Data management in the cloud using Hadoop

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Outline

- Hadoop Basics
- HDFS
 - Goals
 - Architecture
 - Other functions
- MapReduce
 - Basics
 - Word Count Example
 - Handy tools
 - Finding shortest path example
- Related Apache sub-projects (Pig, Hbase, Hive)



Hadoop - Why ?

- Need to process huge datasets on large clusters of computers
- Very expensive to build reliability into each application
- Nodes fail every day
 - Failure is expected, rather than exceptional
 - The number of nodes in a cluster is not constant
- Need a common infrastructure
 - Efficient, reliable, easy to use
 - Open Source, Apache Licence version of Google File System

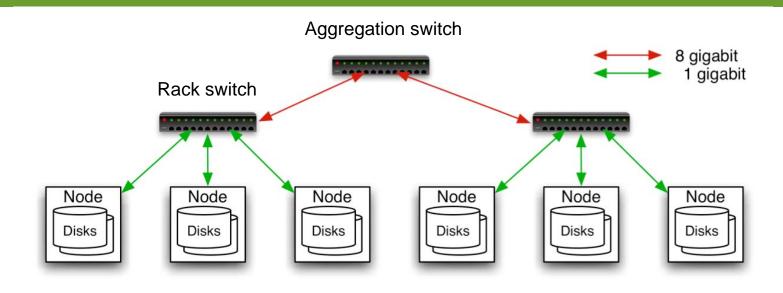


Who uses Hadoop?

- Amazon/A9
- Facebook
- Google
 - It has GFS
- New York Times
- Veoh
- Yahoo!
- many more
- Cloudera
 - Similar to Redhat business model.
 - Added services on Hadoop



Commodity Hardware



- Typically in 2 level architecture
 - Nodes are commodity PCs
 - 30-40 nodes/rack
 - Uplink from rack is 3-4 gigabit
 - Rack-internal is 1 gigabit



Hadoop Distributed File System (HDFS)

Original Slides by Dhruba Borthakur

Apache Hadoop Project Management Committee



Goals of HDFS

- Very Large Distributed File System
 - 10K nodes, 100 million files, 10PB
 - Yahoo is working on a version that can scale to large amounts of data.
- Assumes Commodity Hardware
 - Files are replicated to handle hardware failure
 - Detect failures and recover from them
- Optimized for Batch Processing
 - Data locations exposed so that computations can move to where data resides
 - Remember moving large data is an important bottleneck.
 - Provides very high aggregate bandwidth





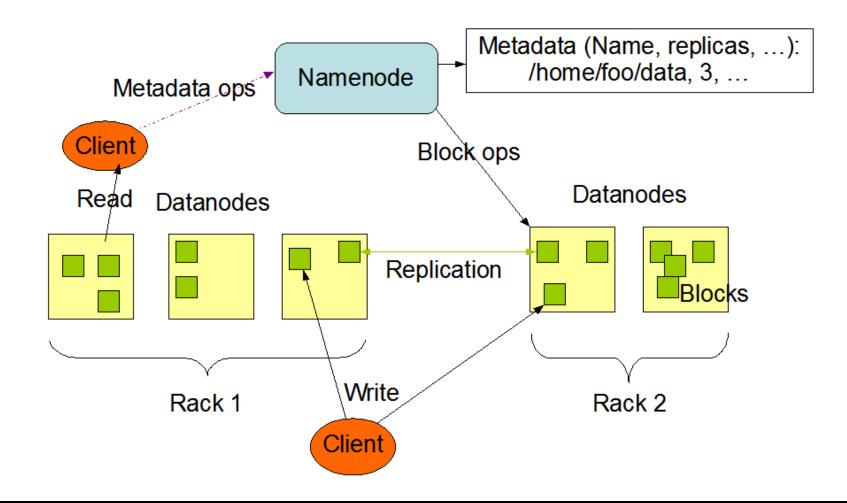
Distributed File System

- Single Namespace for entire cluster
 Again this is changing soon!!!
- Data Coherency
 - Write-once-read-many access model
 - Client can only append to existing files
- Files are broken up into blocks
 - Typically 64MB block size
 - Each block replicated on multiple DataNodes
- Intelligent Client
 - Client can find location of blocks
 - Client accesses data directly from DataNode



HDFS Architecture

HDFS Architecture





Functions of a NameNode

- Manages File System Namespace
 - Maps a file name to a set of blocks
 - Maps a block to the DataNodes where it resides
- Cluster Configuration Management
- Replication Engine for Blocks



