Sampling Techniques for Probabilistic and Deterministic Graphical models

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### **Overview**

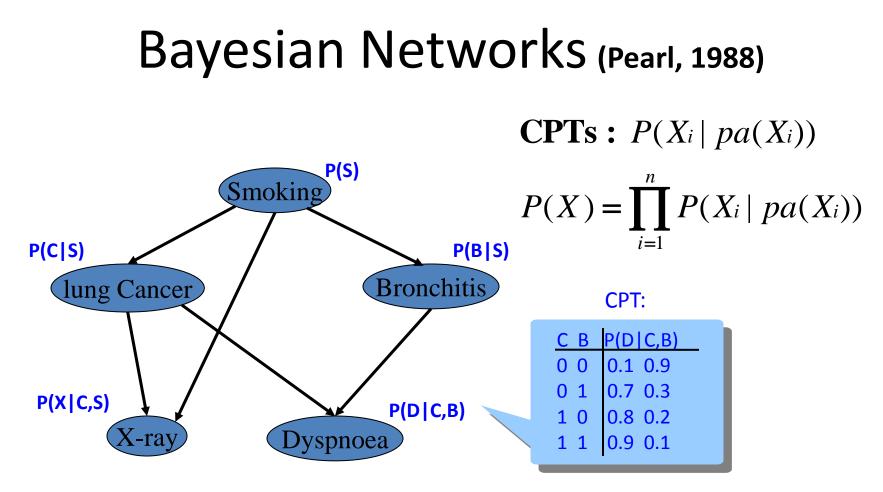
- 1. Probabilistic Reasoning/Graphical models
- 2. Importance Sampling
- 3. Markov Chain Monte Carlo: Gibbs Sampling
- 4. Sampling in presence of Determinism
- 5. Rao-Blackwellisation
- 6. AND/OR importance sampling



- 1. Probabilistic Reasoning/Graphical models
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- 5. Cutset-based Variance Reduction
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# Probabilistic Reasoning; Graphical models

- Graphical models:
  - Bayesian network, constraint networks, mixed network
- Queries
- Exact algorithm
  - using inference,
  - search and hybrids
- Graph parameters:
  - tree-width, cycle-cutset, w-cutset



P(S, C, B, X, D) = P(S) P(C|S) P(B|S) P(X|C,S) P(D|C,B)

#### **Belief Updating:**

P (lung cancer=yes | smoking=no, dyspnoea=yes ) = ?
Probability of evidence:
P ( smoking=no, dyspnoea=yes ) = ?

### Queries

Probability of evidence (or partition function)

$$P(e) = \sum_{X - \operatorname{var}(e)} \prod_{i=1}^{n} P(x_i \mid pa_i) \mid_e \qquad Z = \sum_X \prod_i \psi_i(C_i)$$

Posterior marginal (beliefs):

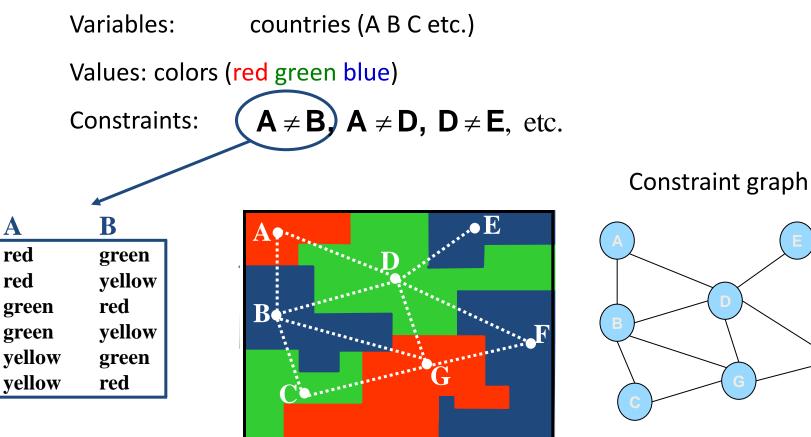
$$P(x_i \mid e) = \frac{P(x_i, e)}{P(e)} = \frac{\sum_{X = \text{var}(e) = X_i} \prod_{j=1}^n P(x_j \mid pa_j)|_e}{\sum_{X = \text{var}(e)} \prod_{j=1}^n P(x_j \mid pa_j)|_e}$$

Most Probable Explanation

$$\overline{\mathbf{x}}^* = \arg\max_{\overline{\mathbf{x}}} \mathbf{P}(\overline{\mathbf{x}}, \mathbf{e})$$

# **Constraint Networks**

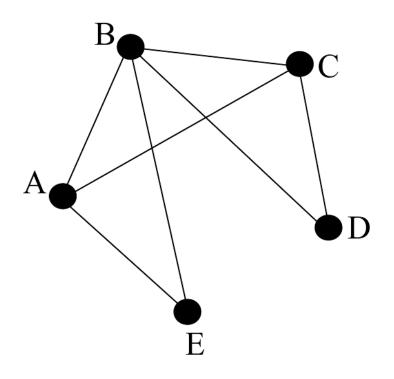
#### Map coloring

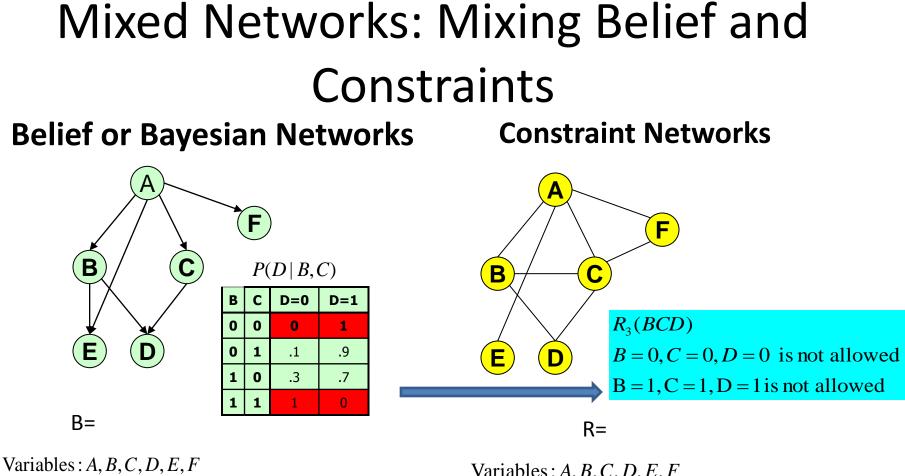


Task: find a solution Count solutions, find a good one

# **Propositional Satisfiability**

 $\varphi = \{ (\neg C), (A \lor B \lor C), (\neg A \lor B \lor E), (\neg B \lor C \lor D) \}.$ 





Domains :  $D_A = D_B = D_C = D_D = D_E = D_F = \{0,1\}$ CPTS : P(A), P(B | A), P(C | A), P(D | B, C)

 $P(E \mid A, B), P(F \mid A)$ 

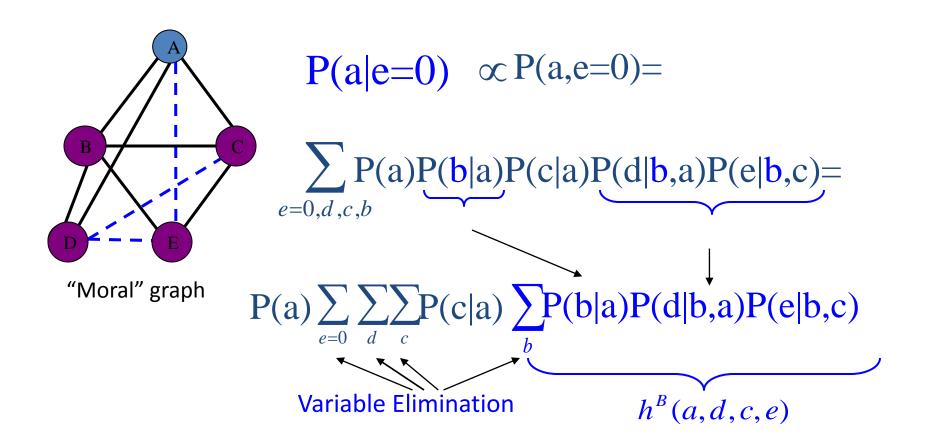
Variables : A, B, C, D, E, FDomains :  $D_A = D_B = D_C = D_D = D_E = D_F = \{0,1\}$ Constraints :  $R_1(ABC), R_2(ACF), R_3(BCD), R_4(A, E)$ 

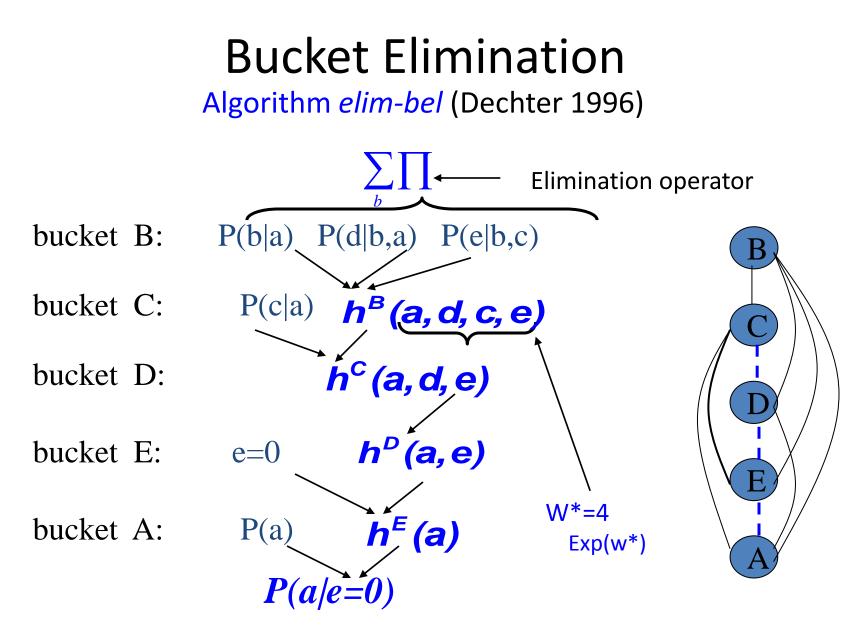
Expresses the set of solutions : sol(R)Constraints could be specified externally or may occur as zeros in the Belief network

Same queries (e.g., weighted counts)

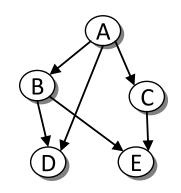
$$M = \sum_{x \in sol(R)} P_B(x)$$

## **Belief Updating**





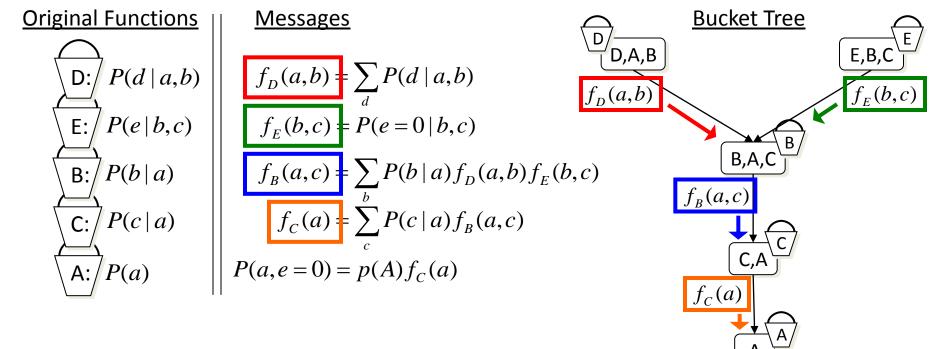
## **Bucket Elimination**



Query: 
$$P(a | e = 0) \propto P(a, e = 0)$$
 Elimination Order: d,e,b,c  

$$P(a, e = 0) = \sum_{c,b,e=0,d} P(a)P(b | a)P(c | a)P(d | a,b)P(e | b,c)$$

$$= P(a)\sum_{c} P(c | a)\sum_{b} P(b | a)\sum_{e=0} P(e | b,c)\sum_{d} P(d | a,b)$$

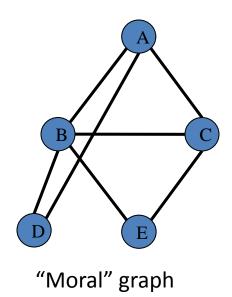


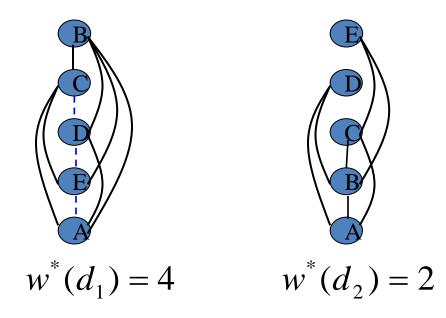
#### **Complexity of Elimination**

#### $O(n \exp(w^*(d)))$

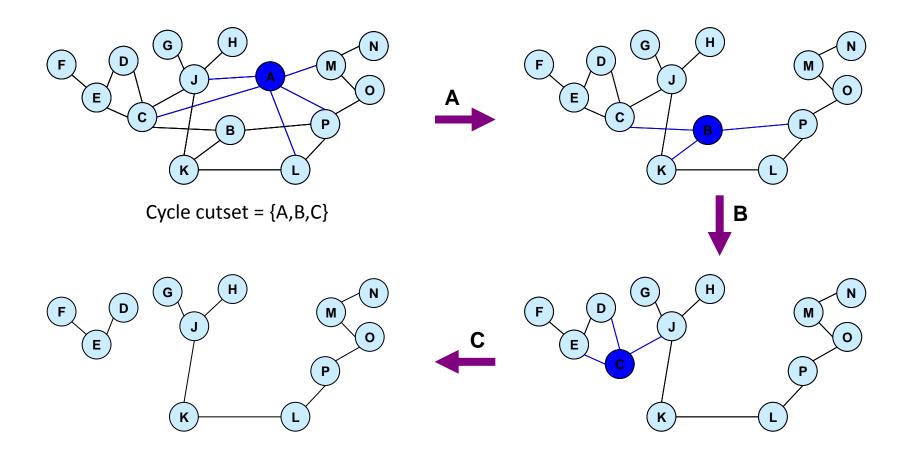
 $w^*(d)$  – the induced width of moral graph along ordering d

The effect of the ordering:



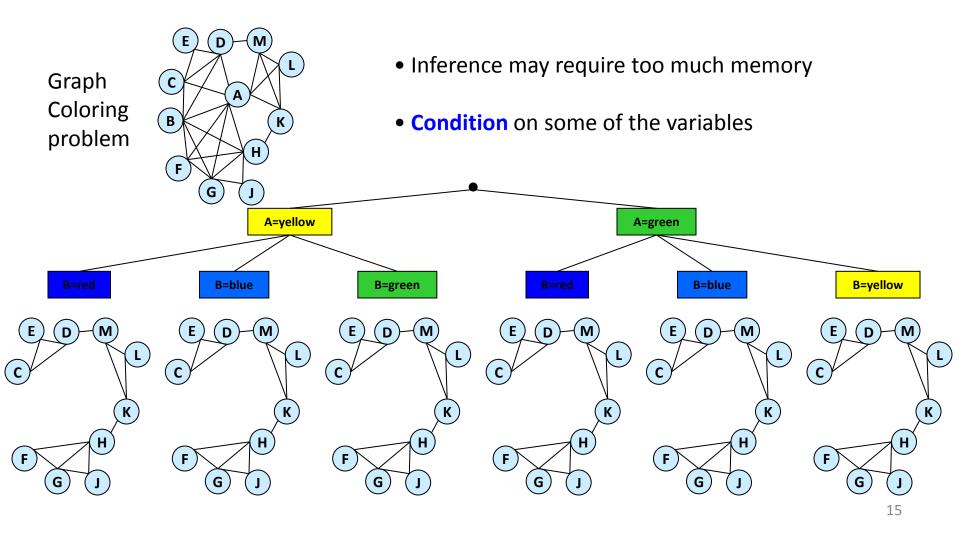


#### **Cutset-Conditioning**

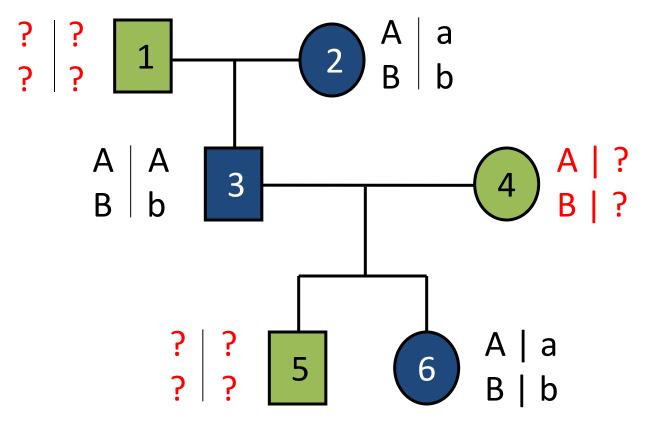


#### Search Over the Cutset

Space: exp(w): w is a user-controled parameter Time: exp(w+c(w))



