

MPE, MAP AND APPROXIMATIONS

Statistical Methods in AI/ML

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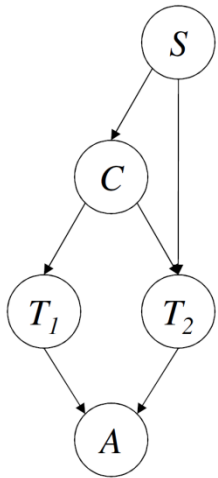


What we will cover?

- MPE= most probable explanation
 - The tuple with the highest probability in the joint distribution $\Pr(X|e)$
- MAP=maximum a posteriori
 - Given a subset of variables Y , the tuple with the highest probability in the distribution $P(Y|e)$
- Exact Algorithms
 - Variable elimination
 - DFS search
 - Branch and Bound Search
- Approximations
 - Upper bounds
 - Local search



Running Example: Cheating in UTD CS Population



S	C	T_2	$\theta_{t_2 c,s}$
male	yes	+ve	.80
male	yes	-ve	.20
male	no	+ve	.20
male	no	-ve	.80
female	yes	+ve	.95
female	yes	-ve	.05
female	no	+ve	.05
female	no	-ve	.95

T_1	T_2	A	$\theta_{a t_1,t_2}$
+ve	+ve	yes	1
+ve	+ve	no	0
+ve	-ve	yes	0
+ve	-ve	no	1
-ve	+ve	yes	0
-ve	+ve	no	1
-ve	-ve	yes	1
-ve	-ve	no	0

S	θ_s
male	.55
female	.45

S	C	$\theta_{c s}$
male	yes	.05
male	no	.95
female	yes	.01
female	no	.99

C	T_1	$\theta_{t_1 c}$
yes	+ve	.80
yes	-ve	.20
no	+ve	.20
no	-ve	.80

Sex (S), Cheating (C), Tests (T1 and T2) and Agreement (A)



Most likely instantiations

- MPE = Most likely assignment to all non-evidence variables (given evidence)
- MAP = Most likely assignment to a subset of non-evidence variables (given evidence)
- A person takes a test and the test administrator says
 - The two tests agree ($A = \text{true}$)
- Query: Most likely instantiation of Sex and Cheating given evidence $A = \text{true}$
- Is this a MAP or an MPE problem?
- Answer: Sex=male and Cheating=no.



MPE vs MAP: Properties

- MPE is a special case of MAP
- Hardness
 - Computing MPE is NP-hard (Max-product problem)
 - Computing MAP is NP^{PP}-hard (Max-sum-product problem believed to be much harder than NP-hard)
- MPE projected on to the MAP variables does not yield the correct answer.
 - MPE given A=yes
 - S=female, C=no, T₁=negative and T₂=negative
 - MPE projected on MAP variables S and C
 - S=female, C=no is incorrect!
 - MAP given A=yes
 - S=male, C=no is correct!
- We will distinguish between
 - MPE and MAP probabilities
 - MPE and MAP instantiations



Bucket Elimination for MPE

- Same schematic algorithm as before
- Replace “elimination operator” by “maximization operator”

MAX_S

S	C	Value
male	yes	0.05
male	no	0.95
female	yes	0.01
female	no	0.99

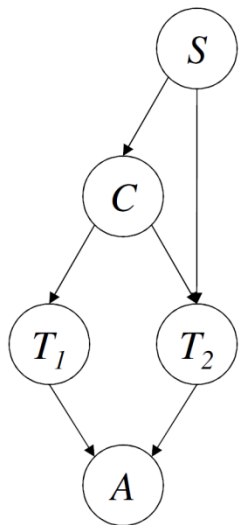
=

C	Value
yes	0.05
no	0.99

Collect all instantiations that agree on all other variables except S and return the maximum value among them.