NameNode Metadata

- Metadata in Memory
 - The entire metadata is in main memory
 - No demand paging of metadata
- Types of metadata
 - List of files
 - List of Blocks for each file
 - List of DataNodes for each block
 - File attributes, e.g. creation time, replication factor
- A Transaction Log
 - Records file creations, file deletions etc



DataNode

- A Block Server
 - Stores data in the local file system (e.g. ext3)
 - Stores metadata of a block (e.g. CRC)
 - Serves data and metadata to Clients
- Block Report
 - Periodically sends a report of all existing blocks to the NameNode
- Facilitates Pipelining of Data
 - Forwards data to other specified DataNodes



Block Placement

- Current Strategy
 - One replica on local node
 - Second replica on a remote rack
 - Third replica on same remote rack
 - Additional replicas are randomly placed
- Clients read from nearest replicas



Heartbeats

- DataNodes send hearbeat to the NameNode
 Once every 3 seconds
- NameNode uses heartbeats to detect DataNode failure



Replication Engine

- NameNode detects DataNode failures
 - Chooses new DataNodes for new replicas
 - Balances disk usage
 - Balances communication traffic to DataNodes



Data Correctness

- Use Checksums to validate data
 Use CRC32
- File Creation
 - Client computes checksum per 512 bytes
 - DataNode stores the checksum
- File access
 - Client retrieves the data and checksum from DataNode
 - If Validation fails, Client tries other replicas



NameNode Failure

- A single point of failure
- Transaction Log stored in multiple directories
 - A directory on the local file system
 - A directory on a remote file system (NFS/CIFS)
- Need to develop a real HA solution
 - Ongoing work to have multiple NameNodes



Data Pipelining

- Client retrieves a list of DataNodes on which to place replicas of a block
- Client writes block to the first DataNode
- The first DataNode forwards the data to the next node in the Pipeline
- When all replicas are written, the Client moves on to write the next block in file



Rebalancer

- Goal: % disk full on DataNodes should be similar
 - Usually run when new DataNodes are added
 - Cluster is online when Rebalancer is active
 - Rebalancer is throttled to avoid network congestion
 - Command line tool



Secondary NameNode

- Copies FsImage and Transaction Log from Namenode to a temporary directory
- Merges FSImage and Transaction Log into a new FSImage in temporary directory
- Uploads new FSImage to the NameNode

Transaction Log on NameNode is purged



User Interface

- Commands for HDFS User:
 - hadoop dfs -mkdir /foodir
 - hadoop dfs -cat /foodir/myfile.txt
 - hadoop dfs -rm /foodir/myfile.txt
- Commands for HDFS Administrator
 - hadoop dfsadmin -report
 - hadoop dfsadmin -decommision datanodename
- Web Interface
 - http://host:port/dfshealth.jsp



MapReduce

Original Slides by Owen O'Malley (Yahoo!)

&

Christophe Bisciglia, Aaron Kimball & Sierra Michells-Slettvet

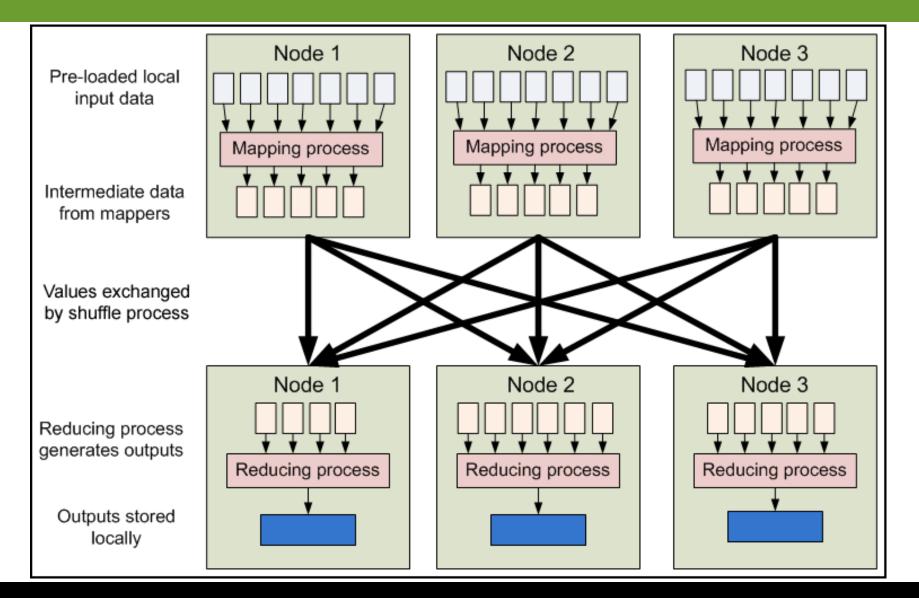


MapReduce - What?