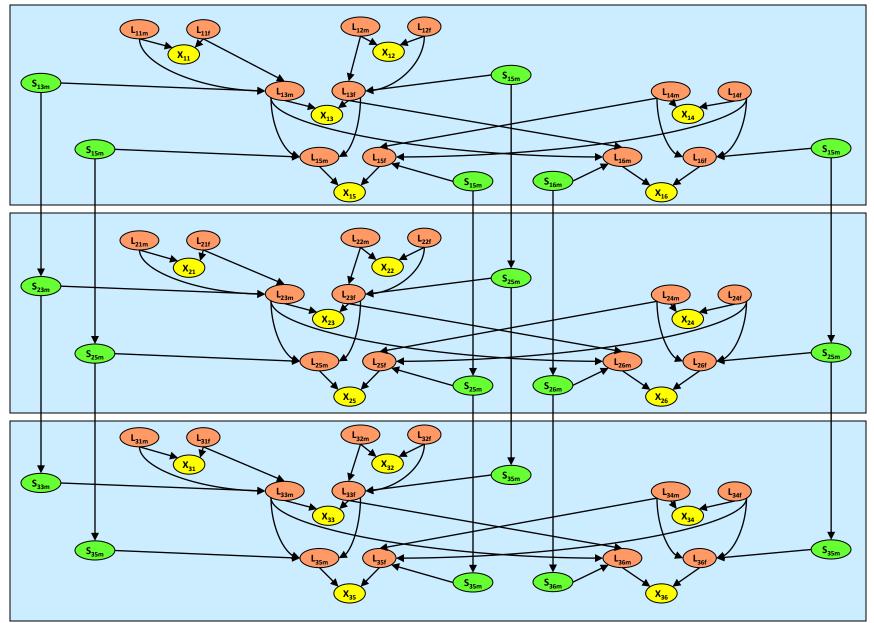
## Linkage Analysis





- 6 individuals
- Haplotype: {2, 3}
- Genotype: {6}
- Unknown

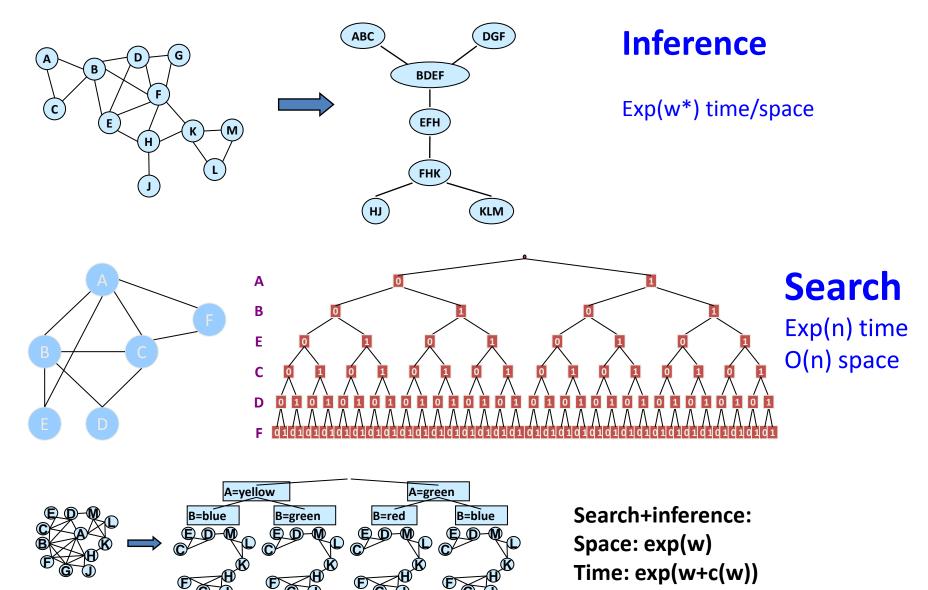
#### Linkage Analysis: 6 People, 3 Markers



# Applications

- Determinism: More Ubiquitous than you may think!
- Transportation Planning (Liao et al. 2004, Gogate et al. 2005)
  - Predicting and Inferring Car Travel Activity of individuals
- Genetic Linkage Analysis (Fischelson and Geiger, 2002)
  - associate functionality of genes to their location on chromosomes.
- Functional/Software Verification (Bergeron, 2000)
  - Generating random test programs to check validity of hardware
- First Order Probabilistic models (Domingos et al. 2006, Milch et al. 2005)
  - Citation matching

#### Inference vs Conditioning-Search



## Approximation

- Since inference, search and hybrids are too expensive when graph is dense; (high treewidth) then:
- Bounding inference:
  - mini-bucket and mini-clustering
  - Belief propagation
- Bounding search:
  - Sampling
- Goal: an anytime scheme

### Approximation

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# Outline

- Definitions and Background on Statistics
- Theory of importance sampling
- Likelihood weighting
- State-of-the-art importance sampling techniques

# A sample

 Given a set of variables X={X<sub>1</sub>,...,X<sub>n</sub>}, a sample, denoted by S<sup>t</sup> is an instantiation of all variables:

$$S^t = (x_1^t, x_2^t, \dots, x_n^t)$$

How to draw a sample ? Univariate distribution

- Example: Given random variable X having domain {0, 1} and a distribution P(X) = (0.3, 0.7).
- Task: Generate samples of X from P.
- How?
  - draw random number  $r \in [0, 1]$
  - If (r < 0.3) then set X=0
  - Else set X=1

# How to draw a sample? Multi-variate distribution

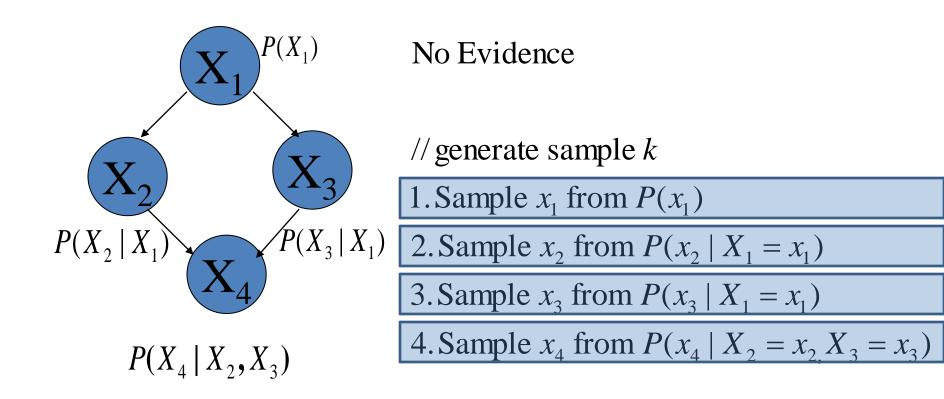
- Let X={X<sub>1</sub>,..,X<sub>n</sub>} be a set of variables
- Express the distribution in product form

 $P(X) = P(X_1) \times P(X_2 | X_1) \times ... \times P(X_n | X_1, ..., X_{n-1})$ 

- Sample variables one by one from left to right, along the ordering dictated by the product form.
- Bayesian network literature: Logic sampling

### Logic sampling (example)

 $P(X_1, X_2, X_3, X_4) = P(X_1) \times P(X_2 | X_1) \times P(X_3 | X_1) \times P(X_4 | X_2, X_3)$ 



#### **Expected value and Variance**

**Expected value**: Given a probability distribution P(X)and a function g(X) defined over a set of variables  $X = \{X_1, X_2, ..., X_n\}$ , the expected value of g w.r.t. P is

$$E_P[g(x)] = \sum_x g(x)P(x)$$

Variance: The variance of g w.r.t. P is:

$$Var_{P}[g(x)] = \sum_{x} [g(x) - E_{P}[g(x)]]^{2} P(x)$$

## Monte Carlo Estimate

#### • Estimator:

- An estimator is a function of the samples.
- It produces an estimate of the unknown parameter of the sampling *distribution*.

Given i.i.d. samples  $S^1, S^2, \dots S^T$  drawn from P,

the Monte carlo estimate of  $E_P[g(x)]$  is given by:

$$\hat{g} = \frac{1}{T} \sum_{t=1}^{T} g(S^t)$$

# Example: Monte Carlo estimate

- Given:
  - A distribution P(X) = (0.3, 0.7).
  - -g(X) = 40 if X equals 0 = 50 if X equals 1.
- Estimate  $E_P[g(x)] = (40x0.3+50x0.7)=47$ .
- Generate k samples from P: 0,1,1,1,0,1,1,0,1,0

$$\hat{g} = \frac{40 \times \# samples(X = 0) + 50 \times \# samples(X = 1)}{\# samples}$$
$$= \frac{40 \times 4 + 50 \times 6}{10} = 46$$

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# Importance sampling: Main idea

- Transform the probabilistic inference problem into the problem of computing the expected value of a random variable w.r.t. to a distribution Q.
- Generate random samples from Q.
- Estimate the expected value from the generated samples.

# Importance sampling for P(e)

Let  $Z = X \setminus E$ ,

Let Q(Z) be a (proposal) distribution, satisfying

 $P(z,e) > 0 \Longrightarrow Q(z) > 0$ 

Then, we can rewrite P(e) as :

$$P(e) = \sum_{z} P(z, e) = \sum_{z} P(z, e) \frac{Q(z)}{Q(z)} = E_{Q} \left[ \frac{P(z, e)}{Q(z)} \right] = E_{Q} [w(z)]$$

Monte Carlo estimate :

$$\hat{P}(e) = \frac{1}{T} \sum_{t=1}^{T} w(z^t)$$
, where  $z^t \leftarrow Q(Z)$ 

# Properties of IS estimate of P(e)

• Convergence: by law of large numbers

$$\hat{P}(e) = \frac{1}{T} \sum_{i=1}^{T} w(z^{i}) \xrightarrow{a.s.} P(e) \text{ for } T \to \infty$$

• Unbiased.

$$E_Q[\hat{P}(e)] = P(e)$$

• Variance:

$$Var_{Q}\left[\hat{P}(e)\right] = Var_{Q}\left[\frac{1}{T}\sum_{i=1}^{N}w(z^{i})\right] = \frac{Var_{Q}[w(z)]}{T}$$

# Properties of IS estimate of P(e)

Mean Squared Error of the estimator

$$MSE_{Q}[\hat{P}(e)] = E_{Q}[(\hat{P}(e) - P(e))^{2}]$$

$$= (P(e) - E_{Q}[\hat{P}(e)])^{2} + Var_{Q}[\hat{P}(e)]$$

$$= Var_{Q}[\hat{P}(e)]$$
This quantity enclosed in the brackets is zero because the expected value of the estimator equals the expected value of g(x)

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# Estimating P(X<sub>i</sub>|e)

Let  $\delta_{x_i}(z)$  be a dirac-delta function, which is 1 if z contains  $x_i$  and 0 otherwise.

$$P(x_{i} | e) = \frac{P(x_{i}, e)}{P(e)} = \frac{\sum_{z} \delta_{x_{i}}(z) P(z, e)}{\sum_{z} P(z, e)} = \frac{E_{Q} \left[ \frac{\delta_{x_{i}}(z) P(z, e)}{Q(z)} \right]}{E_{Q} \left[ \frac{P(z, e)}{Q(z)} \right]}$$

Idea : Estimate numerator and denominator by IS.

Ratio estimate : 
$$\overline{P}(x_i | e) = \frac{\hat{P}(x_i, e)}{\hat{P}(e)} = \frac{\sum_{k=1}^{T} \delta_{x_i}(z^k) w(z^k, e)}{\sum_{k=1}^{T} w(z^k, e)}$$
  
Estimate is biased :  $\mathbb{E}[\overline{P}(x_i | e)] \neq P(x_i | e)$ 

# Properties of the IS estimator for $P(X_i | e)$

Convergence: By Weak law of large numbers

 $\overline{P}(x_i \mid e) \rightarrow P(x_i \mid e) \text{ as } T \rightarrow \infty$ 

• Asymptotically unbiased

$$\lim_{T \to \infty} E_P[\overline{P}(x_i | e)] = P(x_i | e)$$

- Variance
  - Harder to analyze
  - Liu suggests a measure called "Effective sample size"

#### Effective Sample size

$$P(x_i \mid e) = \sum_{z} g_{x_i}(z) P(z \mid e)$$

Given samples from P(z | e), we can estimate  $P(x_i | e)$  using :

$$\hat{P}(x_i/e) = \frac{1}{T} \sum_{j=1}^{T} g_{x_i}(z^t) \qquad \qquad \text{Ideal estimator}$$

$$Define: ESS(Q,T) = \frac{T}{1 + \operatorname{var}_{Q}[w(z)]} \longrightarrow$$
$$\frac{Var_{P}[\hat{P}(x_{i} \mid e)]}{Var_{Q}[\overline{P}(x_{i} \mid e)]} \approx \frac{T}{ESS(Q,T)}$$

Measures how much the estimator deviates from the ideal one.

Thus T samples from P are worth ESS(Q, T) samples from Q.

Therefore, the variance of the weights must be as small as possible.

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#### Likelihood Weighting: Proposal Distribution

$$Q(X \setminus E) = \prod_{X_i \in X \setminus E} P(X_i \mid pa_i, e)$$

Example :

Given a Bayesian network :  $P(X_1, X_2, X_3) = P(X_1) \times P(X_2 | X_1) \times P(X_3 | X_1, X_2)$  and Evidence  $X_2 = x_2$ .  $Q(X_1, X_3) = P(X_1) \times P(X_3 | X_1, X_2 = x_2)$ 

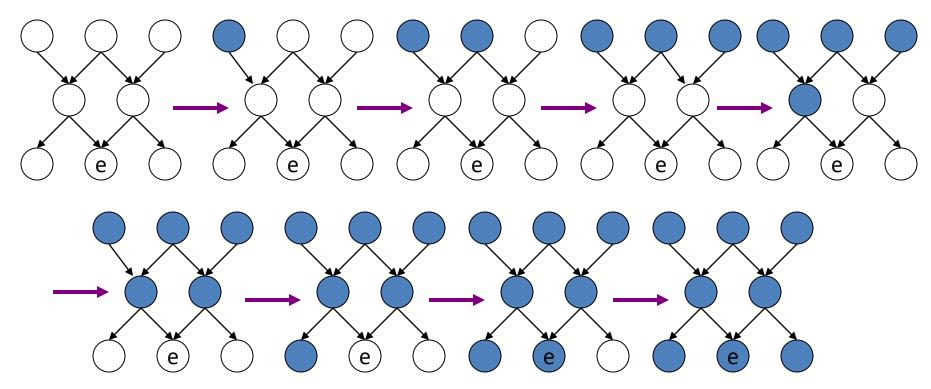
#### Weights:

Given a sample :  $x = (x_1, ..., x_n)$ 

$$w = \frac{P(x,e)}{Q(x)} = \frac{\prod_{X_i \in X \setminus E} P(x_i \mid pa_i, e) \times \prod_{E_j \in E} P(e_j \mid pa_j)}{\prod_{X_i \in X \setminus E} P(x_i \mid pa_i, e)}$$
$$= \prod_{E_j \in E} P(e_j \mid pa_j)$$

# Likelihood Weighting: Sampling

Sample in topological order over X !



Clamp evidence, Sample  $x_i \leftarrow P(X_i | pa_i)$ ,  $P(X_i | pa_i)$  is a look-up in CPT!

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#### **Proposal selection**

 One should try to select a proposal that is as close as possible to the posterior distribution.

$$Var_{Q}\left[\hat{P}(e)\right] = \frac{Var_{Q}[w(z)]}{T} = \frac{1}{N} \sum_{z \in Z} \left(\frac{P(z,e)}{Q(z)} - P(e)\right)^{2} Q(z)$$

 $\frac{P(z,e)}{Q(z)} - P(e) = 0$ , to have a zero - variance estimator

$$\therefore \frac{P(z,e)}{P(e)} = Q(z)$$

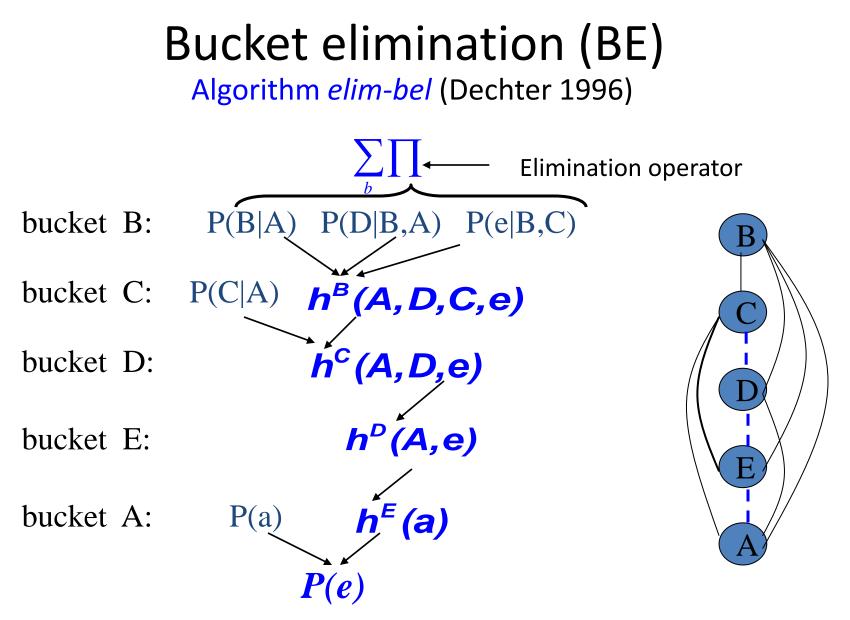
 $\therefore Q(z) = P(z \mid e)$ 

# Proposal Distributions used in Literature

- AIS-BN (Adaptive proposal)
  - Cheng and Druzdzel, 2000
- Iterative Belief Propagation
  - Changhe and Druzdzel, 2003
- Iterative Join Graph Propagation (IJGP) and variable ordering
  - Gogate and Dechter, 2005

# Perfect sampling using Bucket Elimination

- Algorithm:
  - Run Bucket elimination on the problem along an ordering  $o=(X_N,..,X_1)$ .
  - Sample along the reverse ordering:  $(X_1, ..., X_N)$
  - At each variable  $X_i$ , recover the probability  $P(X_i | x_1, ..., x_{i-1})$  by referring to the bucket.



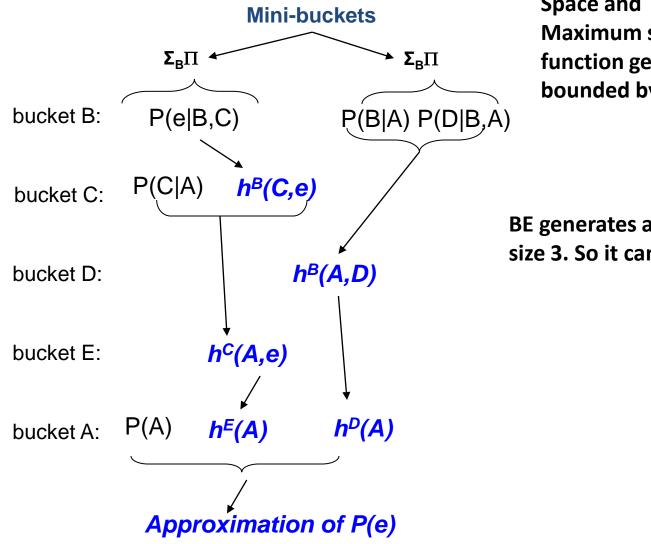
#### Sampling from the output of BE (Dechter 2002)

Set A = a, D = d, C = c in the bucket Sample :  $B = b \leftarrow Q(C \mid a, e, d) \propto P(B \mid a)P(d \mid B, a)P(e \mid b, c)$ bucket B: P(B|A) P(D|B,A) P(e|B,C)bucket C: P(C|A)  $h^{B}(A, D, C, e) \xrightarrow{Set A = a, D = d in the bucket}$ Sample :  $C = c \leftarrow Q(C | a, e, d) \propto h^{B}(a, d, C, e)$ Set A = a in the bucket h<sup>c</sup>(A,D,e) bucket D: Sample :  $D = d \leftarrow Q(D | a, e) \propto h^{C}(a, D, e)$ bucket E: **h<sup>D</sup>(A,e)** Evidence bucket : ignore  $P(A) \quad \boldsymbol{h}^{\boldsymbol{E}}(\boldsymbol{A})$ bucket A:  $\mathbf{Q}(\mathbf{A}) \propto \mathbf{P}(\mathbf{A}) \times \mathbf{h}^{\mathrm{E}}(\mathbf{A})$ Sample :  $A = a \leftarrow Q(A)$ 

#### Mini-buckets: "local inference"

- Computation in a bucket is time and space exponential in the number of variables involved
- Therefore, partition functions in a bucket into "mini-buckets" on smaller number of variables
- Can control the size of each "mini-bucket", yielding polynomial complexity.

