#map2 add [1;2;3] [10;20;30];;
- : int list = [11; 22; 33]
#map2 (+) [1;2;3] [10;20;30];;
-: int list = [11; 22; 33]
#let rec addpairs l =
                                                     Typing annotations are almost never necessary,
  (match 1 with
                                                     but they can help you debug.
     x::y::t \rightarrow (add x)::(addpairs t)
    [] -> []
    [x] -> [x]);;
Toplevel input:
  [x] -> [x]);;
              ^ ^ ^
This expression has type int list,
but is used with type (int -> int) list.
# let rec addpairs (1:int list) : int list =
    (match 1 with
       x::y::t ->
          (add x)::(addpairs t)
      | [] -> []
     | [x] -> [x]);;
Toplevel input:
   x::y::t \rightarrow (add x)::(addpairs t)
This expression has type int list but is
here used with type
  (int -> int) list
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<pre>#let intident (x:int) = x;; intident : int -> int = <fun></fun></pre>	You can also use a typing annotation to restrict the type of a function that would otherwise be polymorphic.
<pre>#let apply (f:'inp->'out) (x:'inp) = (f x);; apply : ('a -> 'b) -> 'a -> 'b = <fun></fun></pre>	Typing annotations may contain type variables.
<pre>#let main () = (print_string "hello\n"; print_int 3; print_newline (); 12);; main : unit -> int = <fun> #main ();; hello 3 - : int = 12</fun></pre>	A ;-separated sequence of expressions is evaluated in order. The last expression is returned as the result of the sequence expression. Use sequence expressions with print statements to debug.
<pre>#"foo" = "foo";; - : bool = true #(3,"foo") = (3,"foo");; - : bool = true #[1;2;3] = [1;2;3];; - : bool = true #Dark (Dark Red) = Dark (Dark Red);; - : bool = true</pre>	OCaml's equality operator (=) tests structural equality. This means that you can use it with ints, bools, strings, tuples, lists, and variants.