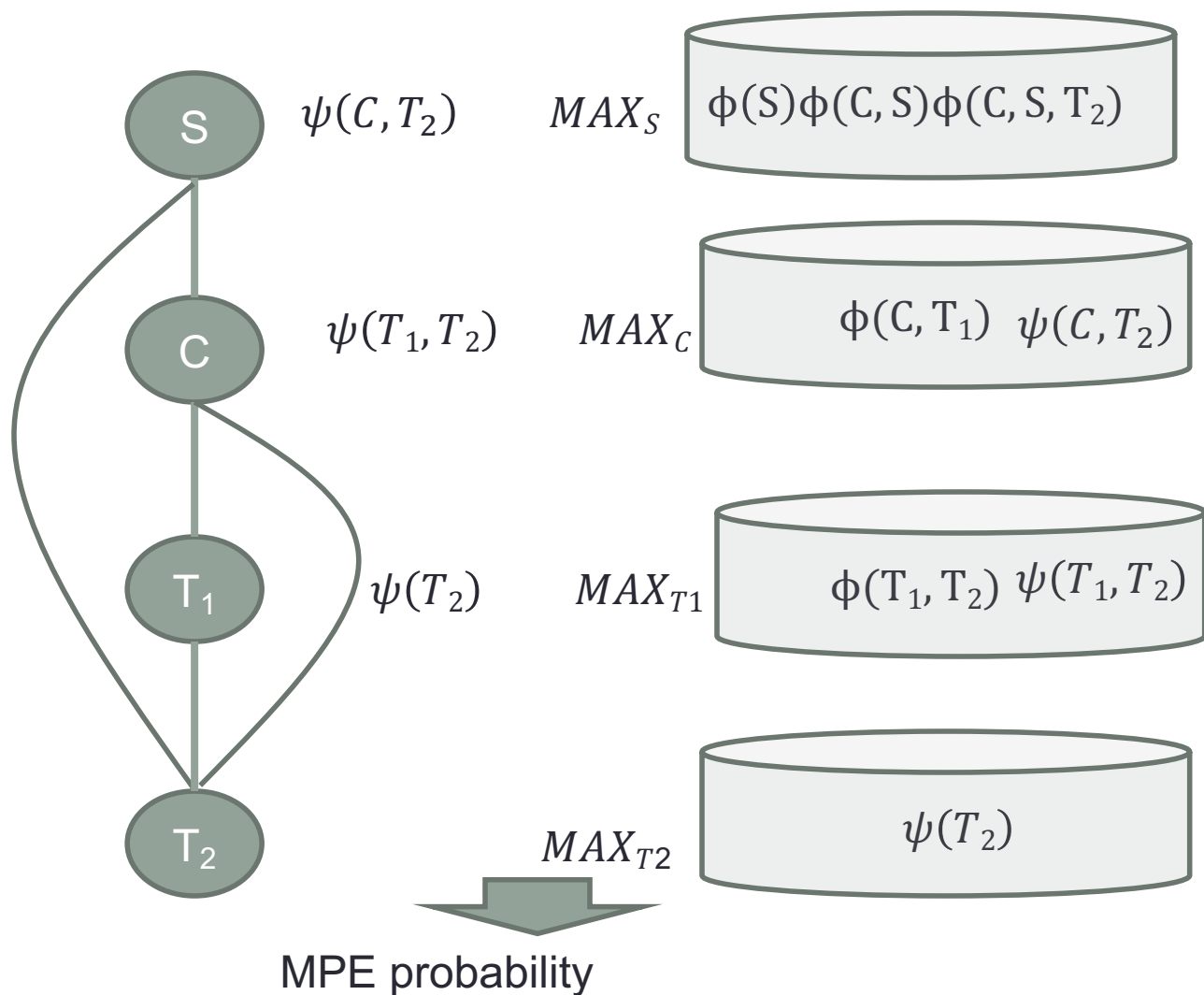


Bucket elimination: order (S, C, T1, T2)



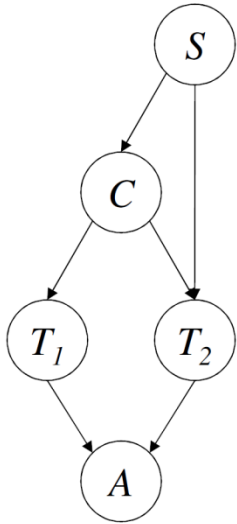
Factors: $\phi(S)$
 $\phi(C, S)$
 $\phi(C, S, T_2)$
 $\phi(C, T_1)$
 $\phi(T_1, T_2)$

