# Phylogenetics 

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(some material based on content by PR in Eric Xing's 10-810 Carnegie Mellon class)


Ernst Haeckel, Tree of life, 1879

## Some nomenclature on trees

- Binary (bifurcating) or multiway (multifurcating)

- A multiway split can be envisioned as a series of binary splits



## Recursively defining trees: subtrees

Trees can be defined in terms of themselves


Jut dallas

Induced subtree


- For a subset of nodes, the induced subtree is the minimal set of edges and nodes to connect them together, taken from the original tree


## Parameterizing a tree

- Topology : connectivity of the tree
- path between any two nodes is unique!
- topology unchanged by "squishing" a drawing of the tree
- Length of edges : the geometry of the tree - notion of "distance" between neighboring nodes
- Labels of nodes (?) : which nodes were not born equal : identifying the ones we are interested in


## Rooted trees

- "Rooting" : a notion of direction of flow - in case of phylogenies, flow of time
- Root : a privileged node
- Direction flows outward from the root
- Rooted tree = Directed edges away from root


Figure 1. Topologies of phylogenetic tress: A. unreoted tree, B. roosed tree.
Perrero \& Lopez, Gen Mol Res, Sep 2005

Directed edges

- Directed edges away from root = Rooting


THIS IS A TREE. (PIC KIT UP

## Is this graph a rooted tree?

- Do a topological sorting on it, is there a unique root?
- Can we order all the nodes in the direction of the edges, and be left with a single topmost node



## Rooted tree nomenclature



Ordered vs unordered branches

- Is the order in which we represent the siblings important ? $2^{\wedge} n$ ways to draw for $n$ interior nodes (rooted)

- Which is different ?




Proper vs improper trees

- Proper : each node has 0 or 2 children (rooted), each node has 1 or 3 neighbors (unrooted)

ROOTED


IMPROPER

HNROOTED



## Phylogeny / phylogenetic tree

- Taxon/taxa/operational taxonomic unit (OTU)
- unit of classification : species, subspecies, individual, etc
insect.eugenes.org
- Phylogeny = evolutionary tree
- Hypothesis concerning evolutionary history of taxa



## Molecular Phylogeny

- Molecular phylogeny = based on models ( or distances based on models) of molecular evolution

insect.eugenes.org

Phylogenetic Trees

- Edges = branches
- Interior nodes: ancestral taxa
- Leaf / exterior nodes : contemporary taxa
- Topology : relation between species
- Branch length = edge weights: amt of change
- Labelled leaves, UT Ranlabelled ancestors


No BRANCH LENGTH


ROOT, No CLOCK

ROOT, WITH CLOCK

Phylogenies have unordered



Ordering has no evolutionary connotation

## Phylogenetic trees are proper trees



Improper phylogenies transformed to proper ones

- Each node = Common ancestor of subset of species
- Unrooting


Most recent common ancestor


## Rooting an unrooted tree

- Root at any branch
- Sometimes, we may not know (or may not care) where the root is



## Counting edges

- For an unrooted tree on n leaves, we have
- n-2 internal nodes
- ( proof by induction, tree on 3 leaves has 1 internal node, every additional leaf incorporated into a tree adds one leaf and one internal node )
$-2 n-3$ edges [ no of nodes minus one ]
- ( think of shrinking the tree one node and edge at a time)

embl.de


## Counting labelled leaf, unlabelled ancestor phylogenetic trees

- Topologically equivalent = a tree changed to another by flipping neighbors, w/o breaking branches
- For a tree of $n-1$ leaves, $2 n-5$ branches : add nth leaf to a branch : $T(n)=T(n-1) X(2 n-5)$
$-T(3)=1$


