

# **MORE**

## **Learning Theory**

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# Last Time

- **Probably approximately correct (PAC)**
  - The only reasonable expectation of a learner is that with high probability it learns a close approximation to the target concept
  - Specify two small parameters,  $0 < \epsilon, 0 < \delta < 1$ 
    - $\epsilon$  is the error of the approximation
    - $(1 - \delta)$  is the probability that, given  $m$  i.i.d. samples, our learning algorithm produces a classifier with error at most  $\epsilon$

# Learning Theory

- We use the observed data in order to learn a classifier
- Want to know how far the learned classifier deviates from the (unknown) underlying distribution
  - With too few samples, we will with high probability learn a classifier whose true error is quite high even though it may be a perfect classifier for the observed data
  - As we see more samples, we pick a classifier from the hypothesis space with low training error & hope that it also has low true error
    - Want this to be true with high probability – can we bound how many samples that we need?

# Haussler, 1988

- What we proved last time:

**Theorem:** For a finite hypothesis space,  $H$ , with  $m$  i.i.d. samples, and  $0 < \epsilon < 1$ , the probability that any consistent classifier has true error larger than  $\epsilon$  is at most  $|H|e^{-\epsilon m}$

- We can turn this into a sample complexity bound

# Sample Complexity

- Let  $\delta$  be an upper bound on the desired probability of not  $\epsilon$ -exhausting the sample space
  - The probability that the version space is not  $\epsilon$ -exhausted is at most  $|H|e^{-\epsilon m} \leq \delta$
  - Solving for  $m$  yields

$$\begin{aligned} m &\geq -\frac{1}{\epsilon} \ln \frac{\delta}{|H|} \\ &= \left( \ln |H| + \ln \frac{1}{\delta} \right) / \epsilon \end{aligned}$$

# PAC Bounds

**Theorem:** For a finite hypothesis space  $H$ ,  $m$  i.i.d. samples, and  $0 < \epsilon < 1$ , the probability that true error of any of the best classifiers (i.e., lowest training error) is larger than its training error plus  $\epsilon$  is at most  $|H|e^{-2m\epsilon^2}$

- Sample complexity (for desired  $\delta \geq |H|e^{-2m\epsilon^2}$ )

$$m \geq \left( \ln|H| + \ln \frac{1}{\delta} \right) / 2\epsilon^2$$

# PAC Bounds

- If we require that the previous error is bounded above by  $\delta$ , then with probability  $(1 - \delta)$ , for all  $h \in H$

$$\epsilon_h \leq \underbrace{\epsilon_h^{train}}_{\text{“bias”}} + \underbrace{\sqrt{\frac{1}{2m} \left( \ln |H| + \ln \frac{1}{\delta} \right)}}_{\text{“variance”}}$$

– For small  $|H|$

- High bias (may not be enough hypotheses to choose from)
- Low variance

# PAC Bounds

- If we require that the previous error is bounded above by  $\delta$ , then with probability  $(1 - \delta)$ , for all  $h \in H$

$$\epsilon_h \leq \underbrace{\epsilon_h^{train}}_{\text{"bias"}} + \underbrace{\sqrt{\frac{1}{2m} \left( \ln |H| + \ln \frac{1}{\delta} \right)}}_{\text{"variance"}}$$

– For large  $|H|$

- Low bias (lots of good hypotheses)
- High variance



# VC Dimension

- Our analysis for the finite case was based on  $|H|$ 
  - If  $H$  isn't finite, this translates into infinite sample complexity
  - We can derive a different notion of complexity for infinite hypothesis spaces by considering only the number of points that can be correctly classified by some member of  $H$
  - We will only consider the binary case for now

# VC Dimension

- How many points in 1-D can be correctly classified by a linear separator?
  - 2 points:



Yes!

# VC Dimension

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Yes!

# VC Dimension

- How many points in 1-D can be correctly classified by a linear separator?
  - 2 points:



# VC Dimension

- How many points in 1-D can be correctly classified by a linear separator?
  - 3 points:



Yes!

# VC Dimension

- How many points in 1-D can be correctly classified by a linear separator?
  - 3 points:



NO!

# VC Dimension

- How many points in 1-D can be correctly classified by a linear separator?

– 3 points:



- 3 points and up: for any collection of three or more there is always some choice of pluses and minuses such that that the points cannot be classified with a linear separator

# VC Dimension

- A set of points is **shattered** by a hypothesis space  $H$  if and only if for every partition of the set of points into positive and negative examples, there exists some consistent  $h \in H$
- The **Vapnik–Chervonenkis (VC) dimension** of  $H$  over inputs from  $X$  is the size of the *largest* finite subset of  $X$  shattered by  $H$



# VC Dimension

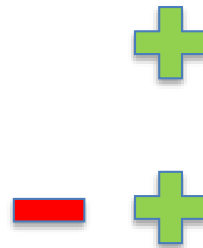
- Common misconception:
  - VC dimension is determined by the largest shattered set of points, not the highest number such that all sets of points that size can be shattered