Evidence: A=true



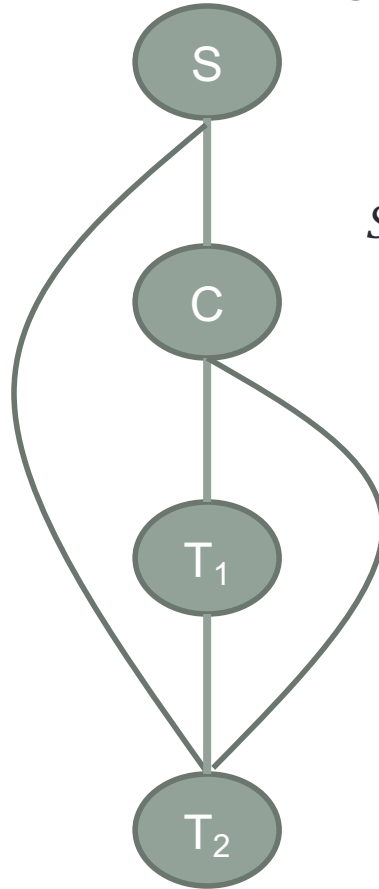


Bucket elimination: Recovering MPE tuple



Factors: $\phi(S)$
 $\phi(C, S)$
 $\phi(C, S, T_2)$
 $\phi(C, T_1)$
 $\phi(T_1, T_2)$

Evidence: $A = \text{true}$



Set $T_2 = -ve, C = no$
 $\max(S \text{ tuple}) = ?$
 $S = female$

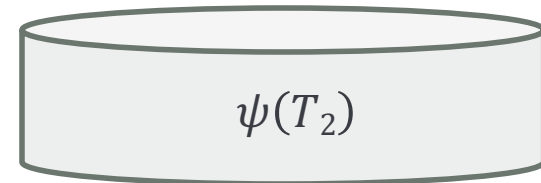
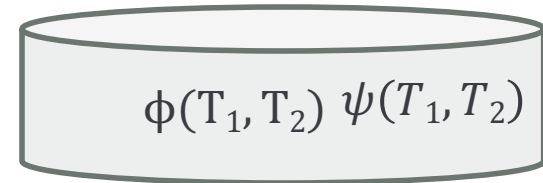
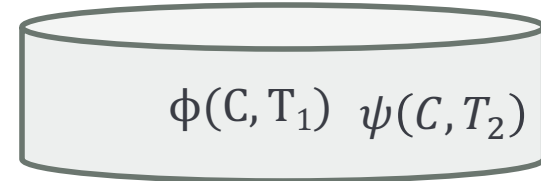
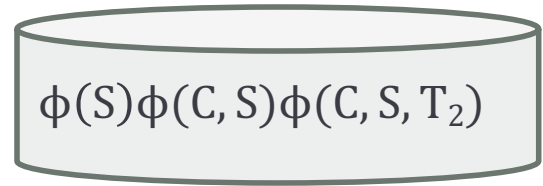
Set $T_2 = -ve, T_1 = -ve$
 $\max(C \text{ tuple}) = ?$
 $C = no$