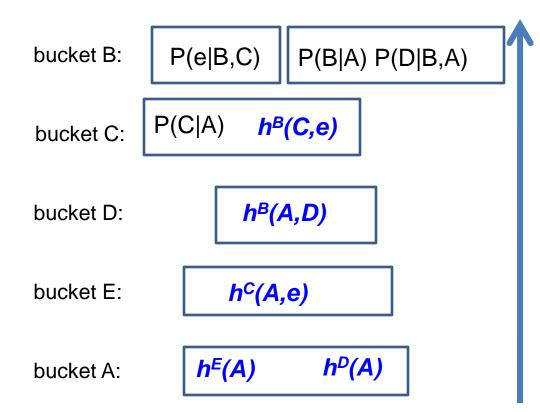
# **Mini-Bucket Elimination**



Space and Time constraints: Maximum scope size of the new function generated should be bounded by 2

BE generates a function having scope size 3. So it cannot be used.

# Sampling from the output of MBE



Sampling is same as in BE-sampling except that now we construct Q from a randomly selected "minibucket"

# IJGP-Sampling (Gogate and Dechter, 2005)

• Iterative Join Graph Propagation (IJGP)

A Generalized Belief Propagation scheme (Yedidia et al., 2002)

 IJGP yields better approximations of P(X|E) than MBE

- (Dechter, Kask and Mateescu, 2002)

- Output of IJGP is same as mini-bucket "clusters"
- Currently the best performing IS scheme!

### Adaptive Importance Sampling

Initial Proposal =  $Q^1(Z) = Q(Z_1) \times Q(Z_2 \mid pa(Z_2)) \times ... \times Q(Z_n \mid pa(Z_n))$  $\hat{P}(E = e) = 0$ 

For i = 1 to k do

Generate samples  $z^1, ..., z^N$  from  $Q^k$ 

$$\hat{P}(E = e) = \hat{P}(E = e) + \frac{1}{N} \sum_{j=1}^{N} w_k(z^i)$$
Update  $Q^{k+1} = Q^k + \eta(k) [Q^k - Q']$ 
End

Return 
$$\frac{\hat{P}(E=e)}{k}$$

# Adaptive Importance Sampling

- General case
- Given k proposal distributions
- Take N samples out of each distribution
- Approximate P(e)

$$\hat{P}(e) = \frac{1}{k} \sum_{j=1}^{k} \left[ Avg - weight - jth - proposal \right]$$

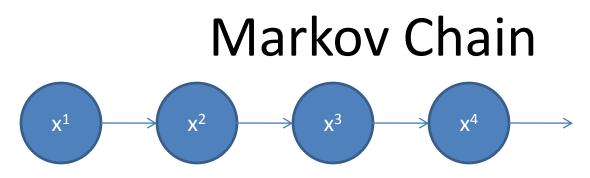
# Estimating Q'(z)

 $Q'(Z) = Q'(Z_1) \times Q'(Z_2 | pa(Z_2)) \times ... \times Q'(Z_n | pa(Z_n))$ where each  $Q'(Z_i | Z_1,..., Z_{i-1})$ 

is estimated by importance sampling

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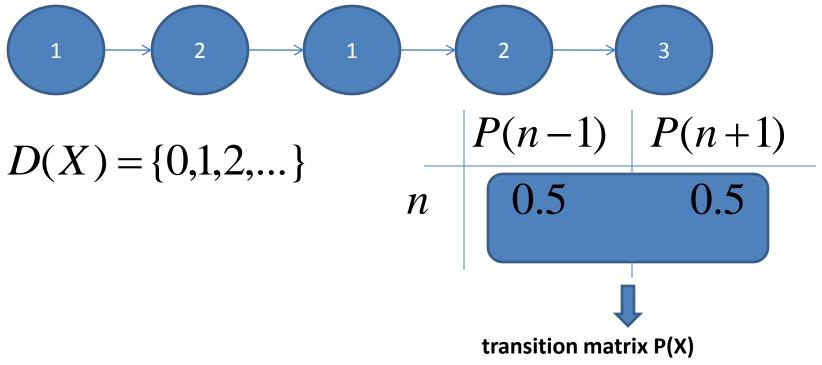
 A Markov chain is a discrete random process with the property that the next state depends only on the current state (Markov Property):

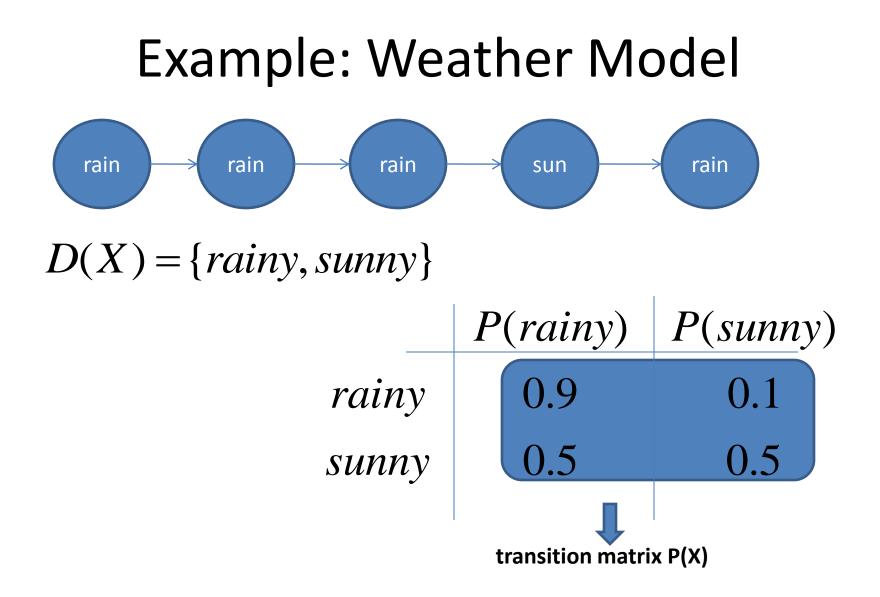
$$P(x^{t} | x^{1}, x^{2}, ..., x^{t-1}) = P(x^{t} | x^{t-1})$$

• If  $P(X^t|x^{t-1})$  does not depend on t (time homogeneous) and state space is finite, then it is often expressed as a transition function (aka transition matrix)  $\sum P(X = x) = 1$ 

# Example: Drunkard's Walk

 a random walk on the number line where, at each step, the position may change by +1 or -1 with equal probability

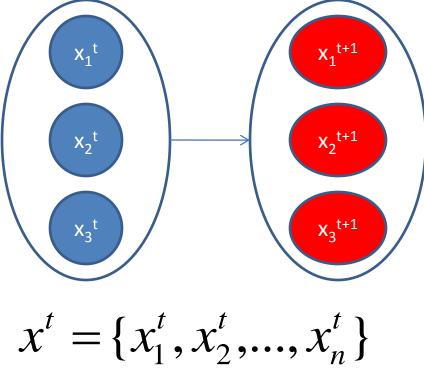




#### Multi-Variable System

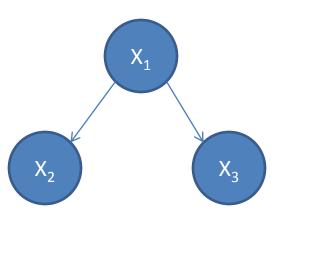
 $X = \{X_1, X_2, X_3\}, D(X_i) = discrete, finite$ 

state is an assignment of values to all the variables

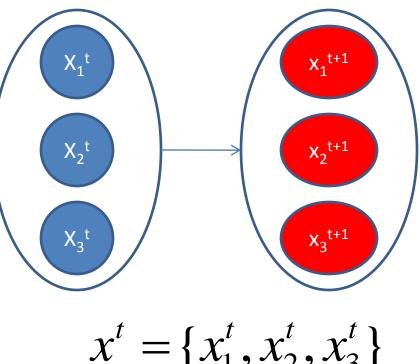


# **Bayesian Network System**

 Bayesian Network is a representation of the joint probability distribution over 2 or more variables



 $X = \{X_1, X_2, X_3\}$ 



# Stationary Distribution Existence

 If the Markov chain is time-homogeneous, then the vector π(X) is a *stationary* distribution (aka *invariant* or *equilibrium* distribution, aka "fixed point"), if its entries sum up to 1 and satisfy:

$$\pi(x_i) = \sum_{x_i \in D(X)} \pi(x_j) P(x_i \mid x_j)$$

- Finite state space Markov chain has a unique stationary distribution if and only if:
  - The chain is irreducible
  - All of its states are positive recurrent

# Irreducible

- A state x is *irreducible* if under the transition rule one has nonzero probability of moving from x to any other state and then coming back in a finite number of steps
- If one state is irreducible, then all the states must be irreducible

(Liu, Ch. 12, pp. 249, Def. 12.1.1)

#### Recurrent

- A state  $\chi$  is *recurrent* if the chain returns to  $\chi$  with probability 1
- Let  $M(\chi)$  be the expected number of steps to return to state  $\chi$
- State  $\chi$  is *positive recurrent* if M( $\chi$ ) is finite The recurrent states in a finite state chain are positive recurrent.

#### **Stationary Distribution Convergence**

• Consider infinite Markov chain:

 $n \rightarrow \infty$ 

$$P^{(n)} = P(x^n | x^0) = P^0 P^n$$

- If the chain is both *irreducible* and *aperiodic*, then:  $\pi = \lim P^{(n)}$
- Initial state is not important in the limit *"The most useful feature of a "good" Markov chain is its fast forgetfulness of its past..."* (Liu, Ch. 12.1)

# Aperiodic

- Define d(i) = g.c.d.{n > 0 | it is possible to go from i to i in n steps}. Here, g.c.d. means the greatest common divisor of the integers in the set. If d(i)=1 for ∀i, then chain is aperiodic
- Positive recurrent, aperiodic states are ergodic

## Markov Chain Monte Carlo

- How do we estimate P(X), e.g., P(X/e) ?
- Generate samples that form Markov Chain with stationary distribution π=P(X/e)
- Estimate π from samples (observed states): visited states x<sup>0</sup>,...,x<sup>n</sup> can be viewed as "samples" from distribution π

$$\overline{\pi}(x) = \frac{1}{T} \sum_{t=1}^{T} \delta(x, x^{t})$$
$$\pi = \lim_{T \to \infty} \overline{\pi}(x)$$

# **MCMC Summary**

- Convergence is guaranteed in the limit
- Initial state is not important, but... typically, we throw away first K samples - "burn-in"
- Samples are dependent, not i.i.d.
- Convergence (*mixing rate*) may be slow
- The stronger correlation between states, the slower convergence!

# Gibbs Sampling (Geman&Geman, 1984)

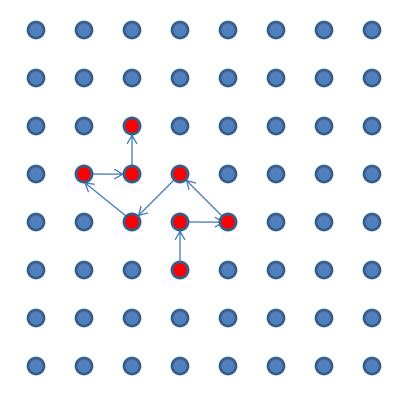
- Gibbs sampler is an algorithm to generate a sequence of samples from the joint probability distribution of two or more random variables
- Sample new variable value one variable at a time from the variable's conditional distribution:

 $P(X_i) = P(X_i \mid x_1^t, ..., x_{i-1}^t, x_{i+1}^t, ..., x_n^t) = P(X_i \mid x^t \setminus x_i)$ 

Samples form a Markov chain with stationary distribution *P(X/e)*

# **Gibbs Sampling: Illustration**

The process of Gibbs sampling can be understood as a *random walk* in the space of all instantiations of X=x (remember drunkard's walk):



In one step we can reach instantiations that differ from current one by value assignment to at most one variable (assume randomized choice of variables X<sub>i</sub>).

#### **Ordered Gibbs Sampler**

Generate sample x<sup>t+1</sup> from x<sup>t</sup> :

Process All Variables In Some Order

In

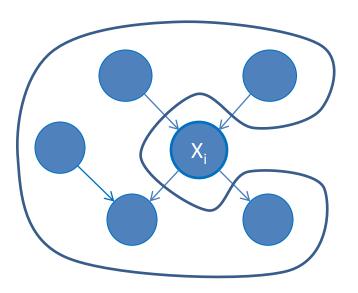
$$X_{1} = x_{1}^{t+1} \leftarrow P(X_{1} | x_{2}^{t}, x_{3}^{t}, ..., x_{N}^{t}, e)$$

$$X_{2} = x_{2}^{t+1} \leftarrow P(X_{2} | x_{1}^{t+1}, x_{3}^{t}, ..., x_{N}^{t}, e)$$

$$...$$

$$X_{N} = x_{N}^{t+1} \leftarrow P(X_{N} | x_{1}^{t+1}, x_{2}^{t+1}, ..., x_{N-1}^{t+1}, e)$$
short, for i=1 to N:
$$X_{i} = x_{i}^{t+1} \leftarrow \text{sampled from } P(X_{i} | x^{t} \setminus x_{i}, e)$$

# **Transition Probabilities in BN**



Given *Markov blanket* (parents, children, and their parents), X<sub>i</sub> is independent of all other nodes

Markov blanket:  $markov(X_i) = pa_i \bigcup ch_i \bigcup (\bigcup_{X_j \in ch_j} pa_j)$ 

$$P(X_i \mid x^t \setminus x_i) = P(X_i \mid markov_i^t):$$
$$P(x_i \mid x^t \setminus x_i) \propto P(x_i \mid pa_i) \prod_{X_j \in ch_i} P(x_j \mid pa_j)$$

Computation is linear in the size of Markov blanket!

# Ordered Gibbs Sampling Algorithm (Pearl,1988)

Input: *X, E=e* 

Output: *T* samples {*x*<sup>*t*</sup> }

*Fix evidence E=e, initialize x<sup>0</sup> at random* 

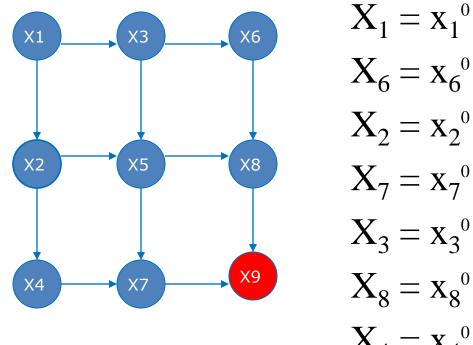
- 1. For t = 1 to T (compute samples)
- 2. For i = 1 to N (loop through variables)

3. 
$$\mathbf{x}_{i}^{t+1} \leftarrow P(X_{i} \mid markov_{i}^{t})$$

- 4. End For
- 5. End For

#### **Gibbs Sampling Example - BN**

$$X = \{X_1, X_2, ..., X_9\}, E = \{X_9\}$$



$$X_{6} - X_{6}$$

$$X_{2} = X_{2}^{0}$$

$$X_{7} = X_{7}^{0}$$

$$X_{3} = X_{3}^{0}$$

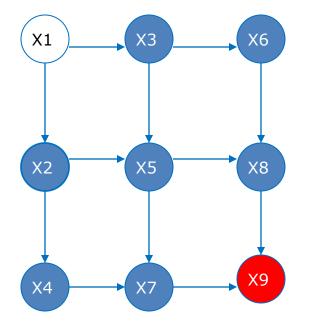
$$X_{8} = X_{8}^{0}$$

$$X_{4} = X_{4}^{0}$$

$$X_{5} = X_{5}^{0}$$

#### **Gibbs Sampling Example - BN**

$$X = \{X_1, X_2, ..., X_9\}, E = \{X_9\}$$



$$x_1^1 \leftarrow P(X_1 | x_2^0, ..., x_8^0, x_9)$$
  
 $x_2^1 \leftarrow P(X_2 | x_1^1, ..., x_8^0, x_9)$ 

# Answering Queries $P(x_i | e) = ?$

• **Method 1**: count # of samples where  $X_i = x_i$  (*histogram estimator*):

$$\overline{P}(X_i = x_i) = \frac{1}{T} \sum_{t=1}^{T} \delta(x_i, x^t)$$
 Dirac delta f-n

- Method 2: average probability (*mixture estimator*):  $\overline{P}(X_i = x_i) = \frac{1}{T} \sum_{t=1}^{T} P(X_i = x_i | markov_i^t)$
- Mixture estimator converges faster (consider estimates for the unobserved values of X<sub>i</sub>; prove via Rao-Blackwell theorem)

## Rao-Blackwell Theorem

**Rao-Blackwell Theorem:** Let random variable set X be composed of two groups of variables, R and L. Then, for the joint distribution π(R,L) and function g, the following result applies

 $Var[E\{g(R) | L\} \leq Var[g(R)]$ for a function of interest g, e.g., the mean or covariance (*Casella&Robert*, 1996, Liu et. al. 1995).

- theorem makes a weak promise, but works well in practice!
- improvement depends the choice of R and L

Importance vs. Gibbs Gibbs:  $x^t \leftarrow \hat{P}(X \mid e)$  $\hat{P}(X \mid e) \xrightarrow{T \to \infty} P(X \mid e)$  $\hat{g}(X) = \frac{1}{T} \sum_{i=1}^{I} g(x^{t})$ Importance:  $X^{t} \leftarrow Q(X \mid e)$ W<sub>t</sub>  $\overline{g} = \frac{1}{T} \sum_{t=1}^{T} \frac{g(x^t)P(x^t)}{Q(x^t)}$ 

## Gibbs Sampling: Convergence

- Sample from  $\overline{P}(X|e) \rightarrow P(X|e)$
- Converges iff chain is irreducible and ergodic
- Intuition must be able to explore all states:
  - if X<sub>i</sub> and X<sub>j</sub> are strongly correlated, X<sub>i</sub>=0↔ X<sub>j</sub>=0, then, we cannot explore states with X<sub>i</sub>=1 and X<sub>i</sub>=1
- All conditions are satisfied when all probabilities are positive
- Convergence rate can be characterized by the second eigen-value of transition matrix

# Gibbs: Speeding Convergence

Reduce dependence between samples (autocorrelation)

- Skip samples
- Randomize Variable Sampling Order
- Employ blocking (grouping)
- Multiple chains

Reduce variance (cover in the next section)

## **Blocking Gibbs Sampler**

- Sample several variables together, as a block
- Example: Given three variables X,Y,Z, with domains of size 2, group Y and Z together to form a variable W={Y,Z} with domain size 4. Then, given sample (x<sup>t</sup>, y<sup>t</sup>, z<sup>t</sup>), compute next sample:

$$x^{t+1} \leftarrow P(X \mid y^t, z^t) = P(w^t)$$

$$(y^{t+1}, z^{t+1}) = w^{t+1} \leftarrow P(Y, Z \mid x^{t+1})$$

- + Can improve convergence greatly when two variables are strongly correlated!
- Domain of the block variable grows exponentially with the #variables in a block!