- MapReduce is a programming model for efficient distributed computing
- It works like a Unix pipeline
 - cat input | grep | sort | uniq -c | cat > output
 - Input | Map | Shuffle & Sort | Reduce | Output
- Efficiency from
 - Streaming through data, reducing seeks
 - Pipelining
- A good fit for a lot of applications
 - Log processing
 - Web index building



MapReduce - Dataflow





MapReduce - Features

- Fine grained Map and Reduce tasks
 - Improved load balancing
 - Faster recovery from failed tasks
- Automatic re-execution on failure
 - In a large cluster, some nodes are always slow or flaky
 - Framework re-executes failed tasks
- Locality optimizations
 - With large data, bandwidth to data is a problem
 - Map-Reduce + HDFS is a very effective solution
 - Map-Reduce queries HDFS for locations of input data
 - Map tasks are scheduled close to the inputs when possible

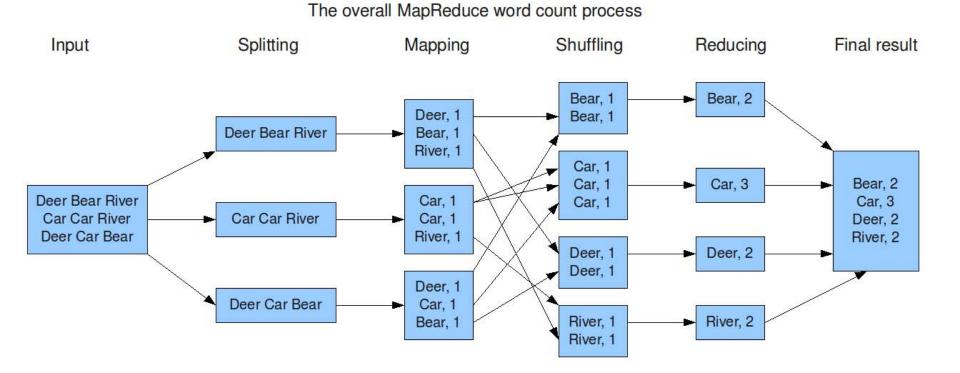


Word Count Example

- Mapper
 - Input: value: lines of text of input
 - Output: key: word, value: 1
- Reducer
 - Input: key: word, value: set of counts
 - Output: key: word, value: sum
- Launching program
 - Defines this job
 - Submits job to cluster



Word Count Dataflow





Word Count Mapper

```
public static class Map extends MapReduceBase implements
Mapper<LongWritable,Text,Text,IntWritable> {
    private static final IntWritable one = new IntWritable(1);
    private Text word = new Text();
```

```
public static void map(LongWritable key, Text value,
    OutputCollector<Text,IntWritable> output, Reporter reporter) throws
    IOException {
    String line = value.toString();
    StringTokenizer = new StringTokenizer(line);
    while(tokenizer.hasNext()) {
        word.set(tokenizer.nextToken());
        output.collect(word,one);
        }
    }
}
```



Word Count Reducer

```
public static class Reduce extends MapReduceBase implements
   Reducer<Text,IntWritable,Text,IntWritable> {
public static void reduce(Text key, Iterator<IntWritable> values,
   OutputCollector<Text,IntWritable> output, Reporter reporter) throws
   IOException {
     int sum = 0;
     while(values.hasNext()) {
       sum += values.next().get();
     }
     output.collect(key, new IntWritable(sum));
```



Word Count Example

- Jobs are controlled by configuring *JobConfs*
- JobConfs are maps from attribute names to string values
- The framework defines attributes to control how the job is executed
 - conf.set("mapred.job.name", "MyApp");
- Applications can add arbitrary values to the JobConf
 - conf.set("my.string", "foo");
 - conf.set("my.integer", 12);
- JobConf is available to all tasks



Putting it all together

- Create a launching program for your application
- The launching program configures:
 - The *Mapper* and *Reducer* to use
 - The output key and value types (input types are inferred from the *InputFormat*)
 - The locations for your input and output
- The launching program then submits the job and typically waits for it to complete



Putting it all together

JobConf conf = new JobConf(WordCount.class); conf.setJobName("wordcount");

conf.setOutputKeyClass(Text.class); conf.setOutputValueClass(IntWritable.class);

conf.setMapperClass(Map.class); conf.setCombinerClass(Reduce.class); conf.setReducer(Reduce.class);

conf.setInputFormat(TextInputFormat.class); Conf.setOutputFormat(TextOutputFormat.class);

FileInputFormat.setInputPaths(conf, new Path(args[0])); FileOutputFormat.setOutputPath(conf, new Path(args[1]));