Z Yang

- If rooted, rooting can happen at any branch : additional factor of $(2 n-5)$


# No of unrooted trees: $\mathrm{T}(\mathrm{k})$ 

$$
T(k)=\frac{(2 k-5)!}{2^{k-3}(k-3)!}
$$

## Grows fast

- Felsenstein, Counting trees

Table 1. The numbers of rooted trees with $n$ labelled tips and with unlabelled interior NODES. THE LEFT COLUMN COUNTS ALL TREES, THE RIGHT COLUMN ONLY BIFURCATING TREES

| $n$ | All trees | Bifurcating trees |
| :---: | :---: | :---: |
| 1 | 1 | 1 |
| 2 | 1 | 1 |
| 3 | 4 | 3 |
| 4 | 26 | 15 |
| 5 | 236 | 105 |
| 6 | 2,752 | 945 |
| 7 | 39,208 | 10,395 |
| 8 | 660,032 | 135,135 |
| 9 | 12,818,912 | 2,027,025 |
| 10 | 282,137,824 | 34,459,425 |
| 11 | 6,939,897,856 | 654,729,075 |
| 12 | 188,666,182,784 | 13,749,310,575 |
| 13 | 5,617,349,020,544 | 316,234,143,225 |
| 14 | 181,790,703,209,728 | 7,905,853,580,625 |
| 15 | 6,353,726,042,486,112 | 213,458,046,676,875 |
| 16 | 238,513,970,965,250,048 | 6,190,283,353,629,375 |
| 17 | 9,571,020,586,418,569,216 | 191,898,783,962,510,625 |
| 18 | 408,837,905,660,430,516,224 | 6,332,659,870,762,850,625 |
| 19 | 18,522,305,410,364,568,764,416 | 221,643,095,476,699,771,875 |
| 20 | 887,094,711,304,094,583,095,296 | 8,200,794,532,637,891,559,375 |
| 21 | 44,782,218,857,751,551,087,214,592 | 319,830,986,772,877,770,815,625 |
| 22 | 2,376,613,641,928,796,906,249,519,104 | 13,113,070,457,687,988,603,440,625 |

- No of trees on 500 taxa ~ $1 \times 10^{\wedge 1074}$
- No of atoms in observable universe ~ $10^{\wedge} 80$


## Minimum evolutionary hypothesis

- In terms of topology : the star topology (not a binary tree)
- only assumes a most recent common ancestor



## Resolution

- Resolution = process of figuring out topology = generating more complicated evolutionary hypotheses
- Partial resolution = intermediate stage



## Outgroups

- When studying a group, we may want a control which is outside that group

bioinf. manchester.ac,uk
- One way to root a tree


## 2 schools of phylogeny reconstruction

- Distance based methods

Linkage tree for 9 population clusters showing genetic distances ( $F_{\text {ST }}$ )
(Cavalli-Sforza et al., 1994:80)


## 2 schools of phylogeny reconstruction

- Character based methods

ABCDE AXCDE AXYDE



Figure: A Phylogenetic Tree
molgen.mpg.de

## What kind of data can we use ?

- Both genetic and phenotype based models of evolution can be either character based or distance based.


## Distance based methods

- How to infer evolutionary relationships on the basis of some similarity measure ?
- Notion of a similarity measure
- eg. Curvature and length
- 2D Euclidean distance



## Genetic distance

- Similar to alignment score
- how far away are two orthologous sequences in sequence space ?
- alignment usually required (compare apple to apple)
- one option : pairwise entropy measure
- similar to Dobs = fraction of sites with substitutions in pairwise alignment


## Model based genetic distances

- Typically such simple measures cant capture complicated evolutionary models ( which model selection)
- However, in JC 69, all mutations are equally likely : can we use Dobs as distance measure ?