## Gibbs: Multiple Chains

- Generate M chains of size K
- Each chain produces independent estimate  $P_m$ :

$$\overline{P}_m(x_i \mid e) = \frac{1}{K} \sum_{t=1}^K P(x_i \mid x^t \setminus x_i)$$

• Estimate  $P(x_i/e)$  as average of  $P_m(x_i/e)$ :  $\hat{P}(\bullet) = \frac{1}{M} \sum_{i=1}^{M} P_m(\bullet)$ 

Treat *P<sub>m</sub>* as independent random variables.

## **Gibbs Sampling Summary**

• Markov Chain Monte Carlo method

(Gelfand and Smith, 1990, Smith and Roberts, 1993, Tierney, 1994)

- Samples are **dependent**, form Markov Chain
- Sample from  $\overline{P}(X \mid e)$  which **converges** to  $\overline{P}(X \mid e)$
- Guaranteed to converge when all *P* > 0
- Methods to improve convergence:
  - Blocking
  - Rao-Blackwellised

# Overview

- 1. Probabilistic Reasoning/Graphical models
- 2. Importance Sampling
- 3. Markov Chain Monte Carlo: Gibbs Sampling
- 4. Sampling in presence of Determinism
- 5. Rao-Blackwellisation
- 6. AND/OR importance sampling

# Outline

- Rejection problem
- Backtrack-free distribution
  - Constructing it in practice
- SampleSearch
  - Construct the backtrack-free distribution on the fly.
- Approximate estimators
- Experiments

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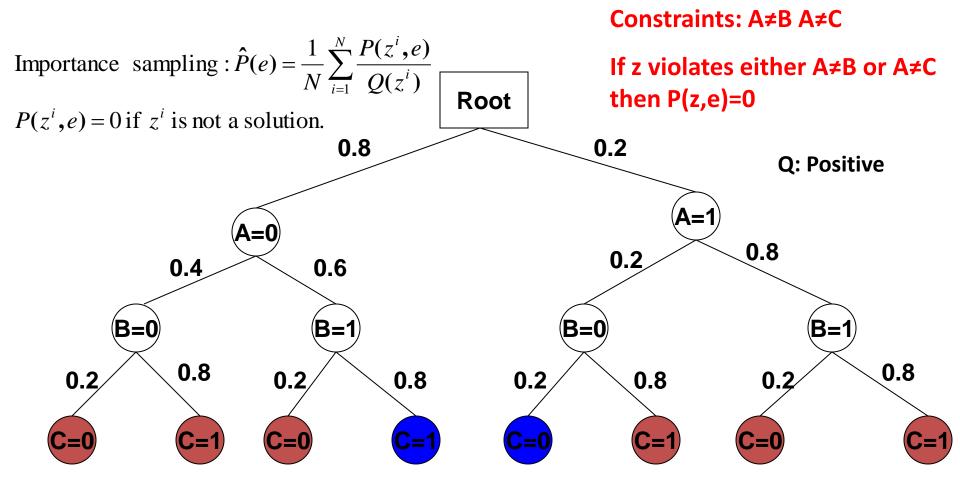
## **Rejection problem**

$$\hat{P}(e) = \frac{1}{N} \sum_{i=1}^{N} \frac{P(z^i, e)}{Q(z^i)}$$

- Importance sampling requirement  $-P(z,e) > 0 \rightarrow Q(z) > 0$
- When P(z,e)=0 but Q(z) > 0, the weight of the sample is zero and it is rejected.
- The probability of generating a rejected sample should be very small.

- Otherwise the estimate will be zero.

## **Rejection Problem**



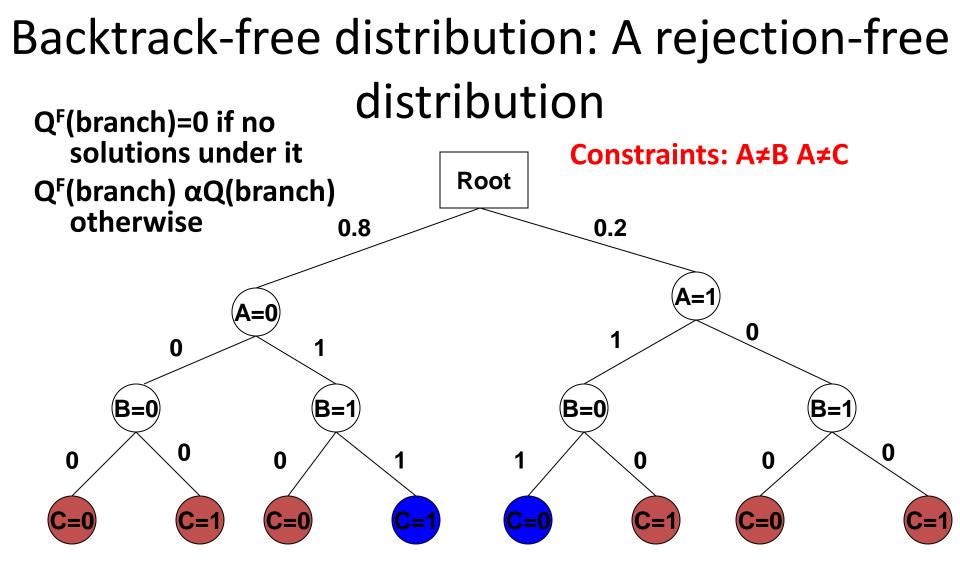
All Blue leaves correspond to solutions i.e. g(x) >0 All Red leaves correspond to non-solutions i.e. g(x)=0

# Outline

- Rejection problem
- Backtrack-free distribution

Constructing it in practice

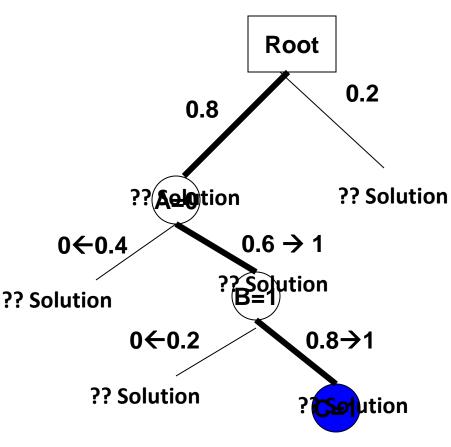
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All Blue leaves correspond to solutions i.e. g(x) >0 All Red leaves correspond to non-solutions i.e. g(x)=0

## Generating samples from Q<sup>F</sup>

**Constraints: A≠B A≠C** 

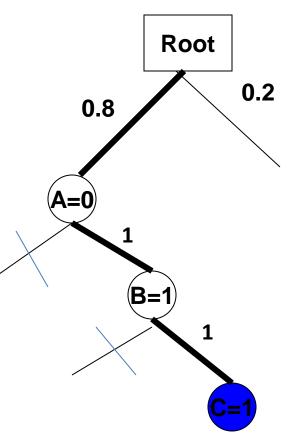


Q<sup>F</sup>(branch)=0 if no solutions under it Q<sup>F</sup>(branch) αQ(branch) otherwise

- Invoke an oracle at each branch.
  - Oracle returns True if there is a solution under a branch
  - False, otherwise

## Generating samples from Q<sup>F</sup>

**Constraints: A≠B A≠C** 

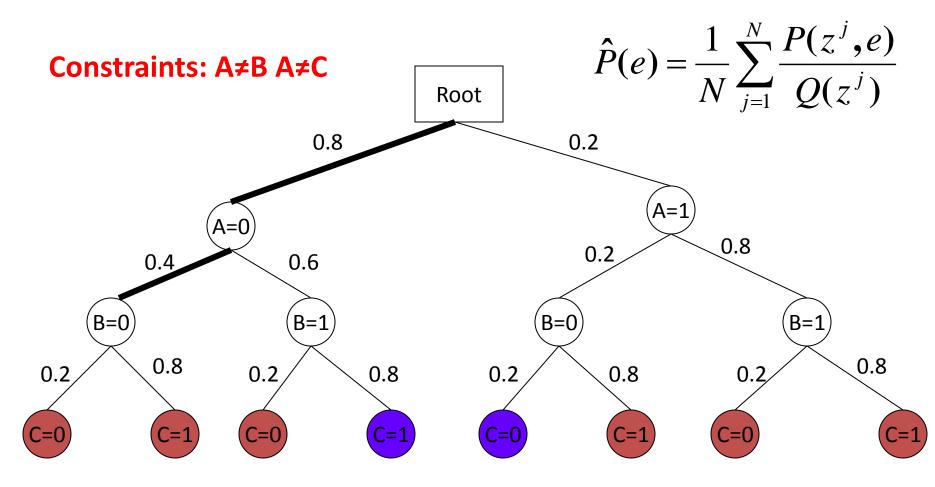


- Oracles
  - Adaptive consistency as preprocessing step
    - Constant time table lookup
    - Exponential in the treewidth of the constraint portion.
  - A complete CSP solver
    - Need to run it at each assignment.

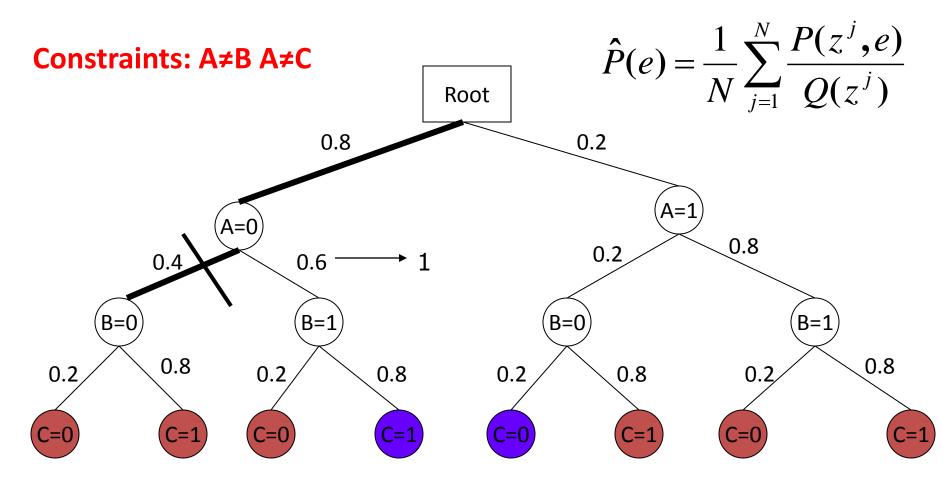
# Outline

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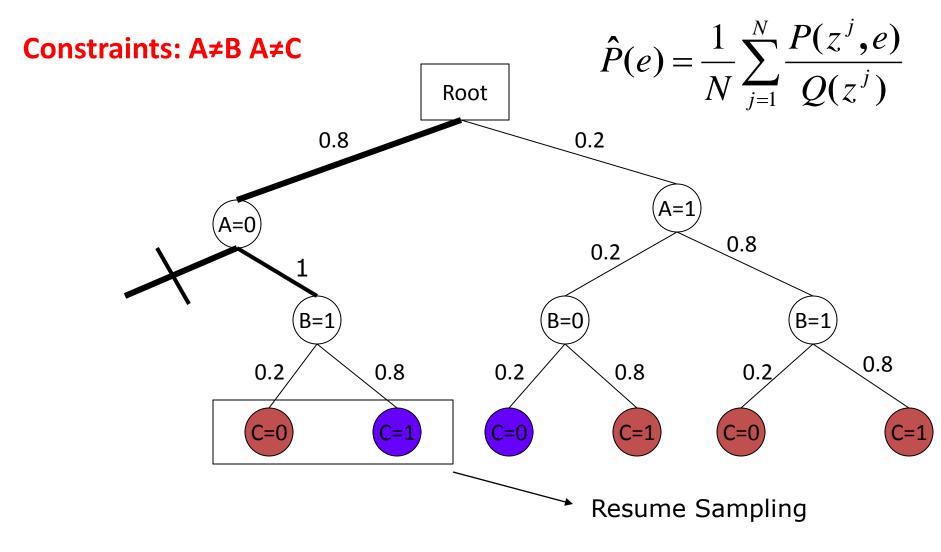
## Algorithm SampleSearch



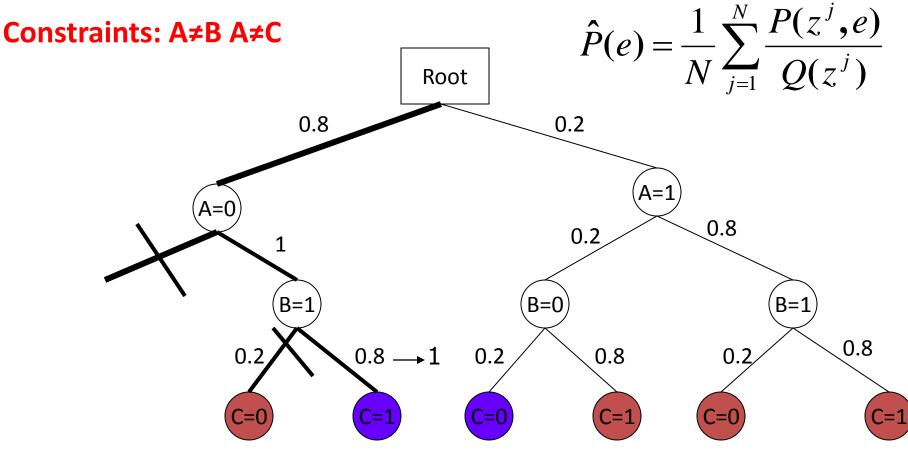
## Algorithm SampleSearch



## Algorithm SampleSearch



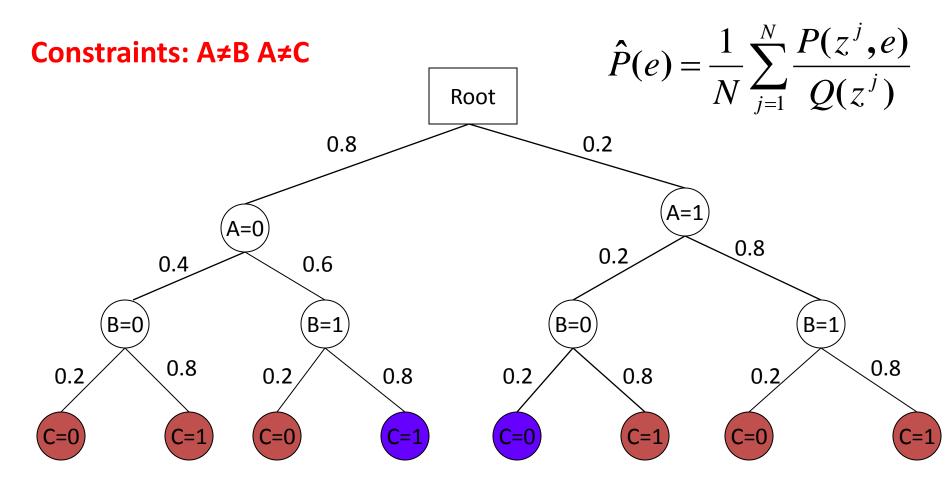
## Algorithm SampleSearch



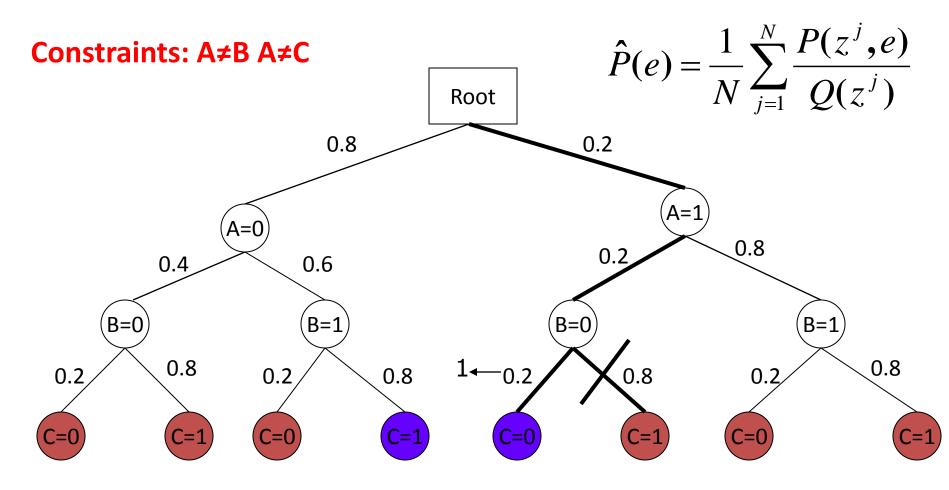
Until P(sample,e)>0

#### Constraint Violated

#### **Generate more Samples**

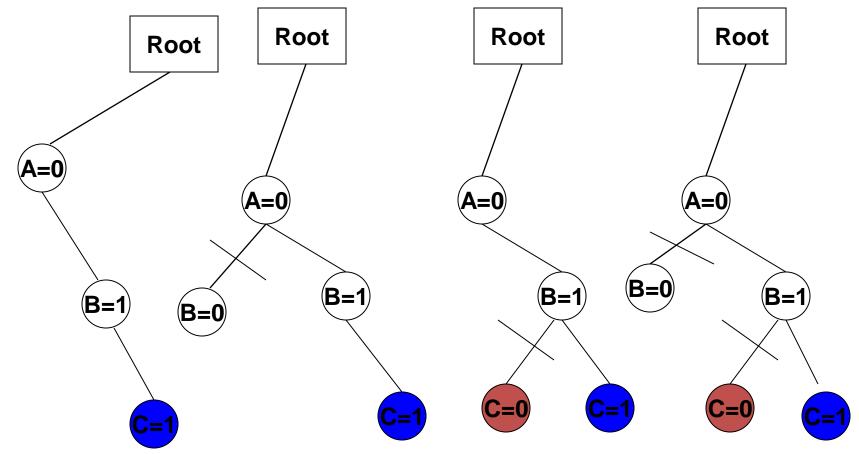


#### **Generate more Samples**



#### Traces of SampleSearch

#### **Constraints:** A≠B A≠C



#### SampleSearch: Sampling Distribution

• Problem: Due to Search, the samples are no longer i.i.d. from Q.