JobClient.runJob(conf);



Input and Output Formats

- A Map/Reduce may specify how it's input is to be read by specifying an *InputFormat* to be used
- A Map/Reduce may specify how it's output is to be written by specifying an *OutputFormat* to be used
- These default to *TextInputFormat* and *TextOutputFormat*, which process line-based text data
- Another common choice is SequenceFileInputFormat and SequenceFileOutputFormat for binary data
- These are file-based, but they are not required to be



How many Maps and Reduces

- Maps
 - Usually as many as the number of HDFS blocks being processed, this is the default
 - Else the number of maps can be specified as a hint
 - The number of maps can also be controlled by specifying the minimum split size
 - The actual sizes of the map inputs are computed by:
 - max(min(block_size,data/#maps), min_split_size)
- Reduces
 - Unless the amount of data being processed is small
 - 0.95*num_nodes*mapred.tasktracker.tasks.maximum



Some handy tools

- Partitioners
- Combiners
- Compression
- Counters
- Zero Reduces
- Distributed File Cache
- Tool



Partitioners

- Partitioners are application code that define how keys are assigned to reduces
- Default partitioning spreads keys evenly, but randomly

 Uses key.hashCode() % num_reduces
- Custom partitioning is often required, for example, to produce a total order in the output
 - Should implement *Partitioner* interface
 - Set by calling conf.setPartitionerClass(MyPart.class)
 - To get a total order, sample the map output keys and pick values to divide the keys into roughly equal buckets and use that in your partitioner



Combiners

- When *maps* produce many repeated keys
 - It is often useful to do a local aggregation following the map
 - Done by specifying a *Combiner*
 - Goal is to decrease size of the transient data
 - Combiners have the same interface as Reduces, and often are the same class
 - Combiners must **not** side effects, because they run an intermdiate number of times
 - In WordCount, conf.setCombinerClass(Reduce.class);



Compression

- Compressing the outputs and intermediate data will often yield huge performance gains
 - Can be specified via a configuration file or set programmatically
 - Set mapred.output.compress to true to compress job output
 - Set *mapred.compress.map.output* to *true* to compress map outputs
- Compression Types (*mapred(.map)?.output.compression.type*)
 - "block" Group of keys and values are compressed together
 - "record" Each value is compressed individually
 - Block compression is almost always best
- Compression Codecs (mapred(.map)?.output.compression.codec)
 - Default (zlib) slower, but more compression
 - LZO faster, but less compression



Counters

- Often Map/Reduce applications have countable events
- For example, framework counts records in to and out of Mapper and Reducer
- To define user counters: static enum Counter {EVENT1, EVENT2}; reporter.incrCounter(Counter.EVENT1, 1);
- Define nice names in a MyClass_Counter.properties file

```
CounterGroupName=MyCounters
```

```
EVENT1.name=Event 1
```

```
EVENT2.name=Event 2
```



Zero Reduces

- Frequently, we only need to run a filter on the input data
 - No sorting or shuffling required by the job
 - Set the number of reduces to 0
 - Output from maps will go directly to OutputFormat and disk



Distributed File Cache

- Sometimes need read-only copies of data on the local computer
 - Downloading 1GB of data for each Mapper is expensive
- Define list of files you need to download in JobConf
- Files are downloaded once per computer
- Add to launching program: DistributedCache.addCacheFile(new URI("hdfs://nn:8020/foo"), conf);
- Add to task:
 - Path[] files =

DistributedCache.getLocalCacheFiles(conf);



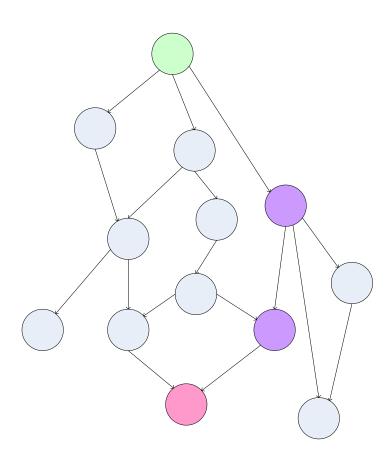
Tool

- Handle "standard" Hadoop command line options
 - -conf file load a configuration file named file
 - D prop=value define a single configuration property prop
- Class looks like:



Example: Finding the Shortest Path

- A common graph search application is finding the shortest path from a start node to one or more target nodes
- Commonly done on a single machine with *Dijkstra's Algorithm*
- Can we use BFS to find the shortest path via MapReduce?





