Evaluation Strategies CS 4301/6371: Advanced Programming Languages

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April 11-16, 2024

Definition (first-class): A type is said to be *first-class* for a programming language if values of that type require no special syntax or encapsulation to be

- assigned to variables,
- passed as arguments,
- returned by functions,
- any other type-agnostic usages.

Which of the following languages have first-class functions?

C C++ SIMPL Java JavaScript Python $\lambda \rightarrow$ System F OCaml Haskell

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Definition (Curried): A multi-argument function is *curried* if it is expressed as a function from each individual argument to a function of the remaining arguments (i.e., has type $\tau_1 \rightarrow \cdots \rightarrow \tau_n$).

Definition (Partial Evaluation): A multi-argument function is *partially evaluated* when it is applied to fewer than its total number of arguments, yielding a function from the remaining arguments to the return value.

Which of the following languages support currying and partial evaluation?

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Eager Evaluation

Definition (Eager Semantics or Call-by-value): An *eager* or *call-by-value* language evaluates all function arguments before passing them as parameters.

Operational semantics look like this:

$$\frac{\forall i \in [1, n], \langle e_i, \sigma \rangle \Downarrow u_i \qquad \sigma(f)(u_1, \dots, u_n) \Downarrow u'}{\langle f(e_1, \dots, e_n), \sigma \rangle \Downarrow u'}$$

Lazy Evaluation

Definition (Lazy Semantics): A *lazy* language evaluates function arguments after the function body has started evaluating. There are two main varieties:

- **Call-by-name** languages (re)evaluate each argument expression each time the function uses it.
 - Can be formalized via capture-avoiding substitution
 - Disadvantage: usually inefficient
 - Advantage: sometimes highly efficient (e.g., unused arguments, highly parallelizable languages)
- **Call-by-need** languages evaluate each argument at first use, then memoize and reuse those values at subsequent uses.
 - Advantage: highest efficiency (usually)
 - Disadvantage: sometimes unintuitive!

Optional Exercises: Devise call-by-value and call-by-need operational semantics for $\lambda\text{-calculus}$

Call-by-reference

Definition (call-by-reference): Languages supporting *call-by-reference* allow callees to destructively modify the values of variables passed as arguments.

Example: Most object-oriented languages pass objects by reference, allowing callees to globally modify the object's fields instead of receiving a local copy of the object.

Note: Call-by-reference does not make sense for immutable variables.

Which evaluation strategies are supported by the following languages?

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С	call-by-value
C++	-
SIMPL	
Java	
JavaScript	
Python	
λ_{\rightarrow}	
System F	
OCaml	
Haskell	

С	call-by-value
C++	call-by-value, call-by-reference
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Java	call-by-value, call-by-reference (objects)
JavaScript	
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λ_{\rightarrow}	
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Evaluation Strategies

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Church-Rosser Property

Definition (Church-Rosser): Languages with the *Church-Rosser Property* are those in which the order of evaluation has no impact on the observable result. More technically, they are those languages whose small-step operational semantics are *confluent*.

Church-Rosser languages typically...

- have strictly immutable variables,
- are *pure* (i.e., free of side-effects).

Languages that are Church-Rosser can have unknown evaluation strategies (unobservable to the user), and offer compilers many optimization opportunities.

Static vs. Dynamic Typing

Definition (static/dynamic typing): A language is *(strictly) statically typed* if all types are erased during compilation. In contrast, a language is *dynamically typed* if types are available at runtime (usually attached to runtime values).

Advantages of strict static typing:

- space- and time-efficiency (no runtime storage or tracking of types)
- types facilitate static debugging
- types facilitate compile-time static code optimization
- types can be more universal (e.g., characterizing all possible executions)

Advantages of dynamic typing:

- type-tag values available at runtime (whether you need them or not)
- sometimes easier patching and bug mitigation
- opportunities for extra security sanity-checking

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Definition (type-safety): A language is *type-safe* if its static semantics preclude all stuck states in its operational semantics.