Set $T_2 = -ve$
 $\max(T_1 \text{ tuple}) = ?$
 $T_1 = -ve$

$\max(T_2 \text{ tuple}) = ?$
 $T_2 = -ve$



MPE probability

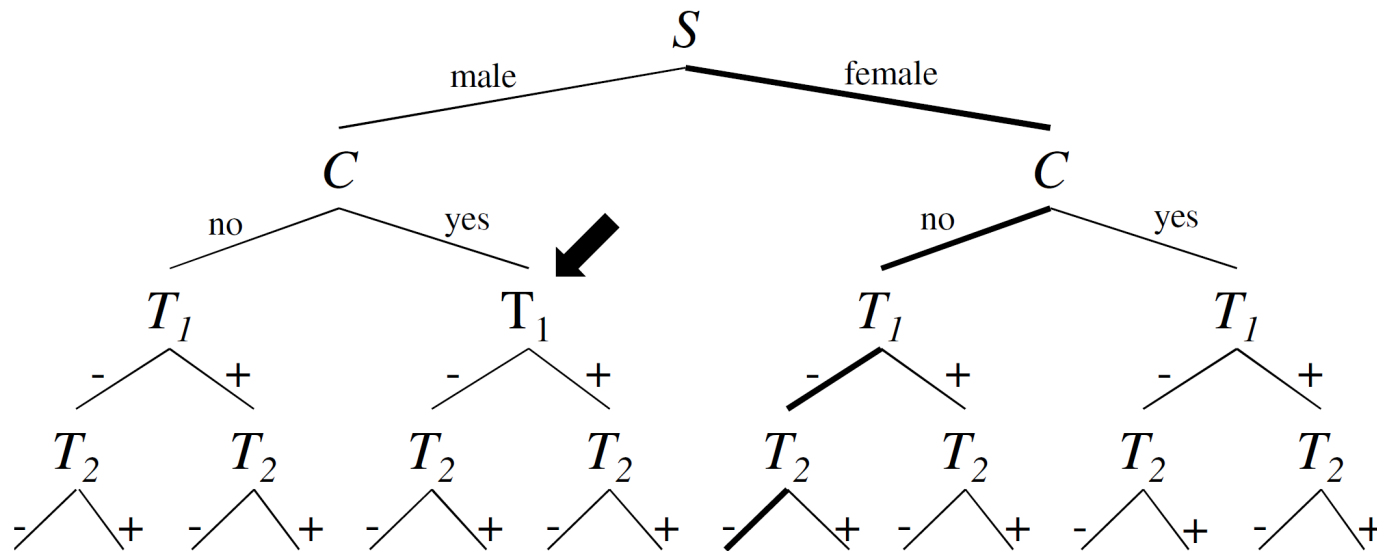


Bucket elimination: MPE vs PE (Z)

- Maximization vs summation
- Complexity: Same
 - Time and Space exponential in the width (w) of the given order:
 $O(n \exp(w+1))$ timewise and $O(n \exp(w))$ spacewise.