Finding the Shortest Path: Intuition

- We can define the solution to this problem inductively
 - DistanceTo(startNode) = 0
 - For all nodes *n* directly reachable from startNode, DistanceTo(n) = 1
 - For all nodes *n* reachable from some other set of nodes S,

DistanceTo(n) = 1 + min(DistanceTo(m), $m \in S$)



From Intuition to Algorithm

- A map task receives a node n as a key, and (D, points-to) as its value
 - *D* is the distance to the node from the start
 - *points-to* is a list of nodes reachable from *n*
- $\Box \forall p \in points-to, emit (p, D+1)$
- Reduces task gathers possible distances to a given p and selects the minimum one



What This Gives Us

- This MapReduce task can advance the known frontier by one hop
- To perform the whole BFS, a non-MapReduce component then feeds the output of this step back into the MapReduce task for another iteration
 - Problem: Where'd the *points-to* list go?
 - Solution: Mapper emits (*n*, *points-to*) as well



Blow-up and Termination

- This algorithm starts from one node
- Subsequent iterations include many more nodes of the graph as the frontier advances
- Does this ever terminate?
 - Yes! Eventually, routes between nodes will stop being discovered and no better distances will be found. When distance is the same, we stop
 - Mapper should emit (*n*,*D*) to ensure that "current distance" is carried into the reducer



Extensions to Map-Reduce

- There are other systems that
 - Use Distributed file systems
 - Manage large number of tasks that are instantiations of user-written functions.
 - Deal with failures gracefully.
 - Example from Ullman's book:

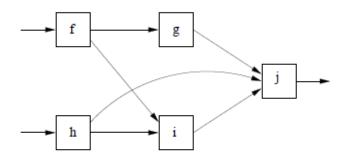


Figure 2.6: An example of a workflow that is more complex than Map feeding Reduce



Extensions to Map-Reduce

- Main applications for Workflow systems
 - Representing a cascade of multiple-map reduce jobs
 - Complex distributed tasks
- Generally more efficient than running multiple mapreduce sequentially.
 - Writing results to hard disks could be problematic.
 - Potential pipelining optimizations.



Hadoop Subprojects



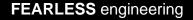
Hadoop Related Subprojects

- Pig
 - High-level language for data analysis
- Hbase
 - Table storage for semi-structured data
- Zookeeper
 - Coordinating distributed applications
- Hive
 - SQL-like Query language and Metastore
- Mahout
 - Machine learning



Pig

Original Slides by Matei Zaharia UC Berkeley RAD Lab





Pig

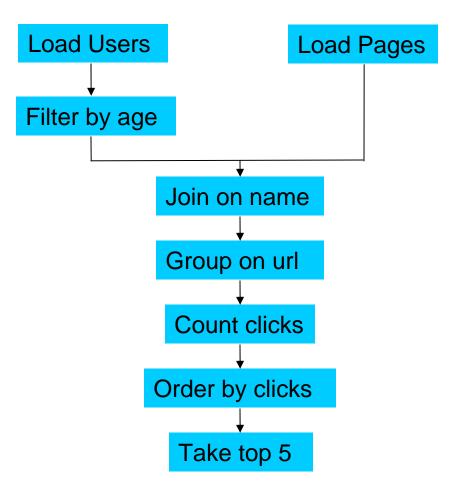
- Started at Yahoo! Research
- Now runs about 30% of Yahoo!'s jobs
- Features
 - Expresses sequences of MapReduce jobs
 - Data model: nested "bags" of items
 - Provides relational (SQL) operators
 (JOIN, GROUP BY, etc.)
 - Easy to plug in Java functions





An Example Problem

 Suppose you have user data in a file, website data in another, and you need to find the top 5 most visited pages by users aged 18-25