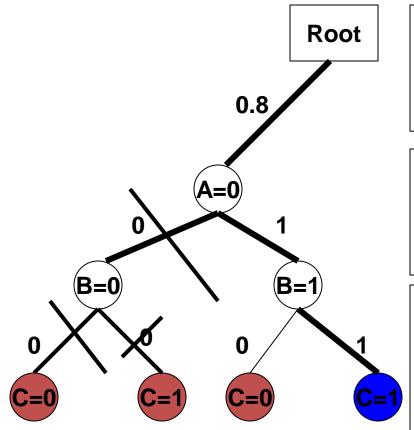
$$\overline{P}(e) = \frac{1}{N} \sum_{j=1}^{N} \frac{P(z^{j}, e)}{Q(z^{j})}, \quad E_{Q}[\overline{P}(e)] \neq P(e)$$

• Thm: SampleSearch generates i.i.d. samples from the **backtrack-free distribution** 

$$\hat{P}_{F}(e) = \frac{1}{N} \sum_{j=1}^{N} \frac{P(z^{j}, e)}{Q^{F}(z^{j})}, \quad E_{Q^{F}}[\hat{P}_{F}(e)] = P(e)$$

#### The Sampling distribution Q<sup>F</sup> of SampleSearch $\hat{P}(e) = \frac{1}{N} \sum_{j=1}^{N} \frac{P(z^{j}, e)}{O(z^{j})}$

**Constraints: A≠B A≠C** 



What is probability of generating A=0?  $Q^{F}(A=0)=0.8$ 

Why? SampleSearch is systematic

What is probability of generating (A=0,B=1)? Q<sup>F</sup>(B=1|A=0)=1

Why? SampleSearch is systematic

What is probability of generating (A=0,B=0)?

Simple: Q<sup>F</sup>(B=0|A=0)=0

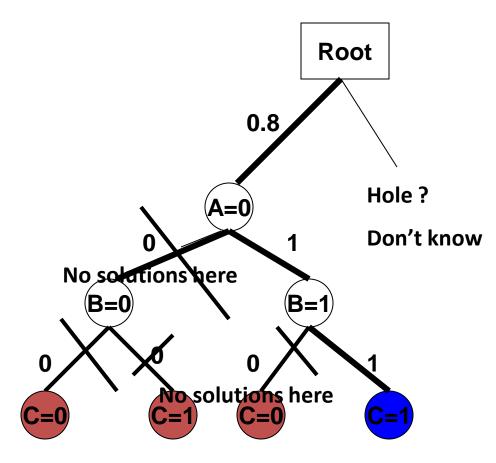
All samples generated by SampleSearch are solutions

**Backtrack-free distribution** 

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- Rejection problem
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## Asymptotic approximations of Q<sup>F</sup>



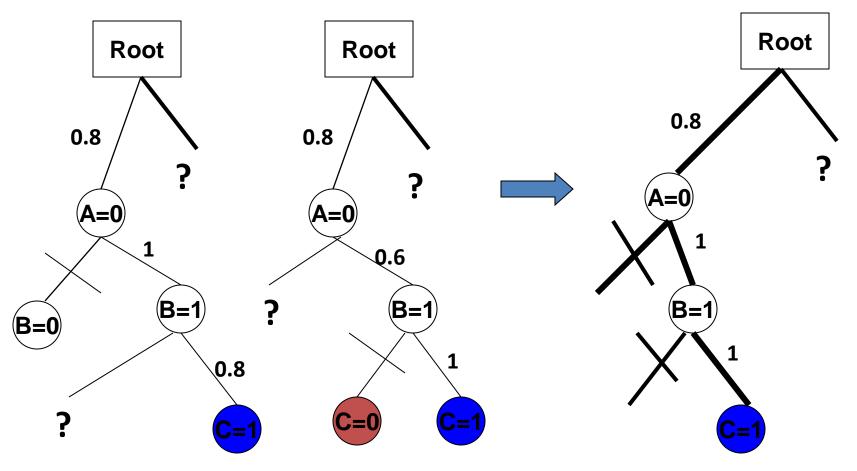
•IF Hole THEN

•U<sup>F</sup>=Q (i.e. there is a solution at the other branch)

•L<sup>F</sup>=0 (i.e. no solution at the other branch)

# Approximations: Convergence in the limit

• Store all possible traces

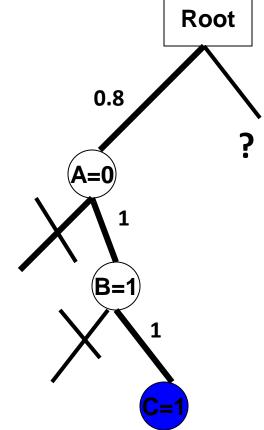


# Approximations: Convergence in the limit

• From the combined sample tree, update U and L. IF Hole THEN  $U_{N}^{F}=Q$  and  $L_{N}^{F}=0$ 

$$\lim_{N \to \infty} E\left[\frac{P(z,e)}{U_N^F(z)}\right] = \lim_{N \to \infty} E\left[\frac{P(z,e)}{L_N^F(z)}\right] = P(e)$$

Asymptotic ally unbiased Bounding  $: U_N^F(z) \le Q^F(z) \le L_N^F(z)$  $\overline{P}_E^U(e) \ge \hat{P}_E(e) \ge \overline{P}_E^L(e)$ 



## **Upper and Lower Approximations**

- Asymptotically unbiased.
- Upper and lower bound on the unbiased sample mean
- Linear time and space overhead
- Bias versus variance tradeoff
  - Bias = difference between the upper and lower approximation.

## Improving Naive SampleSeach

- Better Search Strategy
  - Can use any state-of-the-art CSP/SAT solver e.g. minisat (Een and Sorrenson 2006)
    - All theorems and result hold
- Better Importance Function
  - Use output of generalized belief propagation to compute the initial importance function Q (Gogate and Dechter, 2005)

### Experiments

- Tasks
  - Weighted Counting
  - Marginals
- Benchmarks
  - Satisfiability problems (counting solutions)
  - Linkage networks
  - Relational instances (First order probabilistic networks)
  - Grid networks
  - Logistics planning instances
- Algorithms
  - SampleSearch/UB, SampleSearch/LB
  - SampleCount (Gomes et al. 2007, SAT)
  - ApproxCount (Wei and Selman, 2007, SAT)
  - RELSAT (Bayardo and Peshoueshk, 2000, SAT)
  - Edge Deletion Belief Propagation (Choi and Darwiche, 2006)
  - Iterative Join Graph Propagation (Dechter et al., 2002)
  - Variable Elimination and Conditioning (VEC)
  - EPIS (Changhe and Druzdzel, 2006)

## Results: Solution Counts Langford instances

Problem	$\langle n,k,c,w angle$	Exact	Sample	Approx	REL	SS	SS
			Count	Count	SAT	/LB	/UB
lang12	$\langle 576, 2, 13584, 383 \rangle$	2.16E+5	1.93E+05	2.95E+04	2.16E+05	2.16E+05	2.16E+05
lang16	$\langle 1024, 2, 32320, 639 \rangle$	6.53E+08	5.97E+08	8.22E+06	6.28E+06	6.51E+08	6.99E+08
lang19	$\langle 1444, 2, 54226, 927 \rangle$	5.13E+11	9.73E+10	6.87E+08	8.52E+05	6.38E+11	7.31E+11
lang20	$\langle 1600, 2, 63280, 1023 \rangle$	5.27E+12	1.13E+11	3.99E+09	8.55E+04	2.83E+12	3.45E+12
lang23	$\langle 2116, 2, 96370, 1407 \rangle$	7.60E+15	7.53E+14	3.70E+12	Х	4.17E+15	4.19E+15
lang24	$\langle 2304, 2, 109536, 1535 \rangle$	9.37E+16	1.17E+13	4.15E+11	Х	8.74E+15	1.40E+16
lang27	$\langle 2916, 2, 156114, 1919 \rangle$		4.38E+16	1.32E+14	Х	2.41E+19	2.65E+19

#### Time Bound: 10 hrs

1e+17 1e+16 - - + - -2.5 I. 1e+15 1e+14 1e+13 1e+12 ..×....×....×....× x 1e+11 5000 10000 15000 20000 25000 30000 35000 40000 0 Time in seconds ApproxCount .....X····· SS/UB Exact .... .... SampleCount ---+---SS/LB ---

Number of Solutions

Solution Counts vs Time for lang24.cnf

## Results: Probability of Evidence Linkage instances (UAI 2006 evaluation)

Problem	$\langle n, k, e, w \rangle$	Exact	VEC	EDBP	SS/LB	SS/UB
BN_69	$\langle 777, 7, 78, 47 \rangle$	5.28E-054	1.93E-61	2.39E-57	3.00E-55	3.00E-55
BN_70	$\langle 2315, 5, 159, 87 \rangle$	2.00E-71	7.99E-82	6.00E-79	1.21E-73	1.21E-73
BN_71	$\langle 1740, 6, 202, 70 \rangle$	5.12E-111	7.05E-115	1.01E-114	1.28E-111	1.28E-111
BN_72	$\langle 2155, 6, 252, 86 \rangle$	4.21E-150	1.32E-153	9.21E-155	4.73E-150	4.73E-150
BN_73	$\langle 2140, 5, 216, 101 \rangle$	2.26E-113	6.00E-127	2.24E-118	2.00E-115	2.00E-115
BN_74	$\langle 749, 6, 66, 45 \rangle$	3.75E-45	3.30E-48	5.84E-48	2.13E-46	2.13E-46
BN_75	$\langle 1820, 5, 155, 92 \rangle$	5.88E-91	5.83E-97	3.10E-96	2.19E-91	2.19E-91
BN_76	$\langle 2155, 7, 169, 64 \rangle$	4.93E-110	1.00E-126	3.86E-114	1.95E-111	1.95E-111

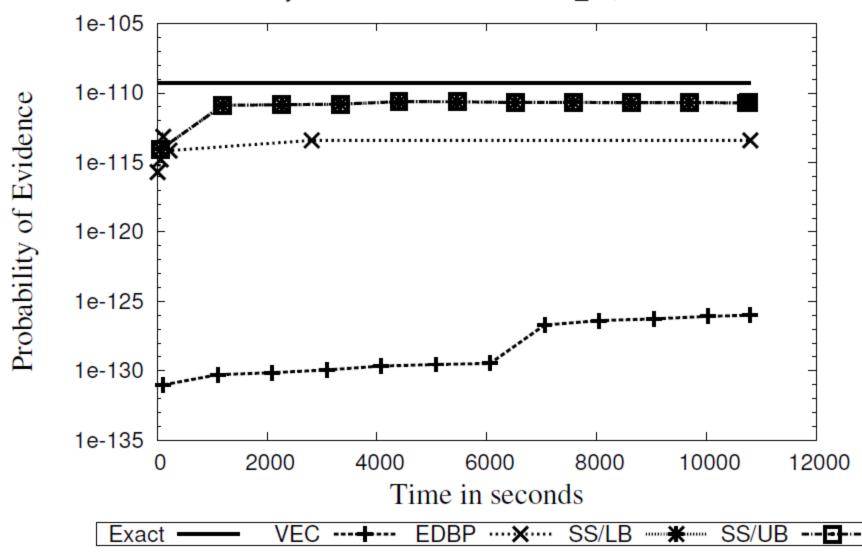
Time Bound: 3 hrs

## **Results: Probability of Evidence**

#### Linkage instances (UAI 2008 evaluation)

<b>U</b>		<b>\</b>				
Problem	$\langle n,k,e,w angle$	Exact	SS/LB	SS/UB	VEC	EDBP
pedigree18	$\langle 1184,1,0,26  angle$	7.18E-79	7.39E-79	7.39E-79	7.18E-79*	7.18E-79*
pedigree1	$\langle 334,2,0,20  angle$	7.81E-15	7.81E-15	7.81E-15	7.81E-15	7.81E-15*
pedigree20	$\langle 437, 2, 0, 25  angle$	2.34E-30	2.31E-30	2.31E-30	2.34E-30*	6.19E-31
pedigree23	$\langle 402, 1, 0, 26  angle$	2.78E-39	2.76E-39	2.76E-39	2.78E-39*	1.52E-39
pedigree25	$\langle 1289, 1, 0, 38 \rangle$	1.69E-116	1.69E-116	1.69E-116	1.69E-116*	1.69E-116*
pedigree30	$\langle 1289, 1, 0, 27 \rangle$	1.84E-84	1.90E-84	1.90E-84	1.85E-84*	1.85E-84*
pedigree37	$\langle 1032, 1, 0, 25 \rangle$	2.63E-117	1.18E-117	1.18E-117	2.63E-117*	5.69E-124
pedigree38	$\langle 724, 1, 0, 18 \rangle$	5.64E-55	3.80E-55	3.80E-55	5.65E-55*	8.41E-56
pedigree39	$\langle 1272, 1, 0, 29 \rangle$	6.32E-103	6.29E-103	6.29E-103	6.32E-103*	6.32E-103*
pedigree42	$\langle 448, 2, 0, 23  angle$	1.73E-31	1.73E-31	1.73E-31	1.73E-31*	8.91E-32
pedigree19	$\langle 793,2,0,23  angle$		6.76E-60	6.76E-60	1.597E-60	3.35E-60
pedigree31	$\langle 1183, 2, 0, 45 \rangle$		2.08E-70	2.08E-70	1.67E-76	1.34E-70
pedigree34	$\langle 1160, 1, 0, 59 \rangle$		3.84E-65	3.84E-65	2.58E-76	4.30E-65
pedigree13	$\langle 1077, 1, 0, 51 \rangle$		7.03E-32	7.03E-32	2.17E-37	6.53E-32
pedigree40	$\langle 1030,2,0,49  angle$		1.25E-88	1.25E-88	2.45E-91	7.02E-17
pedigree41	$\langle 1062, 2, 0, 52 \rangle$		4.36E-77	4.36E-77	4.33E-81	1.09E-10
pedigree44	$\langle 811,1,0,29 angle$		3.39E-64	3.39E-64	2.23E-64	7.69E-66
pedigree51	$\langle 1152, 1, 0, 51 \rangle$		2.47E-74	2.47E-74	5.56E-85	6.16E-76
pedigree7	$\langle 1068, 1, 0, 56 \rangle$		1.33E-65	1.33E-65	1.66E-72	2.93E-66
pedigree9	$\langle 1118, 2, 0, 41 \rangle$		2.93E-79	2.93E-79	8.00E-82	3.13E-89
1 0	X ///////		l I			

#### Time Bound: 3 hrs



#### Probability of Evidence vs Time for BN\_76, num-vars= 2155

## **Results on Marginals**

• Evaluation Criteria

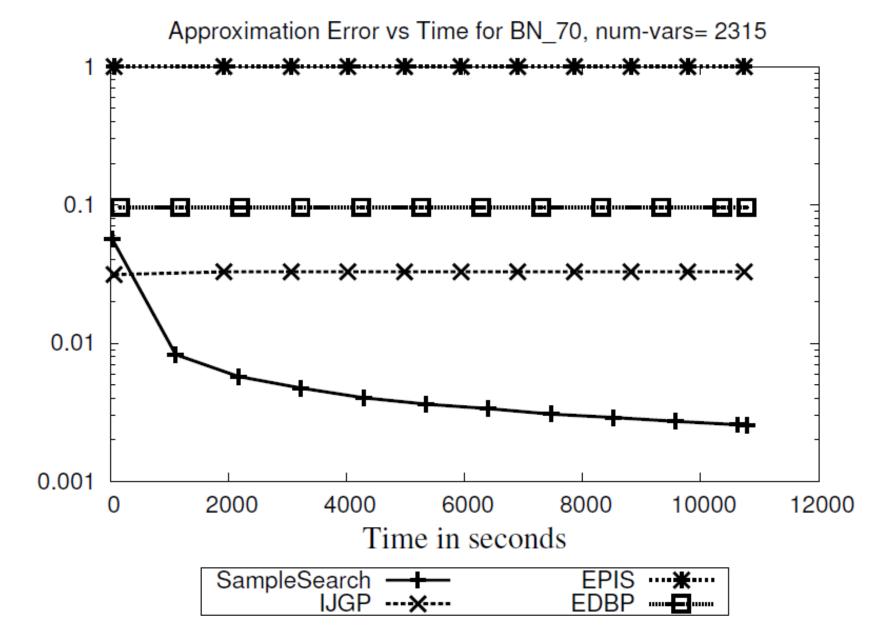
Exact:  $P(x_i)$  Approximate:  $A(x_i)$  $\frac{\sum_{i=1}^{n} \frac{1}{2} \sum_{x_i \in D_i} \left( \sqrt{P(x_i)} - \sqrt{A(x_i)} \right)^2}{n}$ Hellinger distance =  $\frac{n}{n}$ 

- Always bounded between 0 and 1
- Lower Bounds the KL distance
- When probabilities close to zero are present KL distance may tend to infinity.

## Results: Posterior Marginals Linkage instances (UAI 2006 evaluation)

Problem	$\langle n, K, e, w \rangle$	SampleSearch	IJGP	EPIS	EDBP
BN_69	$\langle 777, 7, 78, 47 \rangle$	9.4E-04	3.2E-02	1	8.0E-02
BN_70	$\langle 2315, 5, 159, 87 \rangle$	2.6E-03	3.3E-02	1	9.6E-02
BN_71	$\langle 1740, 6, 202, 70 \rangle$	5.6E-03	1.9E-02	1	2.5E-02
BN_72	$\langle 2155, 6, 252, 86 \rangle$	3.6E-03	7.2E-03	1	1.3E-02
BN_73	$\langle 2140, 5, 216, 101 \rangle$	2.1E-02	2.8E-02	1	6.1E-02
BN_74	$\langle 749, 6, 66, 45 \rangle$	6.9E-04	4.3E-06	1	4.3E-02
BN_75	$\langle 1820, 5, 155, 92 \rangle$	8.0E-03	6.2E-02	1	9.3E-02
BN_76	$\langle 2155, 7, 169, 64 \rangle$	1.8E-02	2.6E-02	1	2.7E-02

Time Bound: 3 hrs Distance measure: Hellinger distance



Average Hellinger Distance

# Summary: SampleSearch

- Manages rejection problem while sampling
  - Systematic backtracking search
- Sampling Distribution of SampleSearch is the backtrack-free distribution Q<sup>F</sup>
  - Expensive to compute
- Approximation of Q<sup>F</sup> based on storing all traces that yields an asymptotically unbiased estimator
  - Linear time and space overhead
  - Bound the sample mean from above and below
- Empirically, when a substantial number of zero probabilities are present, SampleSearch based schemes dominate their pure sampling counter-parts and Generalized Belief Propagation.