Sometimes difficult to tell whether a language is type-safe because:

- Some languages have no formal semantics(?!?!).
- Some languages have an operational semantics that formalizes states most of us would consider stuck states.

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C++ ×

SIMPL ✓

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System F ✓

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Polymorphism

Definition (polymorphism): A language is *polymorphic* if interfaces (e.g., functions) can accommodate entities (e.g., arguments) of multiple different types.

Three main varieties:

I Parametric Polymorphism: type system has type-variables α

- facilitates machine-checked code-reuse idioms
- compatible with strictly static type-safety

2 Subtyping Polymorphism: object types arranged in a hierarchy

- hallmark of object-oriented programming
- static semantics usually characterized by a weakening rule:

$$\frac{\Gamma \vdash e: \tau \quad \tau \preceq \tau'}{\Gamma \vdash e: \tau'}$$

Warning: makes structural induction proofs much harder (Why?)

3 Ad hoc Polymorphism: conditionals can test types at runtime

opens the door for arbitrarily heterogeneous code blocks per type

antithesis of code-reuse (much harder to maintain and debug)

Which forms of polymorphism are supported by the following languages?

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Java	parametric (generics), subtyping, ad hoc
JavaScript	
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OCaml	
Haskell	

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C++	subtyping
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$\lambda_{ ightarrow}$	
System F	
OCaml	
Haskell	

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System F	
OCaml	
Haskell	

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C++	subtyping
SIMPL	none
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Python	parametric (generics), subtyping, ad hoc
λ_{\rightarrow}	none
System F	parametric
OCaml	
Haskell	

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Polymorphism Examples

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Python	parametric (generics), subtyping, ad hoc
λ_{\rightarrow}	none
System F	parametric
OCaml	parametric, subtyping
Haskell	parametric

Definition (shallow types): A shallowly-typed language is one whose type system only supports type quantifiers at the top level of types (not nested within non-quantifiers).

Which of the following languages support non-shallow types:

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$$\lambda_{\rightarrow}$$

System F
OCaml
Haskell

Definition (shallow types): A shallowly-typed language is one whose type system only supports type quantifiers at the top level of types (not nested within non-quantifiers).

C ×
C++ ×
SIMPL ×
Java ×
JavaScript ×
Python ×

$$\lambda_{\rightarrow}$$

System F
OCaml
Haskell

Definition (shallow types): A shallowly-typed language is one whose type system only supports type quantifiers at the top level of types (not nested within non-quantifiers).

C X
C++ X
SIMPL X
Java X
JavaScript X
Python X
$$\lambda_{\rightarrow}$$
 X
System F
OCaml
Haskell

Definition (shallow types): A shallowly-typed language is one whose type system only supports type quantifiers at the top level of types (not nested within non-quantifiers).

C ×
C++ ×
SIMPL ×
Java ×
JavaScript ×
A
$$\rightarrow$$
 ×
System F \checkmark
OCaml
Haskell

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$$\begin{array}{cccc} & \swarrow & \\ \mathsf{C} + + & \checkmark & \\ \mathsf{SIMPL} & \checkmark & \\ \mathsf{Java} & \checkmark & \\ \mathsf{JavaScript} & \leftthreetimes & \\ \mathsf{Python} & \bigstar & \\ \mathsf{\lambda}_{\rightarrow} & \checkmark & \\ \mathsf{System F} & \checkmark & \\ \mathsf{OCaml} & \checkmark (\mathsf{with --rectypes}) \\ \mathsf{Haskell} & \\ \end{array}$$

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$$\begin{array}{cccc} & \swarrow & \\ C++ & \checkmark & \\ SIMPL & \checkmark & \\ Java & \checkmark & \\ JavaScript & \leftthreetimes & \\ A\rightarrow & \checkmark & \\ System F & \checkmark & \\ OCaml & \checkmark (with --rectypes) \\ Haskell & \checkmark & \end{array}$$

Summary Table