Cannot be shattered by a line

# VC Dimension

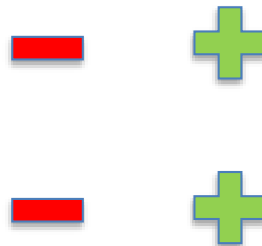
- Common misconception:
  - VC dimension is determined by the largest shattered set of points, not the highest number such that all sets of points that size can be shattered



Can be shattered by a line (no matter the labels), so VC dimension is at least 3

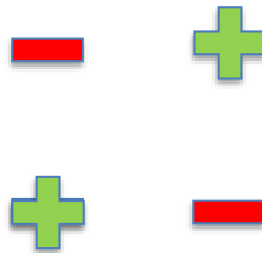
# VC Dimension

- What is the VC dimension of 2-D space under linear separators?
  - It is at least three from the last slide
  - Can some set of four points be shattered?



# VC Dimension

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# VC Dimension

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NO! This means that the VC dimension is at most 3

# VC Dimension

- There exists a linear separator that can shatter any set of size  $d + 1$  in a  $d - \textit{dimensional}$  space, but not  $d + 2$
- The larger the subset of  $X$  that can be shattered, the more expressive the hypothesis space is
- If arbitrarily large finite subsets of  $X$  can be shattered, then  $VC(H) = \infty$

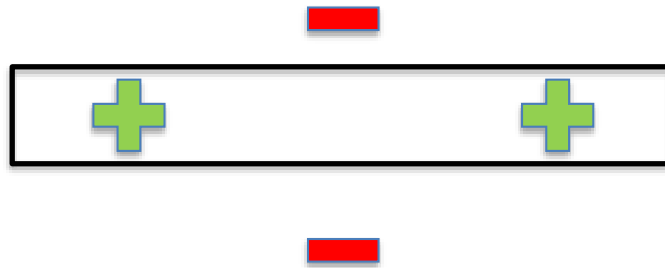
# Axis Parallel Rectangles

- Let  $X$  be the set of all points in  $\mathbb{R}^2$
- Let  $H$  be the set of all axis parallel rectangles in 2-D
  - What is  $VC(H)$ ?

# Axis Parallel Rectangles

- Let  $X$  be the set of all points in  $\mathbb{R}^2$
- Let  $H$  be the set of all axis parallel rectangles in 2-D

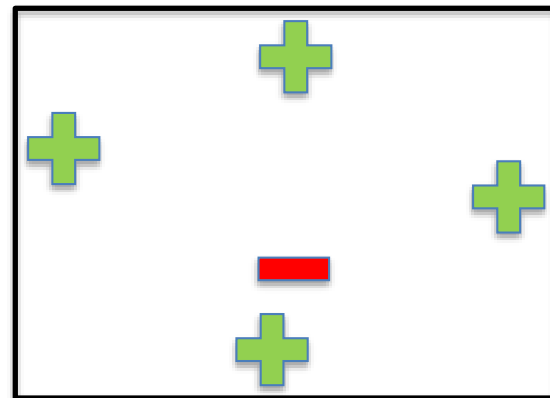
$$-VC(H) \geq 4$$





# Axis Parallel Rectangles

- Let  $X$  be the set of all points in  $\mathbb{R}^2$
- Let  $H$  be the set of all axis parallel rectangles in 2-D
  - $VC(H) = 4$
  - A rectangle can contain at most 4 extreme points, the fifth point must be contained within the rectangle defined by these points



# Examples

- VC dimension of decision trees?
- VC dimension of 1-NN?
- VC dimension of linear separators through the origin?

# PAC Bounds with VC Dimension

- VC dimension can be used to construct PAC bounds

$$m \geq \frac{1}{\epsilon} \left( 4 \ln \frac{2}{\delta} + 8 \cdot VC(H) \ln \frac{13}{\epsilon} \right)$$

- With probability at least  $(1 - \delta)$  every  $h \in H$  satisfies

$$\epsilon_h \leq \epsilon_h^{train} + \sqrt{\frac{1}{m} \left( VC(H) \left( \ln \left( \frac{2m}{VC(H)} \right) + 1 \right) + \ln \frac{4}{\delta} \right)}$$

- These bounds (and the preceding discussion) only work for binary classification, but there are generalizations

# PAC Learning

- **Given:**
  - Set of data  $X$
  - Hypothesis space  $H$
  - Set of target concepts  $C$
  - Training instances from unknown probability distribution over  $X$  of the form  $(x, c(x))$
- **Goal:**
  - Learn the target concept  $c \in C$

# PAC Learning

- Given:
  - A concept class  $C$  over  $n$  instances from the set  $X$
  - A learner  $L$  with hypothesis space  $H$
  - Two constants,  $\epsilon, \delta \in (0, \frac{1}{2})$
- $C$  is said to be PAC learnable by  $L$  using  $H$  iff for all distributions over  $X$ , learner  $L$  by sampling  $n$  instances, will with probability at least  $1 - \delta$  outputs a hypothesis  $h \in H$  such that
  - $\epsilon_h \leq \epsilon$
  - Running time is polynomial in  $\frac{1}{\epsilon}, \frac{1}{\delta}, n, size(c)$