# Overview

- 1. Probabilistic Reasoning/Graphical models
- 2. Importance Sampling
- 3. Markov Chain Monte Carlo: Gibbs Sampling
- 4. Sampling in presence of Determinism
- 5. Rao-Blackwellisation
- 6. AND/OR importance sampling

# Sampling: Performance

• Gibbs sampling

Reduce dependence between samples

- Importance sampling
  - Reduce variance
- Achieve both by sampling a subset of variables and integrating out the rest (reduce dimensionality), aka Rao-Blackwellisation
- Exploit graph structure to manage the extra cost

## Smaller Subset State-Space

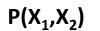
• Smaller state-space is easier to cover

$$X = \{X_1, X_2, X_3, X_4\} \qquad X = \{X_1, X_2\}$$
$$D(X) = 64 \qquad D(X) = 16$$

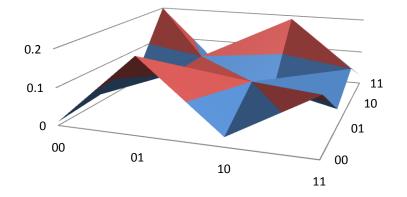
#### **Smoother Distribution**

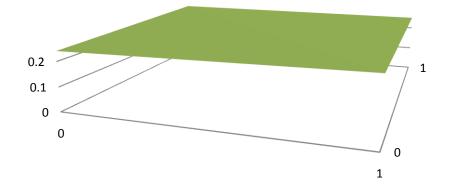
 $P(X_1, X_2, X_3, X_4)$ 

■ 0-0.1 ■ 0.1-0.2 ■ 0.2-0.26



0-0.1 0.1-0.2 0.2-0.26





## Speeding Up Convergence

• Mean Squared Error of the estimator:

$$MSE_{Q}\left[\overline{P}\right] = BIAS^{2} + Var_{Q}\left[\overline{P}\right]$$

- In case of unbiased estimator, BIAS=0  $MSE_{Q}[\hat{P}] = Var_{Q}[\hat{P}] = \left(E_{Q}[\hat{P}]^{2} - E_{Q}[P]^{2}\right)$
- Reduce variance  $\Rightarrow$  speed up convergence !

$$Rao-Blackwellisation$$

$$X = R \bigcup L$$

$$\hat{g}(x) = \frac{1}{T} \{h(x^{1}) + \dots + h(x^{T})\}$$

$$\tilde{g}(x) = \frac{1}{T} \{E[h(x) | l^{1}] + \dots + E[h(x) | l^{T}]\}$$

$$Var\{g(x)\} = Var\{E[g(x) | l]\} + E\{var[g(x) | l]\}$$

$$Var\{g(x)\} \ge Var\{E[g(x) | l]\}$$

$$Var\{\hat{g}(x)\} = \frac{Var\{h(x)\}}{T} \ge \frac{Var\{E[h(x) | l]\}}{T} = Var\{\tilde{g}(x)\}$$

Liu, Ch.2.3

## **Rao-Blackwellisation**

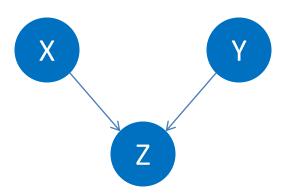
"Carry out analytical computation as much as possible" - Liu

- X=R∪L
- Importance Sampling:

$$Var_{Q}\left\{\frac{P(R,L)}{Q(R,L)}\right\} \ge Var_{Q}\left\{\frac{P(R)}{Q(R)}\right\}$$
Liu, Ch.2.5.5

- Gibbs Sampling:
  - autocovariances are lower (less correlation between samples)
  - if  $X_i$  and  $X_j$  are strongly correlated,  $X_i=0 \leftrightarrow X_j=0$ , only include one fo them into a sampling set

### Blocking Gibbs Sampler vs. Collapsed



Faster Convergence

- Standard Gibbs:  $P(x \mid y, z), P(y \mid x, z), P(z \mid x, y)$  (1)
  - Blocking:  $P(x \mid y, z), P(y, z \mid x)$ (2)
- Collapsed:
  - $P(x \mid y), P(y \mid x)$ (3)

#### Collapsed Gibbs Sampling Generating Samples

Generate sample c<sup>t+1</sup> from c<sup>t</sup> :

. . .

$$C_{1} = c_{1}^{t+1} \leftarrow P(c_{1} \mid c_{2}^{t}, c_{3}^{t}, ..., c_{K}^{t}, e)$$
$$C_{2} = c_{2}^{t+1} \leftarrow P(c_{2} \mid c_{1}^{t+1}, c_{3}^{t}, ..., c_{K}^{t}, e)$$

$$C_{K} = c_{K}^{t+1} \leftarrow P(c_{K} \mid c_{1}^{t+1}, c_{2}^{t+1}, ..., c_{K-1}^{t+1}, e)$$

In short, for i=1 to K:  $C_i = c_i^{t+1} \leftarrow \text{sampled from } P(c_i \mid c^t \setminus c_i, e)$ 

## **Collapsed Gibbs Sampler**

Input: *CX*, *E=e* Output: *T* samples {*c*<sup>*t*</sup> } *Fix evidence E=e, initialize c*<sup>0</sup> at random

- 1. For t = 1 to T (compute samples)
- 2. For i = 1 to N (loop through variables)

3. 
$$c_i^{t+1} \leftarrow P(C_i \mid c^t \setminus c_i)$$

- 4. End For
- 5. End For

# **Calculation Time**

- Computing P(c<sub>i</sub> / c<sup>t</sup>\c<sub>i</sub>, e) is more expensive (requires inference)
- Trading #samples for smaller variance:
  - generate more samples with higher covariance
  - generate fewer samples with lower covariance
- Must control the time spent computing sampling probabilities in order to be timeeffective!

# **Exploiting Graph Properties**

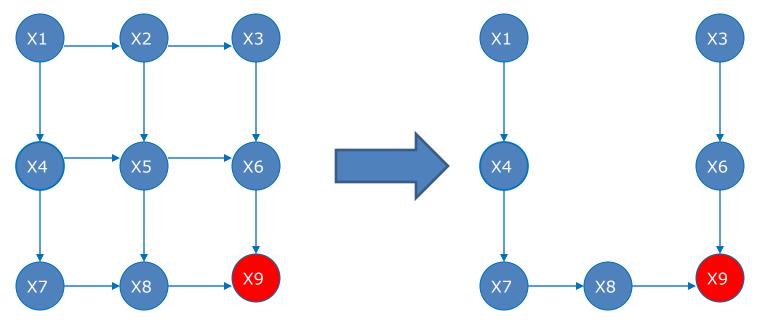
- Recall... computation time is *exponential in the adjusted induced width* of a graph
- *w*-cutset is a subset of variable s.t. when they are observed, induced width of the graph is *w*
- when sampled variables form a *w*-cutset , inference is exp(*w*) (e.g., using *Bucket Tree Elimination*)
- cycle-cutset is a special case of *w*-cutset

Sampling *w*-cutset  $\Rightarrow$  w-cutset sampling!

#### What If C=Cycle-Cutset ?

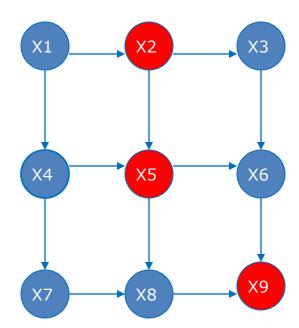
$$c^{0} = \{x_{2}^{0}, x_{5}^{0}\}, E = \{X_{9}\}$$

 $P(x_2, x_5, x_9)$  – can compute using Bucket Elimination



 $P(x_2, x_5, x_9)$  – computation complexity is O(N)

## **Computing Transition Probabilities**



Compute joint probabilities:

$$BE: P(x_2 = 0, x_3, x_9)$$
$$BE: P(x_2 = 1, x_3, x_9)$$

Normalize:

$$\alpha = P(x_2 = 0, x_3, x_9) + P(x_2 = 1, x_3, x_9)$$
$$P(x_2 = 0 | x_3) = \alpha P(x_2 = 0, x_3, x_9)$$
$$P(x_2 = 1 | x_3) = \alpha P(x_2 = 1, x_3, x_9)$$

#### **Cutset Sampling-Answering Queries**

• Query:  $\forall c_i \in C$ ,  $P(c_i | e) = ?$  same as Gibbs:

$$\hat{P}(c_i | e) = \frac{1}{T} \sum_{t=1}^{T} P(c_i | e^t \setminus c_i, e)$$
computed while generating sample t
using bucket tree elimination

• Query:  $\forall x_i \in X \setminus C, P(x_i | e) = ?$ 

$$\overline{P}(x_i/e) = \frac{1}{T} \sum_{t=1}^{T} P(x_i \mid c^t, e)$$

compute after generating sample t using bucket tree elimination

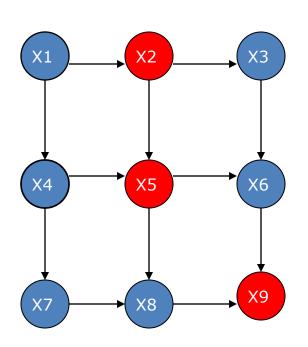
Cutset Sampling vs. Cutset Conditioning

- Cutset Conditioning  $P(x_i|e) = \sum_{c \in D(C)} P(x_i \mid c, e) \times P(c \mid e)$
- Cutset Sampling

$$\overline{P}(x_i/e) = \frac{1}{T} \sum_{t=1}^{T} P(x_i \mid c^t, e)$$

$$= \sum_{c \in D(C)} P(x_i \mid c, e) \times \frac{count(c)}{T}$$
$$= \sum_{c \in D(C)} P(x_i \mid c, e) \times \overline{P(c \mid e)}$$

Cutset Sampling Example Estimating  $P(x_2|e)$  for sampling node  $X_2$ :



$$x_{2}^{1} \leftarrow P(x_{2}/x_{5}^{0}, x_{9}) \quad \text{Sample 1}$$

$$\cdots$$

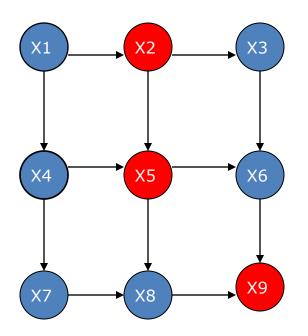
$$x_{2}^{2} \leftarrow P(x_{2}/x_{5}^{1}, x_{9}) \quad \text{Sample 2}$$

$$\cdots$$

$$x_{2}^{3} \leftarrow P(x_{2}/x_{5}^{2}, x_{9}) \quad \text{Sample 3}$$

$$\overline{P}(x_{2} \mid x_{9}) = \frac{1}{3} \begin{bmatrix} P(x_{2}/x_{5}^{0}, x_{9}) \\ + P(x_{2}/x_{5}^{1}, x_{9}) \\ + P(x_{2}/x_{5}^{2}, x_{9}) \end{bmatrix}$$

#### Cutset Sampling Example Estimating $P(x_3 | e)$ for non-sampled node $X_3$ :



$$c^{1} = \{x_{2}^{1}, x_{5}^{1}\} \Longrightarrow P(x_{3} \mid x_{2}^{1}, x_{5}^{1}, x_{9})$$

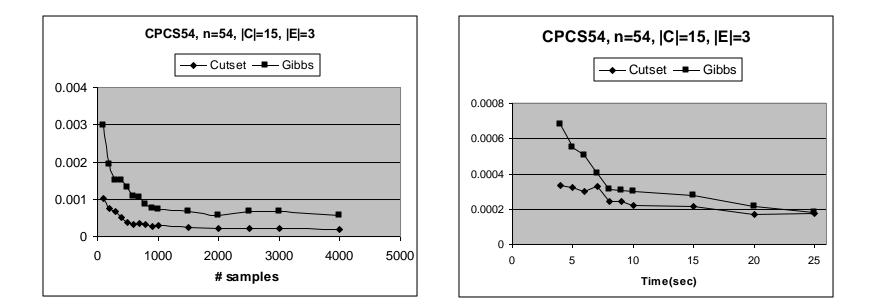
$$c^{2} = \{x_{2}^{2}, x_{5}^{2}\} \Longrightarrow P(x_{3} \mid x_{2}^{2}, x_{5}^{2}, x_{9})$$

$$c^{3} = \{x_{2}^{3}, x_{5}^{3}\} \Longrightarrow P(x_{3} \mid x_{2}^{3}, x_{5}^{3}, x_{9})$$

$$\Box P(x_{3} \mid x_{2}^{1}, x_{5}^{1}, x_{9})$$

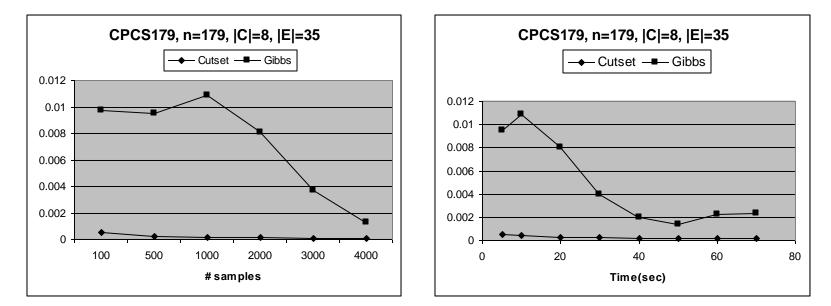
$$P(x_3 \mid x_9) = \frac{1}{3} \begin{bmatrix} P(x_3 \mid x_2^1, x_5^1, x_9) \\ + P(x_3 \mid x_2^2, x_5^2, x_9) \\ + P(x_3 \mid x_2^3, x_5^3, x_9) \end{bmatrix}$$

#### **CPCS54** Test Results



MSE vs. #samples (left) and time (right) Ergodic, |X|=54,  $D(X_i)=2$ , |C|=15, |E|=3Exact Time = 30 sec using Cutset Conditioning

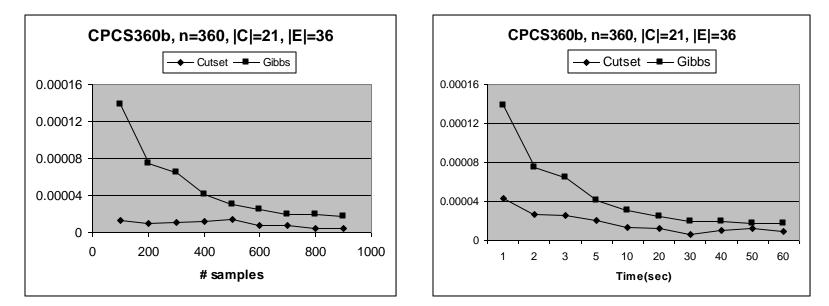
#### **CPCS179** Test Results



MSE vs. #samples (left) and time (right) Non-Ergodic (1 deterministic CPT entry) |X| = 179, |C| = 8,  $2 \le D(X_i) \le 4$ , |E| = 35

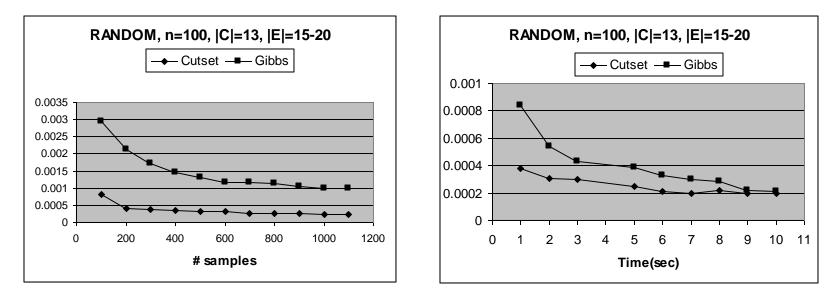
Exact Time = 122 sec using Cutset Conditioning

#### **CPCS360b** Test Results



MSE vs. #samples (left) and time (right) Ergodic, |X| = 360,  $D(X_i)=2$ , |C| = 21, |E| = 36Exact Time > 60 min using Cutset Conditioning Exact Values obtained via Bucket Elimination

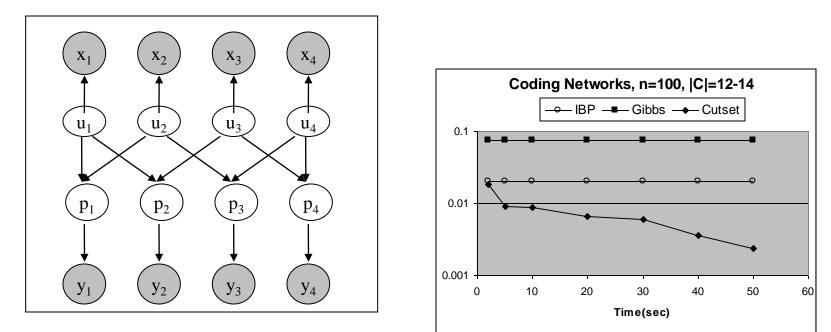
#### **Random Networks**



MSE vs. #samples (left) and time (right) |X| = 100,  $D(X_i) = 2$ , |C| = 13, |E| = 15-20Exact Time = 30 sec using Cutset Conditioning

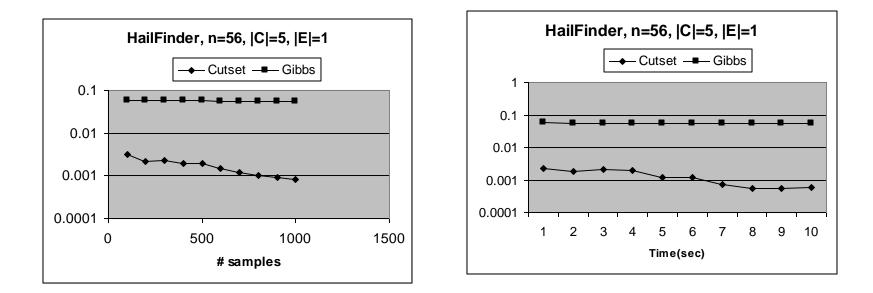
## **Coding Networks**

**Cutset Transforms Non-Ergodic Chain to Ergodic** 



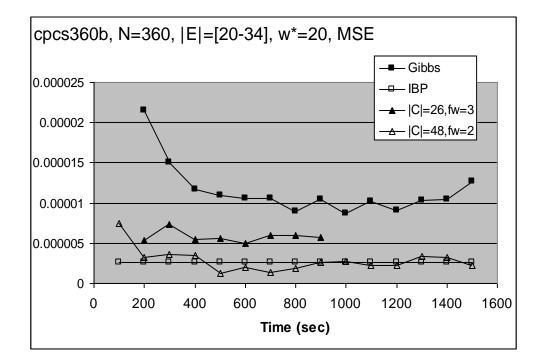
MSE vs. time (right) Non-Ergodic, |X| = 100,  $D(X_i)=2$ , |C| = 13-16, |E| = 50Sample Ergodic Subspace  $U=\{U_1, U_2, ..., U_k\}$ Exact Time = 50 sec using Cutset Conditioning

#### Non-Ergodic Hailfinder



MSE vs. #samples (left) and time (right) Non-Ergodic, |X| = 56, |C| = 5,  $2 <=D(X_i) <=11$ , |E| = 0Exact Time = 2 sec using Loop-Cutset Conditioning

### CPCS360b - MSE



MSE vs. Time

Ergodic, |X| = 360, |C| = 26,  $D(X_i)=2$ 

Exact Time = 50 min using BTE

### **Cutset Importance Sampling**

(Gogate & Dechter, 2005) and (Bidyuk & Dechter, 2006)