In MapReduce

import java.io.IOException; import java.util.ArrayList; import java.util.Iterator; import java.util.List; support org.papeche.hadoop.fs.pihj import org.apache.hadoop.io.jong/ttable; import org.apache.hadoop.io.fs.pitable import org.apache.hadoop.io.fs.pitable import org.apache.hadoop.io.fs.pitable import org.apache.hadoop.mapred.lieoutputFormat; import org.apache.hadoop.mapred.judrextInputFormat; import org.apache.hadoop.mapred.judrextInputFormat; import org.apache.hadoop.mapred.lieoutputFormat; import org.apache.hadoop.mapred.lieoutputFormat; import org.apache.hadoop.mapred.lieoutputFormat; import org.apache.hadoop.mapred.lieoutputFormat; import org.apache.hadoop.mapred.lieoutputFormat; import org.apache.hadoop.mapred.kecordReader; import org.apache.hadoop.mapred.kecordReader; import org.apache.hadoop.mapred.securefileOutputFormat; import org.apa import org.apache.hadoop.fs.Path: public class MRExample { public static class LoadPages extends MapReduceBase implements Mapper<LongWritable, Text, Text, Text> { // Pull the key out
String line = val.toString();
int firstComma = line.indexOf(','); int firstComma = line.indexOf(','); String key = line.substring(0, firstComma); String value = line.substring(firstComma + 1); (// Frepend an index to the value so we know which file // ic came from. Text outVal = new Text('); * value); oc.collect(outKey, outVal); } public static class LoadAndFilterUsers extends MapReduceBase implements Mapper<LongWritable, Text, Text, Text> { String line = val.toString(); int firstComma = line.indexof(','); String value = line.substring(firstComma + 1); Int age = Integr.parsen(tvalue); String key = line.substring(0, firstComma); Text outKey = new Text(key); // Frepend an index to the value so we know which file // it came from. Text outVal = new Text(key); Text outVal = new Text(subs); oc:collect(outKey, outVal); public static class Join extends MapReduceBase implements Reducer<Text, Text, Text, Text> { public void reduce(Text key, Iterator<Text> iter, OutputCollector<Text, Text> oc, Reporter reporter) throws IOException { // For each value, figure out which file it's from and store it // accordingly. List<String> first = new ArrayList<String>(); List<String> second = new ArrayList<String>(); while (iter.hasNext()) { Text t = iter.next(); String value = t.toString(); if (value.charAt(0) == '1')

// Do the cross product and collect the values for (String sl : first) (reporter.setStatus("OK"); >) } public static class LoadJoined extends MapReduceBase implements Mapper<Text, Text, Text, LongWritable> { public void map(Text k, Text val, OutputCollector<Text, LongWritable> oc, Reporter reporter) throws IOException { // Find the url String line = val.toString(); pint of the second 3 public static class ReduceUrls extends MapReduceBase implements Reducer<Text, LongWritable, WritableComparable, Writable> { public void reduce(ilc void reduce(
 Text key, n
 Textkey, n
 terator<LongWritable> iter,
 OutputCollector<WritableComparable, Writable> oc,
 Reporter reporter) throws IOException {
 // Add up all the values we see
 long sum = 0; while (iter.hasNext()) { sum += iter.next().get(); reporter.setStatus("OK"); oc.collect(key, new LongWritable(sum)); } public static class LoadClicks extends MapReduceBase implements Mapper<WritableComparable, Writable, LongWritable, Text> { public void map(WritableComparable key, writablecomparable key, Writable val, OutputCollector<LongWritable, Text> oc, Reporter reporter) throws IOException (oc.collect((LongWritable)val, (Text)key); }
public static class LimitClicks extends MapReduceBase
implements Reducer<LongWritable, Text, LongWritable, Text> { int count = 0; int count = 0; public void reduce(LongWritable key, Iterator<Text> iter, OutputCollector<LongWritable, Text> oc, Reporter reporter) throws IOException { // Only output the first 100 records
while (count < 100 && iter.hasNext()) {
 co.collect(key, iter.next());
 count+;;</pre> } }
public static void main(String[] args) throws IOException {
 JobConf 1p = new JobConf(NRExample.class);
 lp.setJobName("Load Pages");
 lp.setInputFormat(TextInputFormat.class);

reporter.setStatus("OK");

lp.setOutputKeyClass(Text.class); lp.setOutputValueClass(Text.class); lp.setMapperClass(LoadPages.class); FileInputFormat.addInputPath(lp, new Path("/user/gates/pages")); FileOutputFormat.setOutputPath(1p, new Path(*/user/gates/tmp/indexed_pages")); lp.setNumReduceTasks(0); Job loadPages = new Job(lp);

JobConf lfu = new JobConf(MRExample.class); lfu.setJobRame("Load and Filter Users"); lfu.setInputFormat(TextInputFormat.class); lfu.setOutputFeytClass(Text.class); lfu.setUnputFayLClass(Text.class); lfu.setHApperClass(LoadAndFilterUsers.class); lieInputFormat.addInputFah(lfu, new Path("/user/gates/users")); FileOutputFormat.setOutputPath(lfu, new Path("/user/gates/tmp/filtered_users"));
lfu.setNumReduceTasks(0); Job loadUsers = new Job(lfu); JobConf join - new JobConf (MERxample.class); Join.setIputFormat(KeyVals.class); join.setIputFormat(KeyVals.class); join.setUtputFormat(KeyVals.class); join.setUtputFormat(KeyVals.class); join.setKeduerclass(Join.class); FileInputFormat.addInputFot(join, new Fakt('Jusr/gates/tep/Indexad_pages')); put