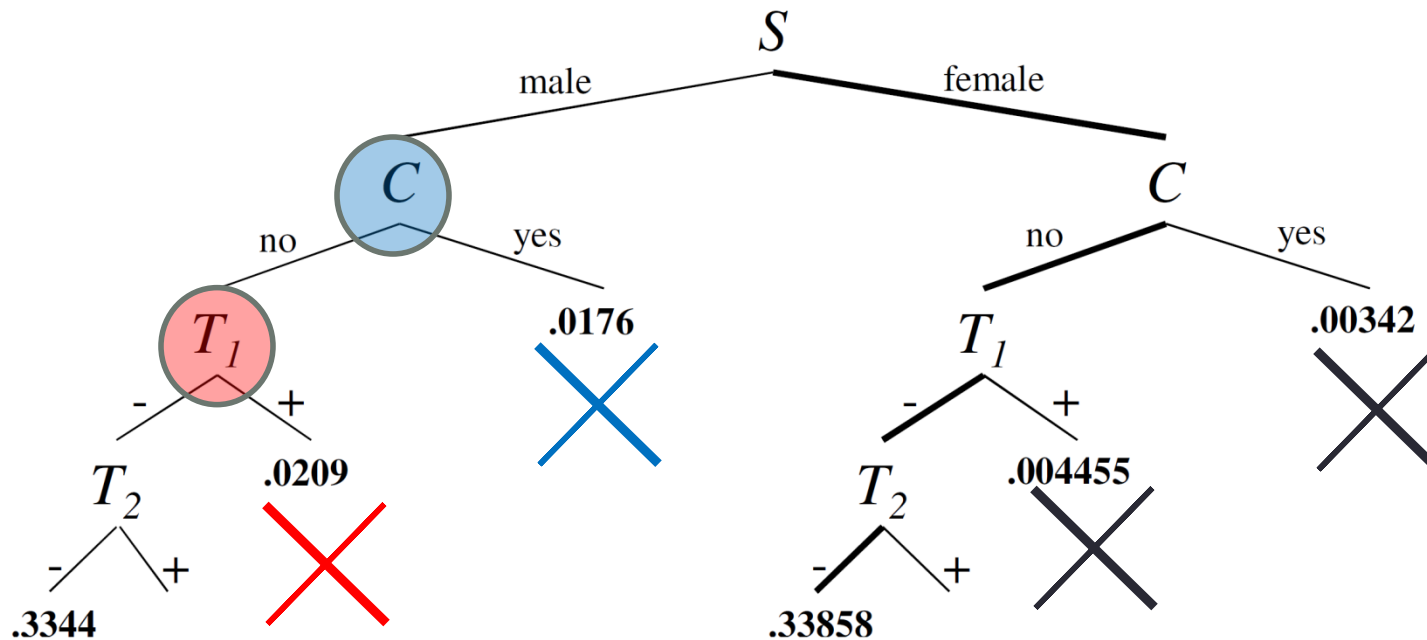


OR search for MPE



- At leaf nodes compute probabilities by taking product of factors
- Select the path with the highest leaf probability

Branch and Bound Search



- Let us say we have a method to upper bound MPE at each node
- Prune nodes which have smaller upper bound than the current MPE solution
- Amount of pruning depends on the quality of the upper bound. Lower the upper bound (i.e., better the upper bound), better the pruning.

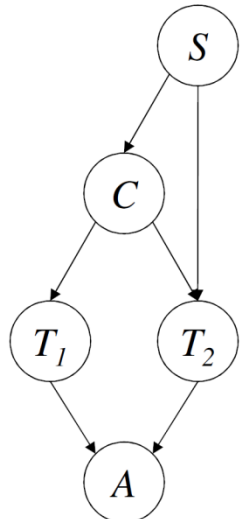


Mini-Bucket Approximation: Idea

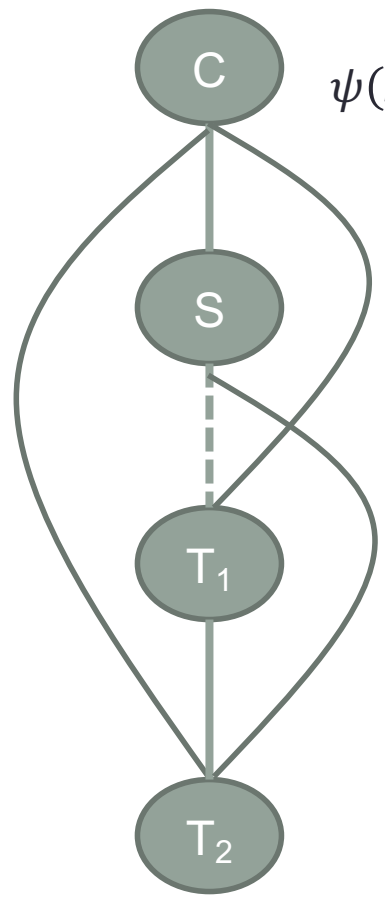
Split a bucket into mini-buckets => bound complexity

$$\begin{array}{c}
 \mathbf{bucket(Y) =} \\
 \{ \phi_1, \dots, \phi_r, \phi_{r+1}, \dots, \phi_n \} \\
 \underbrace{\hspace{10em}} \\
 \begin{array}{ccc}
 \swarrow & g = MAX_Y \left(\prod_{i=1}^n \phi_i \right) & \searrow \\
 \{ \phi_1, \dots, \phi_r \} & & \{ \phi_{r+1}, \dots, \phi_n \} \\
 h_1 = MAX_Y \left(\prod_{i=1}^r \phi_i \right) & & h_2 = MAX_Y \left(\prod_{i=r+1}^n \phi_i \right)
 \end{array} \\
 \\
 g \leq h_1 \times h_2
 \end{array}$$