• Apply Importance Sampling over cutset C

$$\hat{P}(e) = \frac{1}{T} \sum_{t=1}^{T} \frac{P(c^{t}, e)}{Q(c^{t})} = \frac{1}{T} \sum_{t=1}^{T} w^{t}$$

where  $P(c^t, e)$  is computed using Bucket Elimination

$$\overline{P}(c_i \mid e) = \alpha \frac{1}{T} \sum_{t=1}^T \delta(c_i, c^t) w^t$$

$$\overline{P}(x_i \mid e) = \alpha \frac{1}{T} \sum_{t=1}^{T} P(x_i \mid c^t, e) w^t$$

# Likelihood Cutset Weighting (LCS)

- Z=Topological Order{C,E}
- Generating sample t+1:

For  $Z_i \in Z$  do: If  $Z_i \in E$   $z_i^{t+1} = z_i, z_i \in e$ Else  $z_i^{t+1} \leftarrow P(Z_i \mid z_1^{t+1}, \dots, z_{i-1}^{t+1})$ End If

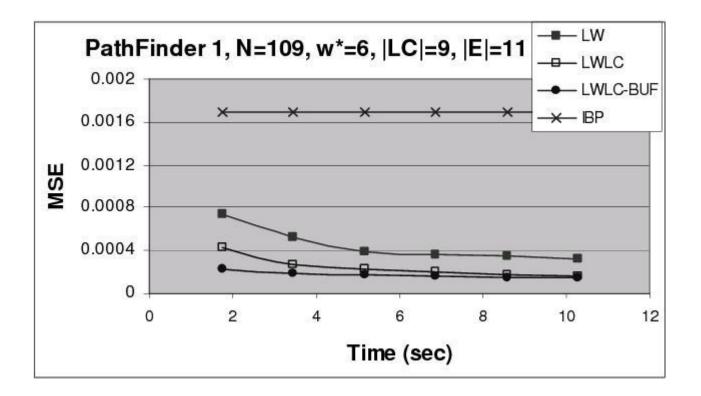
 computed while generating sample t using bucket tree elimination

 can be memoized for some number of instances K
 (based on memory available)

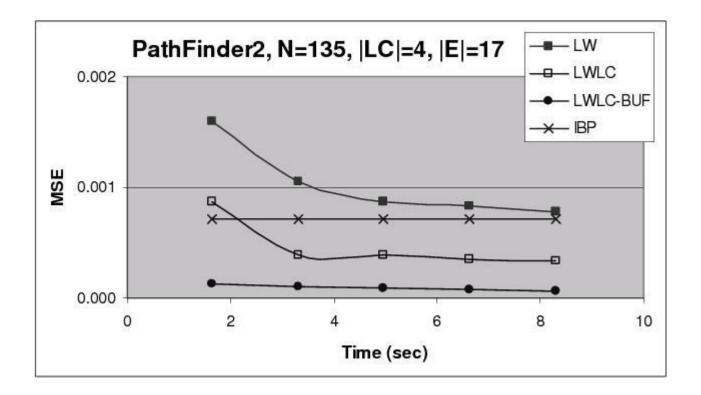
End For

 $KL[P(C|e), Q(C)] \le KL[P(X|e), Q(X)]$ 

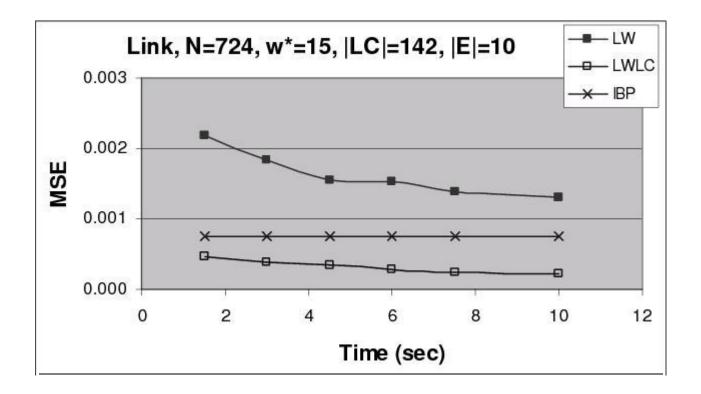
### Pathfinder 1



### Pathfinder 2



### Link



### Summary

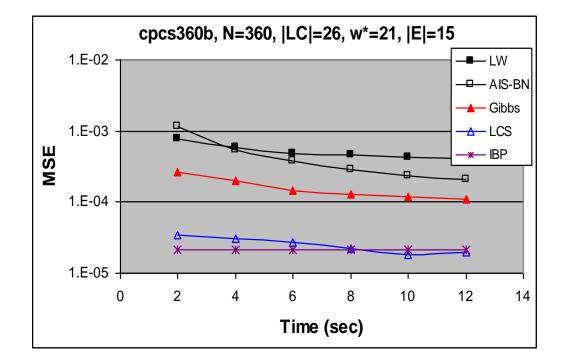
**Importance Sampling** 

- i.i.d. samples
- Unbiased estimator
- Generates samples fast
- Samples from Q
- Reject samples with zero-weight
- Improves on cutset

**Gibbs Sampling** 

- Dependent samples
- Biased estimator
- Generates samples slower
- Samples from  $\overline{P}(X|e)$
- Does not converge in presence of constraints
- Improves on cutset

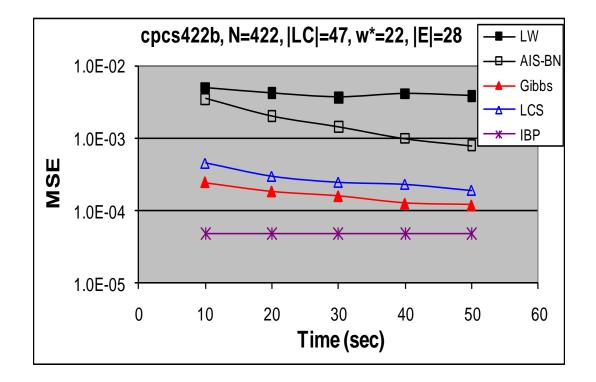
### CPCS360b



LW – likelihood weighting

LCS – likelihood weighting on a cutset

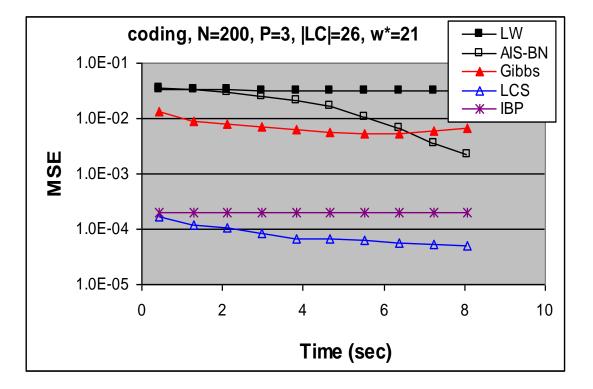
### CPCS422b



LW – likelihood weighting

LCS – likelihood weighting on a cutset

### **Coding Networks**



LW – likelihood weighting

LCS – likelihood weighting on a cutset

### Overview

- 1. Probabilistic Reasoning/Graphical models
- 2. Importance Sampling
- 3. Markov Chain Monte Carlo: Gibbs Sampling
- 4. Sampling in presence of Determinism
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### Motivation



Expected value of the number on the face of a die:  $\frac{1+2+3+4+5+6}{6} = 3.5$ 

What is the expected value of the product of the numbers on the face of "k" dice?

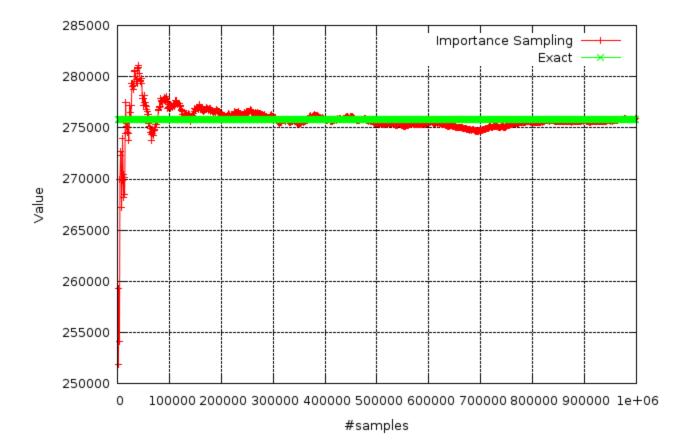
 $(3.5)^{k}$ 

### Monte Carlo estimate

- Perform the following experiment N times.
  - Toss all the k dice.
  - Record the product of the numbers on the top face of each die.
- Report the average over the N runs.

 $\hat{Z} = \frac{1}{N} \sum_{i=1}^{N} \prod_{j=1}^{k} \text{ (the number on the face of the "j<sup>th</sup>" dice in the N<sup>th</sup> run)}$ 

### How the sample average converges?



10 dice. Exact Answer is (3.5)<sup>10</sup>

### But This is Really Dumb?

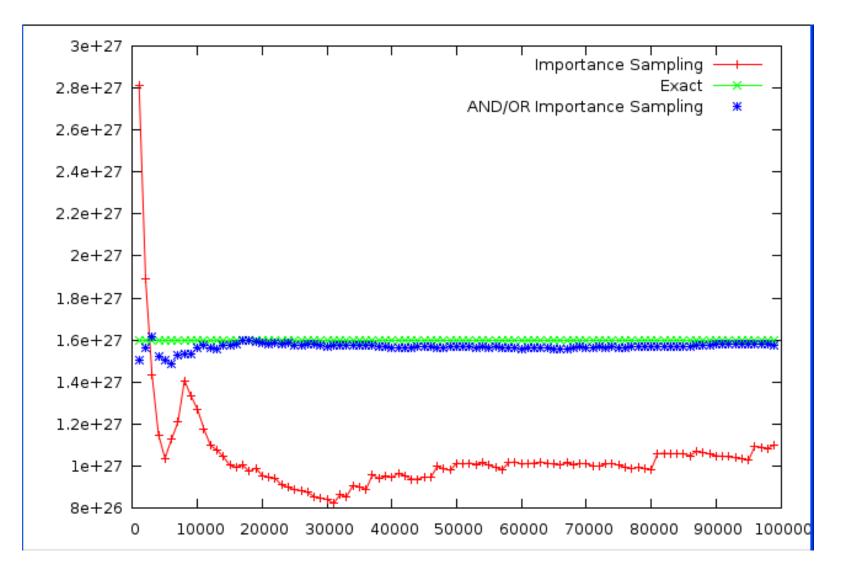
- The dice are independent.
- A better Monte Carlo estimate
  - 1. Perform the experiment N times
  - 2. For each dice record the average
  - 3. Take a product of the averages

 $\hat{Z}_{new} = \prod_{j=1}^{k} \frac{1}{N} \sum_{i=1}^{N} (\text{the number on the face of the "j<sup>th</sup>" dice in the N<sup>th</sup> run)}$ 

### • Conventional estimate: Averages of products.

 $\hat{Z} = \frac{1}{N} \sum_{i=1}^{N} \prod_{j=1}^{k} \text{ (the number on the face of the "j<sup>th</sup>" dice in the N<sup>th</sup> run)}$ 

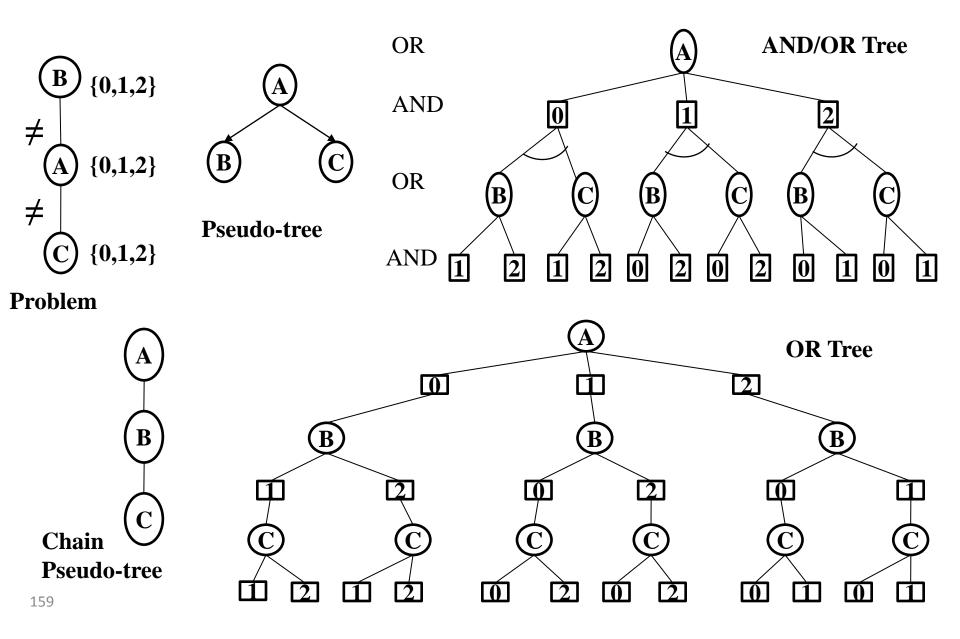
### How the sample Average Converges



# Moral of the story

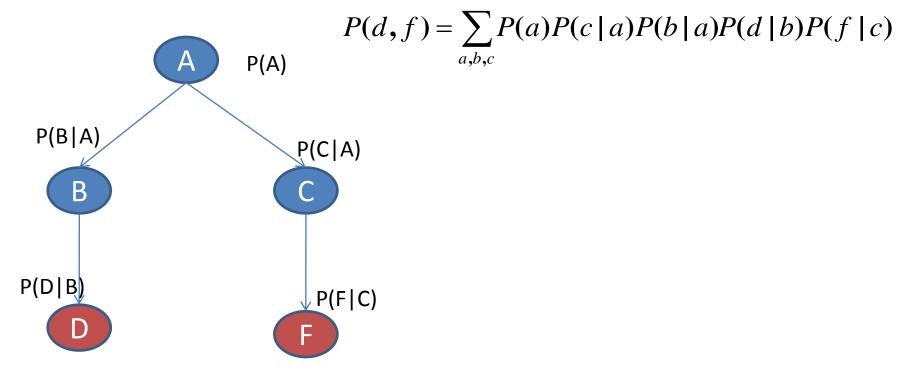
- Make use of (conditional) independence to get better results
- Used for exact inference extensively
  - Bucket Elimination (Dechter, 1996)
  - Junction tree (Lauritzen and Speigelhalter, 1988)
  - Value Elimination (Bacchus et al. 2004)
  - Recursive Conditioning (Darwiche, 2001)
  - BTD (Jegou et al., 2002)
  - AND/OR search (Dechter and Mateescu, 2007)
- How to use it for sampling?
   AND/OR Importance sampling

### Background: AND/OR search space



### AND/OR sampling: Example

a,b,c



Gogate and Dechter, UAI 2008, CP2008

### AND/OR Importance Sampling (General Idea)

• Decompose Expectation  $P(d, f) = \sum_{a,b,c} P(a)P(c \mid a)P(b \mid a)P(d \mid b)P(f \mid c)$ 

 $Q(A, B, C) = Q(A)Q(B \mid A)Q(C \mid A)$ 

**Pseudo-tree** 

С

$$P(d, f) = \sum_{a,b,c} \frac{P(a)P(c \mid a)P(b \mid a)P(d \mid b)P(f \mid c)}{Q(a)Q(b \mid a)Q(c \mid a)} Q(a)Q(b \mid a)Q(c \mid a)$$
  
=  $E_{Q} \left[ \frac{P(a)P(c \mid a)P(b \mid a)P(d \mid b)P(f \mid c)}{Q(a)Q(b \mid a)Q(c \mid a)} \right]$ 

# AND/OR Importance Sampling (General Idea)

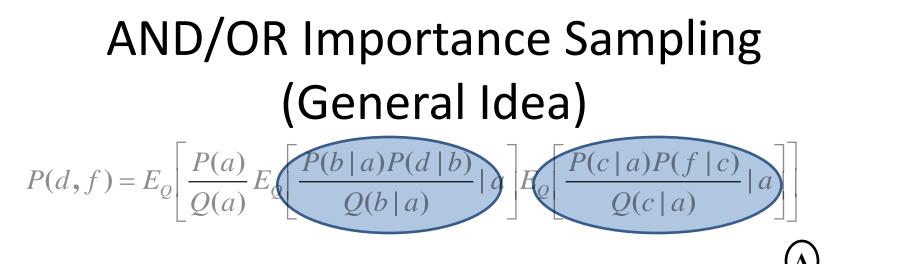
Decompose Expectation

**Pseudo-tree** 

 $P(d, f) = \sum_{a,b,c} \frac{P(a)P(c \mid a)P(b \mid a)P(d \mid b)P(f \mid c)}{Q(a)Q(b \mid a)Q(c \mid a)} Q(a)Q(b \mid a)Q(c \mid a)$ 

$$P(d,f) = \sum_{a} \frac{P(a)}{Q(a)} Q(a) \sum_{b} \frac{P(b \mid a) P(d \mid b)}{Q(b \mid a)} Q(b \mid a) \sum_{c} \frac{P(c \mid a) P(f \mid c)}{Q(c \mid a)} Q(c \mid a)$$

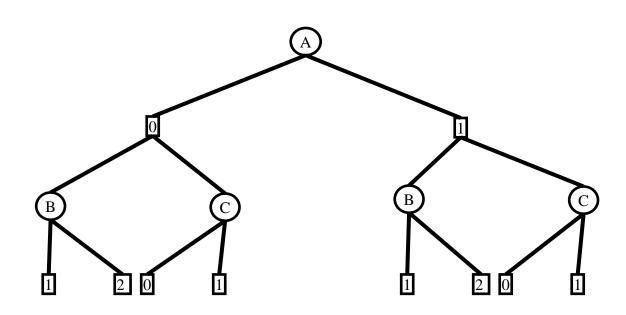
$$P(d,f) = E_{Q} \begin{bmatrix} \frac{P(a)}{Q(a)} E_{Q} \begin{bmatrix} P(b \mid a)P(d \mid b) \\ Q(b \mid a) \end{bmatrix} B_{Q} \begin{bmatrix} P(c \mid a)P(f \mid c) \\ Q(c \mid a) \end{bmatrix}$$

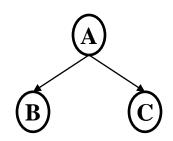


- Compute all expectations separately
- How?
  - Record all samples
  - For each sample that has A=a
    - Estimate the conditional expectations separately using the generated samples
    - Combine the results

