

# CS 6363.003 Homework 2

Due Sunday March 7th on eLearning

February 21, 2021

Please answer the following 4 questions, some of which have multiple parts.

- Suppose we are given two sorted arrays  $A[1..n]$  and  $B[1..n]$ . Describe an algorithm to find the median element (the element of rank  $n$ ) in the union of  $A$  and  $B$  in  $O(\log n)$  time. You may assume that the arrays contain no duplicate elements. [Hint: Compare  $A[\lfloor n/2 \rfloor]$  and  $B[\lfloor n/2 \rfloor]$ . How can you reduce your search space to two sorted arrays of size  $\lfloor n/2 \rfloor$ ?]
  - Now suppose we are given two sorted arrays  $A[1..m]$  and  $B[1..n]$  with no duplicate elements and an integer  $k$  where  $1 \leq k \leq m+n$ . Describe an algorithm to find the  $k$ th smallest element in  $A \cup B$  in  $O(\log(m+n))$  time. [Hint: Now compare  $A[\lfloor m/2 \rfloor]$  and  $B[\lfloor n/2 \rfloor]$  but only reduce the search space in one of the two arrays.]
- Suppose you are given a set  $P = \{p_1, \dots, p_n\}$  of  $n$  points in the plane, represented by two arrays  $X[1..n]$  and  $Y[1..n]$ . Specifically, point  $p_i$  has coordinates  $(X[i], Y[i])$ . A point  $p_i \in P$  is called **Pareto optimal** if there exists no point  $p_j \in P$  with  $i \neq j$  such that both  $X[j] \geq X[i]$  and  $Y[j] \geq Y[i]$ . In other words, each Pareto optimal point of  $P$  has no other point of  $P$  both above and to its right.

Describe and analyze an  $O(n \log n)$  time algorithm to output the set of Pareto optimal members of  $P$ . (Any reasonable output describing these points is fine; for example, you could output an array  $Z[1..h]$  where each element of  $Z$  is the index  $i$  of a Pareto optimal point  $p_i$ .) [Hint: Use divide-and-conquer.]

- Suppose we are given an array  $A[1..n]$  of numbers, which may be positive, negative, or zero, and which are **not** necessarily integers. We are going to design a dynamic programming algorithm that finds the largest sum of elements in a contiguous subarray  $A[i..j]$ . For example, if we are given the array  $[-6, 12, -7, 0, 14, -7, 5]$ , our algorithm should return 19 for the contiguous subarray  $A[2..5]$ . Given the one-element array  $[-374]$  as input, our algorithm should return 0 (the empty interval is still an interval!) For the sake of analysis, we'll assume that comparing, adding, or multiplying any pair of numbers takes  $O(1)$  time.
  - Unless it is empty, the maximum contiguous subarray must consist of some *last* element  $A[j]$  along with 0 or more elements preceding  $A[j]$ . Accordingly, let  $maxSum(j)$  equal the largest sum of elements in a contiguous subarray of  $A[1..j]$  whose *last member* is  $A[j]$ .  
Give a recursive definition for  $maxSum(j)$ . Don't forget the base cases!

- (b) What would be the running time of a dynamic programming algorithm that computes  $\text{maxSum}(j)$  for all  $j$  from 1 to  $n$  using your recursive definition? [Hint: You should be able to answer this question without having to describe an iterative algorithm.]
- (c) Describe and analyze an efficient algorithm that finds the largest sum of elements in a contiguous subarray of  $A[1 .. n]$ . [Hint: Use parts (a) and (b).]
- (d) Now suppose in addition to  $A[1 .. n]$ , you are given an additional integer  $X \geq 0$ . Describe and analyze an algorithm that finds the largest sum of elements in a contiguous subarray of  $A$  whose length is at most  $X$ . [Hint: You'll want to start over by slightly modifying the function from part (a).]
4. For each of the following problems, the input consists of two arrays  $X[1 .. k]$  and  $Y[1 .. n]$  where  $k \leq n$ .
- (a) Describe and analyze an algorithm to decide whether  $X$  is a subsequence of  $Y$ . For example, the string **PPAP** is a subsequence of the string PENPINEAPPLEAPPLEPEN.
- (b) Suppose the input also includes a third array  $C[1 .. n]$  of numbers, which may be positive, negative, or zero, where  $C[i]$  is the *cost* of  $Y[i]$ . Describe and analyze an algorithm to compute the minimum cost of any occurrence of  $X$  as a subsequence of  $Y$ . That is, we want to find the minimum total cost  $\sum_{j=1}^k C[I[j]]$  among all arrays  $I[1 .. k]$  such that  $I[j] < I[j + 1]$  and  $X[j] = Y[I[j]]$  for every index  $j$ .
- (c) Describe and analyze an algorithm to compute the total number of (possibly overlapping) occurrences of  $X$  as a subsequence of  $Y$ . For purposes of analysis, assume we can add two arbitrary integers in  $O(1)$  time. For example, the string **PPAP** appears exactly 23 times as a subsequence of the string PENPINEAPPLEAPPLEPEN. If all characters in  $X$  and  $Y$  are equal, your algorithm should return  $\binom{n}{k}$ .