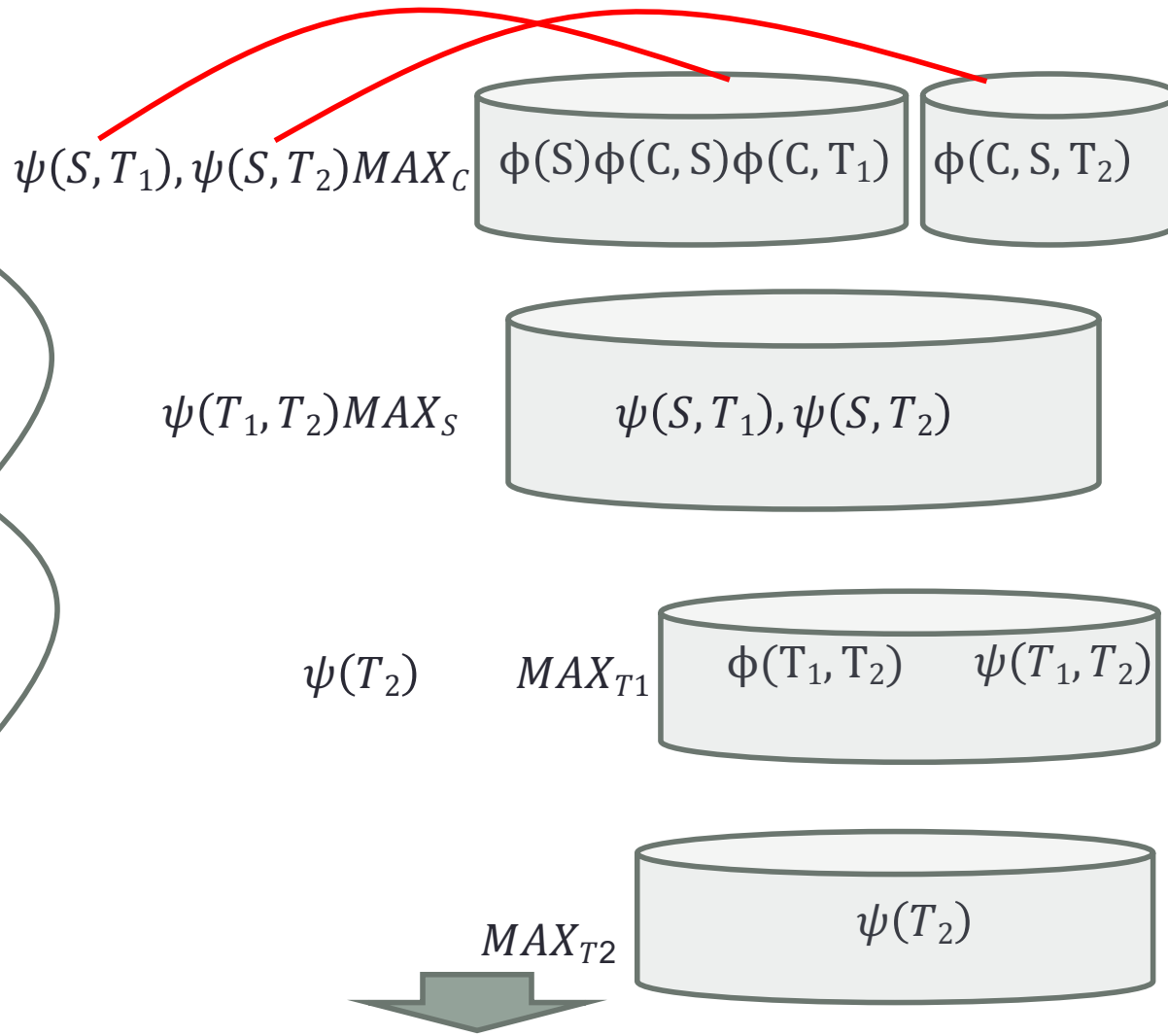
Mini Bucket elimination: (max-size=3 vars)



Factors: $\phi(S)$
 $\phi(C, S)$
 $\phi(C, S, T_2)$
 $\phi(C, T_1)$
 $\phi(T_1, T_2)$