FileInputFormat.addInputPath(join, new Path("/user/gates/tmp/filtered_users")); FileOutputFormat.setOutputPath(join, new Path("/user/gates/tmp/joined")); join.setNumReduceTasks(50); Job joinJob = new Job(join); joinJob.addDependingJob(loadPages); joinJob.addDependingJob(loadUsers);

JobConf group = new JobConf(MRExample.class); group.setJobName("Group URLs"); JobConf group = new JobConf (MEXample.class); group.setDupUtFormat (KeyNuleFacthputFormat.class); group.setDupUtFormat(KeyNuleFacthputFormat.class); group.setDupUtFormat(Legal (LongWritable.class); group.setDupUtFoluc(Lass(LongWritable.class); group.setDubutFoluc(Lass(MeducUtIs.class); group.setEdbuerClass(MeducUtIs.class); group.setEdbuerClass(MeducUtIs.class); group.setEdbuerClass(MeducUtIs.class); group.setEdbuerClass(MeducUtIs.class); Path("/user/gates/tmp/joined")); FileOutputFormat.setOutputPath(group, new Path("/user/gates/tmp/grouped")); group.setNumReduceTasks(50); Job groupJob = new Job(group); groupJob.addDependingJob(joinJob);

JobConf top100 = new JobConf(MRExample.class): JOBCONT top100 = new JOBCONT (MEXAMPLe.class); top100.setInputFormat(SequenceFileInputFormat.class); top100.setOutputFeyClass(LongWritable.class); top100.setOutputFeyClass(Text.class); top100.setOutputValueClass(Text.class); top100.setOutputFormat.CsequenceFileOutputFormat.class); top100.setMapperClass(LoadClicks.class); top100.setReducerClass(LimitClicks.class); top100.setReducerClass(LimitClicks.class); FileInputFormat.addInputPath(top100, new Path("Juser/gates/imp/grouped")); Path("/user/gates/imp/grouped")); FileOutputFormat.setOutputPath(top100, new Path("/user/gates/top100sitesforusers18to25")); top100.setNumReduceTasks(1); Job limit = new Job(ton100) limit.addDependingJob(groupJob);

JobControl jc = new JobControl("Find top 100 sites for users JobLoncies 18 to 25"); jc.addJob(loadPages); jc.addJob(loadJsers); jc.addJob(joinJob); - add.Tob(groupJob); jc.addJob(limit); jc.run(); }

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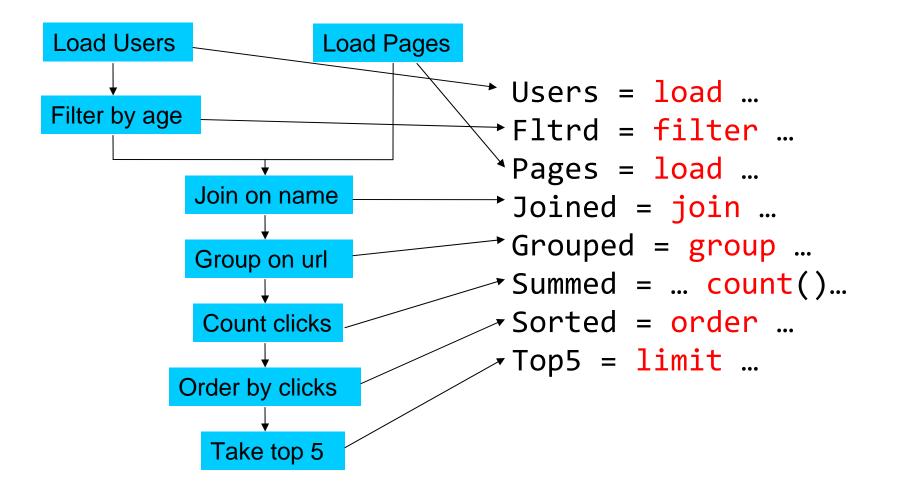


In Pig Latin

Users = load 'users' as (name, age); Filtered = filter Users by age >= 18 and age <= 25; Pages = load 'pages' as (user, url); Joined = join Filtered by name, Pages by user; Grouped = group Joined by url; Summed = foreach Grouped generate group, count(Joined) as clicks; Sorted = order Summed by clicks desc; Top5 = limit Sorted 5; store Top5 into 'top5sites';