## Correction for multiple substitutions

- $\mathrm{A} \rightarrow \mathrm{T} \rightarrow \mathrm{C}$ (multiple mutations when we see one substitution)
- $\mathrm{A} \rightarrow \mathrm{T} \rightarrow \mathrm{A}$ (multiple mutations when we see no substitution) P( CORRECT fOR:


P( 2 or more mutation at each site)

$$
=\text { infinite series }[H W]
$$

$$
D_{J C}=-\frac{3}{4} \ln \left[1-\frac{4}{3} D_{\text {OBS }}\right]
$$



## Metric

- Mathematical notion of "distance"
- Intuitive properties

$$
\begin{array}{rlrl}
d(x, y) & >0 \quad \text { for } x \neq y \\
d(x, y) & =0 \quad \text { for } x=y \\
d(x, y) & =d(y, x) \quad \forall x, y \\
d(x, y) & \leq d(x, z)+d(y, z) \quad \forall x, y, z & & \text { (triangle inequality) } \quad \text { for all } \mathrm{C}
\end{array}
$$

- Space in which $x, y, z$ lives + metric definition
- Metric space

Additive metrics \& ultrametrics

- Additive / tree metric

- remember, route between 2 points on a tree is unique
- Ultrametric : $d(x, y) \leq \max (d(x, z), d(y, z))$


UT DALLXhere should we depict the set of metrics corr. to rooted trees?

## Tree metrics : 4 point condition

- Consider every quartet of leaves in the tree

- 3 ways to split into two subsets of 2 nodes:
$-A B$ \& CD, $A C$ \& BD, $A D$ \& $B C$
- $d(A B)+d(C D)<=\max (d(A C)+d(B D), d(A D)+d(B C))$
- $d(A C)+d(B D)<=\max (d(A B)+d(C D), d(A D)+d(B C))$
$-d(A D)+d(B C)<=\max (d(A C)+d(B D), d(A B)+d(C D))$

4PC jiff additivity
$\binom{n}{2}$ equations

$2 n-3$ variables

$$
\begin{aligned}
u+v & =d(A B) \\
u+w+x & =d(A C) \\
u+w+y & =d(A D) \quad \Rightarrow \\
v+w+x=d(B C) & \text { have } \\
v+w+y=d(B D) & \text { soln } \\
v+y & =d(C D)
\end{aligned}
$$

## Ultrametric $=3$ point condition

$$
d(x, y) \leq \max (d(x, z), d(y, z))
$$

- Consider every triplet of leaves A, B, C

- $\mathrm{d}(\mathrm{AB})<=\max (\mathrm{d}(\mathrm{AC}), \mathrm{d}(\mathrm{BC}))^{x+y}$
- $d(B C)<=\max (d(A B), d(A C))<(x+\omega+z$
- $d(A C)<=\max (d(A B), d(B C))=y+\omega+z)$


## 3 PC satisfied (hence 4 PC satisfied)


cloaked tree

## A few changes to 3 PC matrix may still be 3 PC matrix !



Some change mary still keel the clock

Some changes to 3 PC matrix may only satisfy 4 PC, not 3 PC


Is it clocked?

## Some changes to 4 PC matrix can preserve 4 PC condition



Sane changes can

Some changes to 4 PC : only triangle inequality satisfied


So, what changed ?

but, we can plot them in a metric space!

Arbitrary changes to the distance matrix
 We can no longer plot the prints in any metric space

## Comparison

## Distance metrics

## Stricter condition, set of matrices satisying this shrinks

## Triangle inequality

 / distance metric / additive metricTopology recovered, branch lengths may be inconsistent

Points may still be plotted in metric space

4 point condition

Fits unique unrooted (or one of its many equivalent rooted) tree with unique branch lengths

Tree may be drawn

3 point condition
/ ultrametric

Fits unique rooted, clocked tree with branch lengths

Tree drawn, one point on tree (root) is
equidistant from all leaves

URigynomial time greedy algorithm to reconstruct tree for all

## Clustering

- Partitioning some data points ( species or individuals in our case ) based on some metric or distance measure



## Clustering is hard

- k-means is NP hard



## Hierarchical clustering

- We may want to find such clusters at different levels : all the way from the whole data set to individual data points
- Notion of distance between clusters, between data points, and inbetween both categories



## Clustering $\leftarrow \rightarrow$ Tree topology



नtpoulus

Tree notation : the Newick format


## Top down \& bottom up

- Bottom up : Start with each element in its own cluster. Join clusters until one is left.