Evidence: A=true



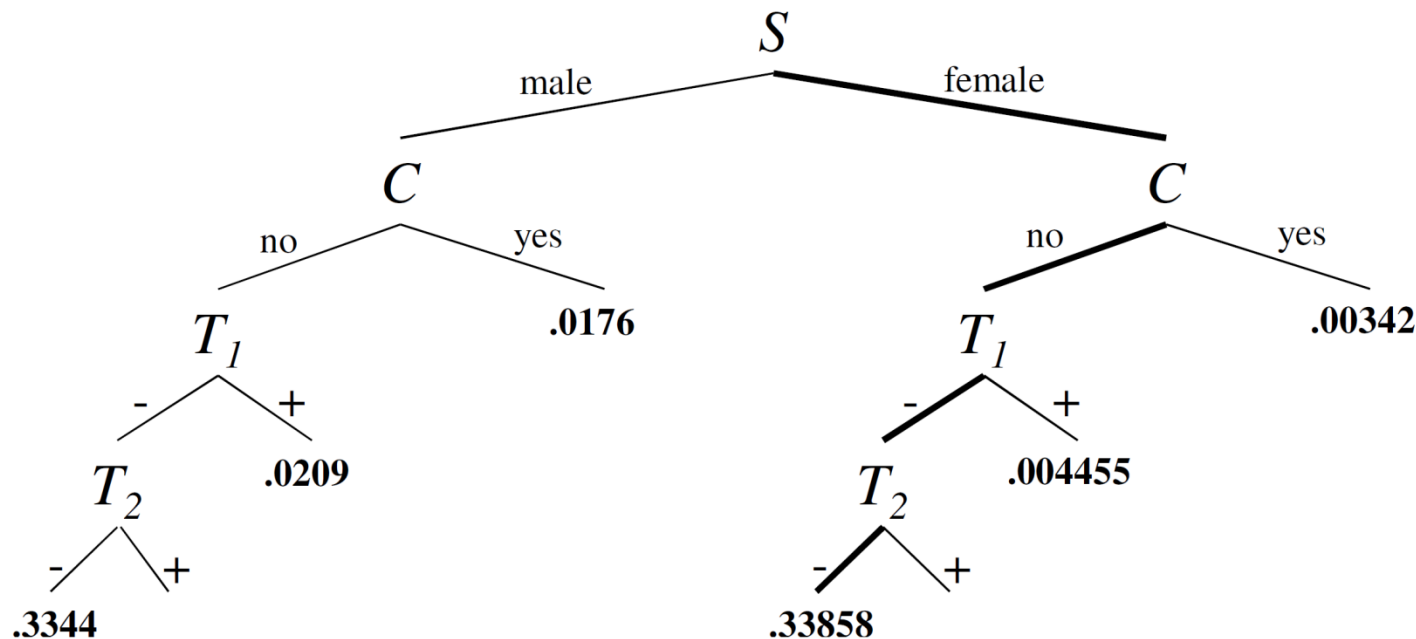
Upper bound on the MPE probability

Mini-bucket (i-bounds)

- A parameter “i” which controls the size of (number of variables in) each mini-bucket
- Algorithm exponential in “i” : $O(n \exp(i))$
- Example
 - $i=2$, quadratic
 - $i=3$, cubed
 - etc
- Higher the i-bound, better the upper bound
- In practice, can use i-bounds as high as 22-25.



Branch and Bound Search



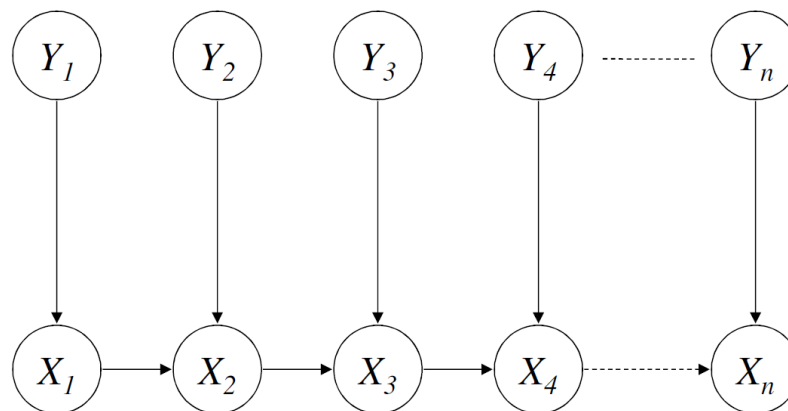
- Run MBE at each branch point.
- Prune nodes which have smaller upper bound than the current MPE solution
- Radu Marinescu's PhD thesis (AND/OR branch and bound search plus more):
<https://www.ics.uci.edu/~dechter/publications/r158.pdf>