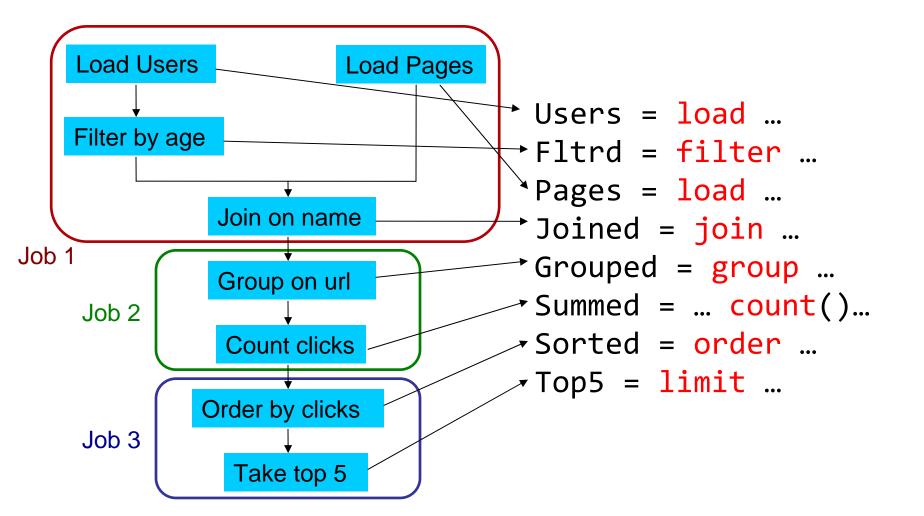


Ease of Translation





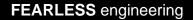
Ease of Translation





HBase

Original Slides by Tom White Lexeme Ltd.





HBase - What?

- Modeled on Google's Bigtable
- Row/column store
- Billions of rows/millions on columns
- Column-oriented nulls are free
- Untyped stores byte[]



Row	Timestamp	Column family: animal:		Column family repairs:
		animal:type	animal:size	repairs:cost
enclosure1	t2	zebra		1000 EUR
	t1	lion	big	
enclosure2				



Column family animal:

(enclosure1, t2, animal:type)	zebra
(enclosure1, t1, animal:size)	big
(enclosure1, t1, animal:type)	lion

Column family repairs:

(enclosure1, t1, repairs:cost)	1000 EUR
--------------------------------	----------



HBase - Code

Htable table = ... Text row = new Text("enclosure1"); Text col1 = new Text("animal:type"); Text col2 = new Text("animal:size"); BatchUpdate update = new BatchUpdate(row); update.put(col1, "lion".getBytes("UTF-8")); update.put(col2, "big".getBytes("UTF-8)); table.commit(update);

update = new BatchUpdate(row); update.put(col1, "zebra".getBytes("UTF-8")); table.commit(update);



HBase - Querying

• Retrieve a cell

Cell = table.getRow("enclosure1").getColumn("animal:type").getValue();

• Retrieve a row

RowResult = table.getRow("enclosure1");

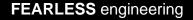
• Scan through a range of rows

Scanner s = table.getScanner(new String[] { "animal:type" });



Hive

Original Slides by Matei Zaharia UC Berkeley RAD Lab

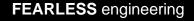




Hive

- Developed at Facebook
- Used for majority of Facebook jobs
- "Relational database" built on Hadoop
 - Maintains list of table schemas
 - SQL-like query language (HiveQL)
 - Can call Hadoop Streaming scripts from HiveQL
 - Supports table partitioning, clustering, complex data types, some optimizations







Creating a Hive Table

CREATE TABLE page_views(viewTime INT, userid BIGINT, page_url STRING, referrer_url STRING, ip STRING COMMENT 'User IP address') COMMENT 'This is the page view table' PARTITIONED BY(dt STRING, country STRING) STORED AS SEQUENCEFILE;

 Partitioning breaks table into separate files for each (dt, country) pair
 Ex: /hive/page_view/dt=2008-06-08,country=USA /hive/page_view/dt=2008-06-08,country=CA



A Simple Query

 Find all page views coming from xyz.com on March 31st:

```
SELECT page_views.*
FROM page_views
WHERE page_views.date >= '2008-03-01'
AND page_views.date <= '2008-03-31'
AND page_views.referrer_url like '%xyz.com';</pre>
```

 Hive only reads partition 2008-03-01,* instead of scanning entire table



Aggregation and Joins

• Count users who visited each page by gender:

SELECT pv.page_url, u.gender, COUNT(DISTINCT u.id)
FROM page_views pv JOIN user u ON (pv.userid = u.id)
GROUP BY pv.page_url, u.gender
WHERE pv.date = '2008-03-03';

• Sample output:

page_url	gender	count(userid)
home.php	MALE	12,141,412
home.php	FEMALE	15,431,579
photo.php	MALE	23,941,451
photo.php	FEMALE	21,231,314

Using a Hadoop Streaming Mapper Script

```
FROM page_views;
```

