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# Rabin Crypto System Overview

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## The Rabin Cryptosystem

- Computationally secure against a chosen plaintext attack
  - Provided that the modulus  $n = pq$  can not be factored.

$$p \equiv q \equiv 3 \pmod{4}$$

- $n$  is the public key. The primes  $p$  and  $q$  are the private key.
- Choose  $a$  to simplify the computation of square roots modulo  $p$  and  $q$



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- B encrypts a message  $m$  and sends the ciphertext  $c$  to A
- Encryption:
  - Obtain A's public key  $n$ .
  - Represent the message as an integer  $m$  in the range  $\{0, 1, \dots, n-1\}$ .
  - Compute  $c = m^2 \bmod n$
  - Send the ciphertext  $c$  to A

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- A decrypts the ciphertext  $c$  as follows:
- Decryption:
  - Compute  $\sqrt{c} \bmod n$
  - There are four square roots  $m_1, m_2, m_3, m_4$  of  $c$  modulo  $n$ .
  - The message  $m$  is equal to one of these four messages

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- When  $p \equiv 3 \pmod{4}$  there is a simple formula to compute the square root of  $c$  in mod  $p$ .

$$\begin{aligned} (\pm c^{(p+1)/4})^2 &\equiv c^{(p+1)/2} \pmod{p} \\ &\equiv c^{(p-1)/2} c \pmod{p} \\ &\equiv c \pmod{p} \end{aligned}$$

- Here we have made use of Euler's criterion to claim that  $c^{(p-1)/2} \equiv 1 \pmod{p}$

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- Hence the two square roots of  $c \pmod{p}$  are

$$\pm c^{(p+1)/4} \pmod{p}$$

- In a similar fashion, the two square roots of  $c \pmod{q}$  are  $\pm c^{(q+1)/4} \pmod{q}$

- Then we can obtain the four square roots of  $c \pmod{n}$  using the Chinese Remainder Theorem

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- Example:  $n = 77 = 7 \times 11$ 
  - Suppose  $c = m^2 \pmod{77}$
  - Then for message  $m$  the ciphertext  $c$  is computed as  $\sqrt{c} \pmod{77}$
  - And for decryption we need to compute
$$c \equiv 10^2 \equiv 23 \pmod{77}$$
  - Suppose Alice wants to send message  $m = 10$

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- To find the square roots of 23 in mod 7 and in mod 11 we can use the formula since 7 and 11 are congruent to 3 mod 4.

$$23^{(7+1)/4} \equiv 2^2 \equiv 4 \pmod{7}$$

$$23^{(11+1)/4} \equiv 1^3 \equiv 1 \pmod{11}$$

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- Using the Chinese Remainder Theorem, we compute the four square roots of 23 mod 77 to be

$$\pm 10, \pm 32 \pmod{77}$$

- Therefore the four possible plaintexts are

$$m_1 = 10, m_2 = 67, m_3 = 32, m_4 = 45$$