Computing MAP probabilities: Bucket Elimination

- Given MAP variables “M” and evidence be “e”
- Can compute the MAP probability using bucket elimination by first summing out all non-MAP variables, and then maximizing out MAP variables.
- By summing out non-MAP variables we are effectively computing the joint marginal $\Pr(M, e)$ in factored form.
- By maximizing out MAP variables M, we are effectively solving an MPE problem over the resulting marginal.
- The variable order used in `BE_MAP` is constrained as it requires MAP variables M to appear last in the order.
- Best case: `BE_MAP` is exponential in constrained treewidth which is the minimum width over (constrained) orders in which non-MAP variables are ordered before MAP variables.



MAP and constrained width



- Treewidth = 2
- MAP variables = $\{Y_1, \dots, Y_n\}$
- Any order in which M variables come first has width greater than or equal to n
- BE_MPE is exponential in 3 and BE_MAP is exponential in $O(\bar{n})$.



MAP by branch and bound search

- MAP can be solved using depth-first branch-and-bound search, just as we did for MPE.
- Algorithm BB_MAP resembles the one for computing MPE with two exceptions.
- Exception 1: The search space consists only of the MAP variables
- Exception 2: We use a version of MBE_MAP for computing the bounds
 - Order all MAP variables after the non-MAP variables.



MAP by Local Search

- Given a network with n variables and an elimination order of width w
 - Complexity: $O(r \text{ nexp}(w+1))$ where “ r ” is the number of local search steps
- Start with an initial random instantiation of MAP variables
- Neighbors of the instantiation “ m ” are instantiations that result from changing the value of one variable in “ m ”
- Score for neighbor “ m ”: $\text{Pr}(m,e)$
- How to compute $\text{Pr}(m,e)$?
 - Bucket elimination.

MAP: Local search algorithm

LS_MAP(\mathcal{N} , \mathbf{M} , \mathbf{e})

input:

\mathcal{N} : Bayesian network
 \mathbf{M} : some network variables
 \mathbf{e} : evidence ($\mathbf{E} \cap \mathbf{M} = \emptyset$)

output: instantiation \mathbf{m} of \mathbf{M} which (approximately) maximizes $\Pr(\mathbf{m}|\mathbf{e})$.

main:

```
1:  $r \leftarrow$  number of local search steps
2:  $P_f \leftarrow$  probability of randomly choosing a neighbor
3:  $\mathbf{m}^* \leftarrow$  some instantiation of variables  $\mathbf{M}$  {best instantiation}
4:  $\mathbf{m} \leftarrow \mathbf{m}^*$  {current instantiation}
5: for  $r$  times do
6:    $p \leftarrow$  random number in  $[0, 1]$ 
7:   if  $p \leq P_f$  then
8:      $\mathbf{m} \leftarrow$  randomly selected neighbor of  $\mathbf{m}$ 
9:   else
10:    compute the score  $\Pr(\mathbf{m} - X, x, \mathbf{e})$  for each neighbor  $\mathbf{m} - X, x$ 
11:    if no neighbor has a higher score than the score for  $\mathbf{m}$  then
12:       $\mathbf{m} \leftarrow$  randomly selected neighbor of  $\mathbf{m}$ 
13:    else
14:       $\mathbf{m} \leftarrow$  a neighbor of  $\mathbf{m}$  with a highest score
15:    end if
16:  end if
17:  if  $\Pr(\mathbf{m}, \mathbf{e}) > \Pr(\mathbf{m}^*, \mathbf{e})$ , then  $\mathbf{m}^* \leftarrow \mathbf{m}$ 
18: end for
19: return  $\mathbf{m}^*$ 
```

Recap

- Exact MPE and MAP
 - Bucket elimination
 - Branch and Bound Search
- Approximations
 - Mini bucket elimination
 - Branch and Bound Search
 - Local Search