- Top down : Start with a single cluster of all elements. Split one cluster at each step until no clusters can be split further
DDALlas

$$
\Rightarrow
$$



## Simple bottom-up hierarchical clustering

- Start with every taxa in its own cluster
- Join the two nearest clusters at each step
- Recompute distance matrix
- Repeat until only one cluster is left

Distance between 2 clusters, cluster and an element based on distance betw. elements :

- Max -link
- Min - link



## Distance between clusters

- For single elements, consider them to be in their own cluster, then compute :
- max link $\max _{i, j} d\left(\lambda_{i}, y_{j}\right)$

$$
i, \dot{j}
$$

$-\min$ link $\min _{i, j} d\left(x_{i}, y_{j}\right)$


## Input / output of hierarchical clustering

- Input : a metric, the data points / heatmap
- Output : tree relation of the clusters, if possible branch lengths of tree (if further possible a root)



## Lets build some trees ...

## UPGMA

- Unweighted pair-group method with arithmetic mean
- In our language : bottom up hierarchical clustering with average-link distance between clusters
- Gives a clocked (hence rooted) tree ! (how ?)


## UPGMA

- See a beautiful handtrace here :
http://www.southampton.ac.uk/~re1u06/teaching/ upgma/
- However, UPGMA is clocked : lineage specific evolutionary rates (if 3 PC is violated ) cant be handled


## Neighbor joining

Make each taxa a tree of one species

- At each step identify 2 trees which are the most similar to each other AND further apart from the rest
- Make a new ancestral node and connect these 2 trees to it with different branch lengths
- Accounts for lineage specific evolutionary rates : makes an unclocked tree : so rooting is arbitrary if performed
- Update distance matrix

URepeatuntil two trees are left, then connect them

## Neighbor joining



Neighbor joining

- Notion of how different from the rest of the leaves one element is :

$$
u_{i}=\sum_{j} d_{i j} /(n-2)
$$

d

$u$



PICK i \& $j$ based on Lowest $m_{i j}$

Neighbor joining

- Joining the two trees :



## Neighbor joining

- If $i$ th and $j$ th trees were joined, then the distance from the new ancestor to any other remaining tree $k$ is updated, and the original 2 trees are removed from the distance matrix :


$$
d_{n, k}=\left(d_{i k}+d_{j k}-d_{i j}\right) / 2
$$

- Originally, when each tree has single element, -UT Dhe distances matrix are just distances betw taxa


## Neighbor joining

- Another beautiful handtrace :
- http://www.cbs.dtu.dk/dtucourse/cookbooks/gor m/27615/molevol.powerpoints/MoIEvolClass05.p pt


## Bottom up hierarchical clustering

- Decisions:
- how to identify neighbors and merge them together into a single tree / cluster
- how to assign branch lengths
- how to update the distance matrix once merging two clusters is performed


## In practice ...

- Most distance matrices don't follow 3PC or 4PC
- Converting an arbitrary distance matrix to 3PC or 4 PC is difficult and may change the nature of relationships unfaithfully
- Forcing a hierarchical clustering on a non conforming matrix leads to wrong branch lengths and / or topology
- Squared error = sum of squares of difference between distance matrix and distance on tree for each taxon pair


## Then, what?

- Pick any tree (!!! ) - you may start with NJ tree
- Calculate squared error between the original distance matrix and distances on the tree
- Change the tree slightly : if the squared error decreases, keep the changes
- Keep doing this step till no more improvement

How to change a tree "slightly" ? (coming up next class)

## Notion of a "better" tree

- Distance methods do not provide a rational notion of whether one tree is better than the other
- Most distance methods produce one tree
- Comparing trees across methods is difficult
- Minimum evolution : minimal sum of branch lengths is better
- Premise may not be true for distant species


## Figuring out the ancestral state

- How to figure out the ancestral state ?
- going from phenotype to distance is relatively easy (define a metric)
- going from distance to phenotype is difficult ( may be ambiguous )



## To the rescue ...

- Character based methods


## ABCDE AXCDE AXYDE



Figure: A Phylogenetic Tree

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