Pseudo-tree

### **AND/OR Importance Sampling**





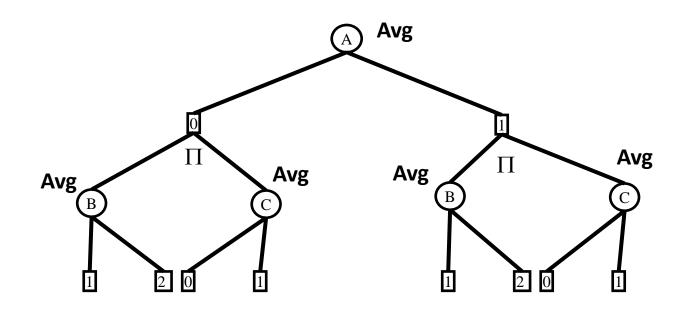
**Pseudo-tree** 

Sample #	A	B	С
1	0	1	0
2	0	2	1
3	1	1	1
4	1	2	0

Estimate of 
$$E\left[\frac{P(b \mid A=0)P(d \mid b)}{Q(b \mid A=0)} \mid A=0\right]$$

= Average Weight of samples of B having A = 0 =  $\frac{w(B=1, A=0) + w(B=2, A=0)}{2}$ 

### AND/OR Importance Sampling



Sample #	Z	X	Y
1	0	1	0
2	0	2	1
3	1	1	1
4	1	2	0

All AND nodes: Separate Components. Take Product

#### **Operator: Product**

All OR nodes: Conditional Expectations given the assignment above it

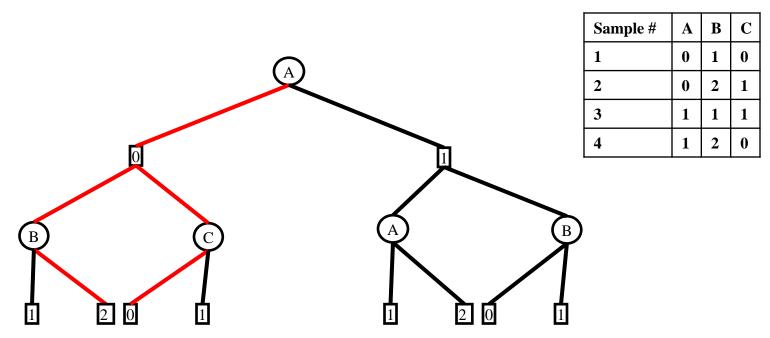
#### **Operator: Weighted Average**

### Algorithm AND/OR Importance Sampling

- 1. Construct a pseudo-tree.
- 2. Construct a proposal distribution along the pseudo-tree
- 3. Generate samples  $\mathbf{x}_1, \ldots, \mathbf{x}_N$  from Q along O.
- 4. Build a AND/OR sample tree for the samples  $x_1, \ldots, x_N$  along the ordering O.
- 5. FOR all leaf nodes *i of AND-OR tree do* 
  - 1. IF AND-node v(i)= 1 ELSE v(i)=0
- 6. FOR every node *n* from leaves to the root do
  - 1. IF AND-node v(n)=product of children
  - 2. IF OR-node v(n) = Average of children
- 7. Return v(root-node)

Gogate and Dechter, UAI 2008, CP2008

### # samples in AND/OR vs Conventional



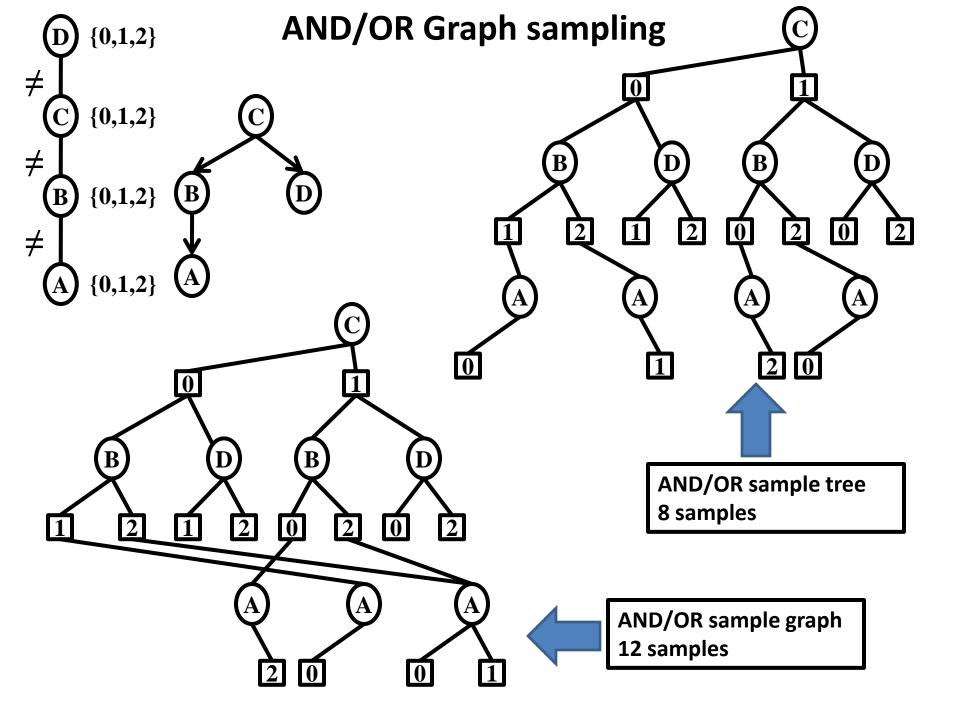
- 8 Samples in AND/OR space versus 4 samples in importance sampling
- Example: A=0, B=2, C=0 is not generated but still considered in the AND/OR space

# Why AND/OR Importance Sampling

- AND/OR estimates have smaller variance.
- Variance Reduction
  - Easy to Prove for case of complete independence (Goodman, 1960)

$$V[\overline{xy}] = \frac{V[x]E[y]^{2}}{N} + \frac{V[y]E[x]^{2}}{N} + \frac{V[x]V[y]}{N}, \text{ not independent}$$
Note the squared  
$$V[\overline{xy}] = \frac{V[x]E[y]^{2}}{N} + \frac{V[y]E[x]^{2}}{N} + \frac{V[x]V[y]}{N^{2}}, \text{ independent}$$
term.

 Complicated to prove for general conditional independence case (See Vibhav Gogate's thesis)!



Combining AND/OR sampling and w-cutset sampling  $Var_{Q}[\hat{P}(e)] = Var_{Q}\left[\frac{1}{N}\sum_{i=1}^{N}w(z^{i})\right] = \frac{Var_{Q}[w(z)]}{N}$ 

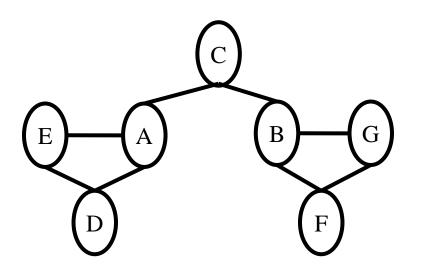
- Reduce the variance of weights
  - Rao-Blackwellised w-cutset sampling (Bidyuk and Dechter, 2007)
- Increase the number of samples; kind of
  - AND/OR Tree and Graph sampling (Gogate and Dechter, 2008)
- Combine the two

### Algorithm AND/OR w-cutset sampling

Given an integer constant w

- 1. Partition the set of variables into K and R, such that the treewidth of R is bounded by w.
- 2. AND/OR sampling on K
  - 1. Construct a pseudo-tree of K and compute Q(K) consistent with K
  - 2. Generate samples from Q(K) and store them on an AND/OR tree
- 3. Rao-Blackwellisation (Exact inference) at each leaf
  - 1. For each leaf node of the tree compute Z(R|g) where g is the assignment from the leaf to the root.
- 4. Value computation: Recursively from the leaves to the root
  - 1. At each AND node compute product of values at children
  - 2. At each OR node compute a weighted average over the values at children
- 5. Return the value of the root node

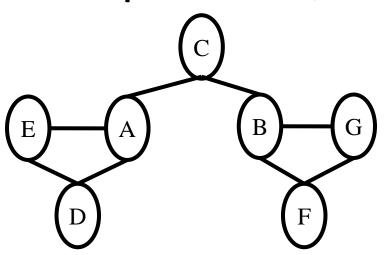
### AND/OR w-cutset sampling: Step 1: Partition the set of variables

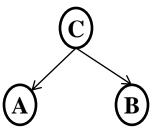


Practical constraint: Can only perform exact inference if the treewidth is bounded by 1.

**Graphical model** 

### AND/OR w-cutset sampling: Step 2: AND/OR sampling over {A,B,C}

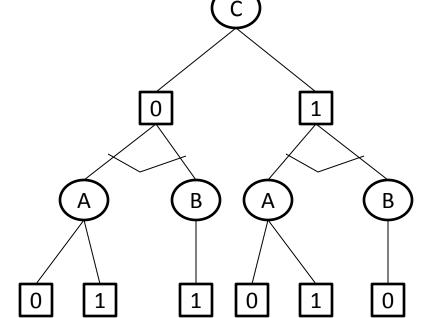


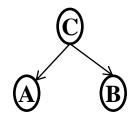


**Pseudo-tree** 

Graphical model

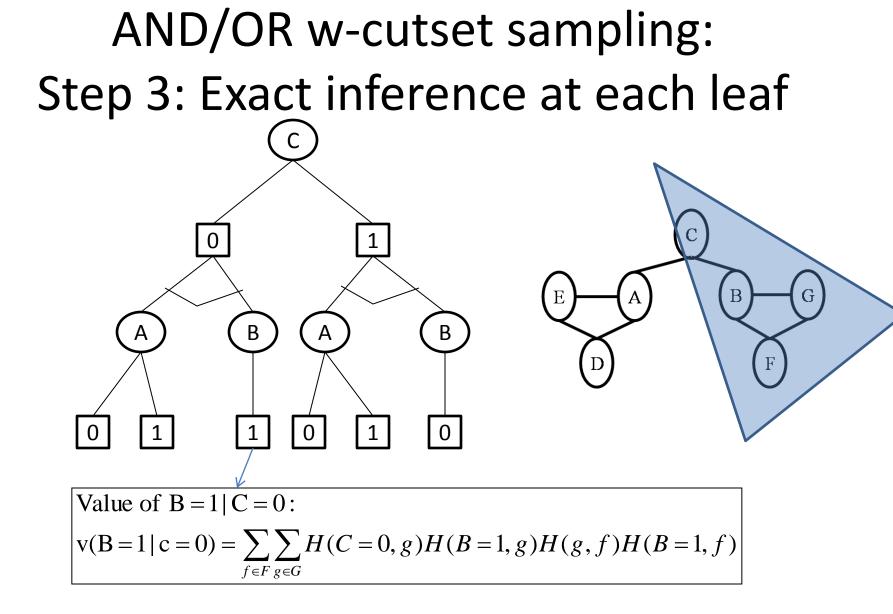
# AND/OR w-cutset sampling: Step 2: AND/OR sampling over {A,B,C}



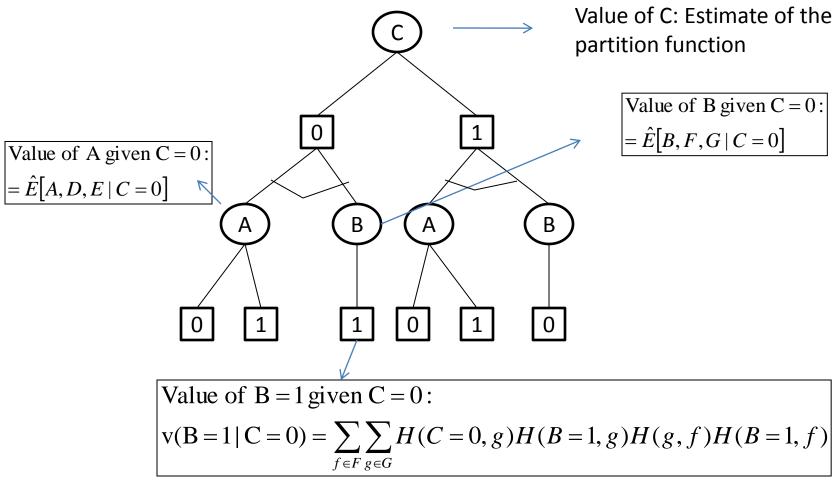


**Pseudo-tree** 

Samples: (C=0,A=0,B=1), (C=0,A=1,B=1), (C=1,A=0,B=0), (C=1,A=1,B=0)



### AND/OR w-cutset sampling: Step 4: Value computation



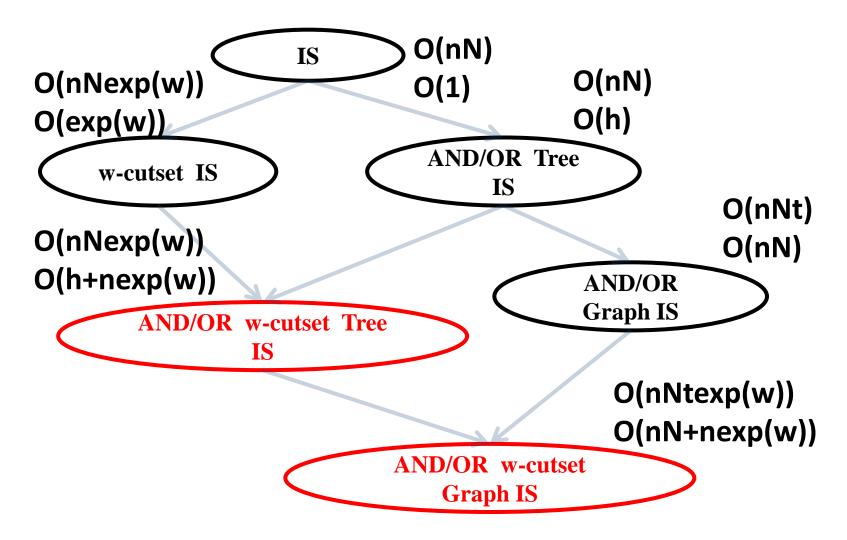
### **Properties and Improvements**

- Basic underlying scheme for sampling remains the same
  - The only thing that changes is what you estimate from the samples
  - Can be combined with any state-of-the-art importance sampling technique
- Graph vs Tree sampling
  - Take full advantage of the conditional independence properties uncovered from the primal graph

### AND/OR w-cutset sampling Advantages and Disadvantages

- Advantages
  - Variance Reduction
  - Relatively fewer calls to the Rao-Blackwellisation step due to efficient caching (Lazy Rao-Blackwellisation)
  - Dynamic Rao-Blackwellisation when context-specific or logical dependencies are present
    - Particularly suitable for Markov logic networks (Richardson and Domingos, 2006).
- Disadvantages
  - Increases time and space complexity and therefore fewer samples may be generated.

### Take away Figure: Variance Hierarchy and Complexity



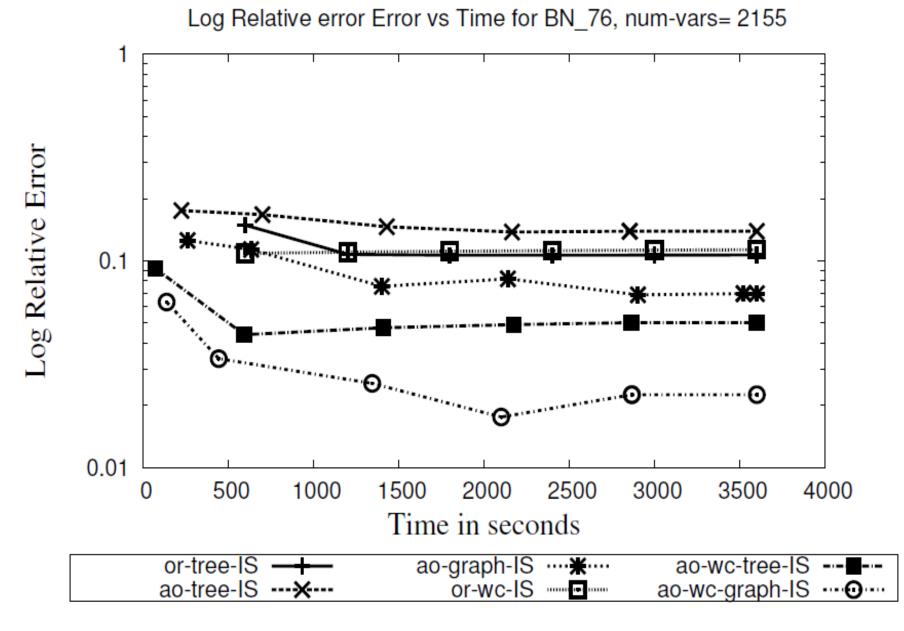
### Experiments

- Benchmarks
  - Linkage analysis
  - Graph coloring
- Algorithms
  - OR tree sampling
  - AND/OR tree sampling
  - AND/OR graph sampling
  - w-cutset versions of the three schemes above

### Results: Probability of Evidence Linkage instances (UAI 2006 evaluation)

Problem	$\langle n, k, E, t^*, c \rangle$	Exact	or-	<b>ao-</b>	<b>ao-</b>	or-wc-	ao-wc-	ao-wc-
			tree-IS	tree-IS	graph-IS	tree-IS	tree-IS	graph-IS
			$\Delta$	$\Delta$	$\Delta$	$\Delta$	$\Delta$	Δ
BN_69.uai	$\langle 777, 7, 78, 47, 59 \rangle$	5.28E-54	2.26E-02	2.46E-02	2.43E-02	2.42E-02	2.34E-02	4.22E-03
BN_70.uai	$\langle 2315, 5, 159, 87, 98 \rangle$	2.00E-71	6.32E-02	7.25E-02	5.12E-02	8.18E-02	5.36E-02	2.62E-02
BN_71.uai	$\langle 1740, 6, 202, 70, 139 \rangle$	5.12E-111	6.74E-02	5.51E-02	2.35E-02	8.58E-02	9.46E-03	1.21E-02
BN_72.uai	$\langle 2155, 6, 252, 86, 88 \rangle$	4.21E-150	3.19E-02	4.61E-02	2.46E-03	6.12E-02	1.41E-03	2.63E-03
BN_73.uai	$\langle 2140, 5, 216, 101, 149 \rangle$	2.26E-113	1.18E-01	1.12E-01	4.55E-02	1.58E-01	3.54E-02	3.95E-02
BN_74.uai	$\langle 749, 6, 66, 45, 72 \rangle$	3.75E-45	5.34E-02	4.31E-02	2.87E-02	8.08E-02	2.83E-02	2.76E-02
BN_75.uai	$\langle 1820, 5, 155, 92, 131 \rangle$	5.88E-91	4.47E-02	8.15E-02	4.73E-02	7.28E-02	4.20E-02	7.60E-03
BN_76.uai	$\langle 2155, 7, 169, 64, 239 \rangle$	4.93E-110	1.07E-01	1.39E-01	6.95E-02	1.13E-01	5.03E-02	2.26E-02
BN_77.uai	$\langle 1020,9,135,22,97\rangle$	6.88E-79	1.06E-01	9.38E-02	8.26E-02	1.24E-01	6.75E-02	3.27E-02

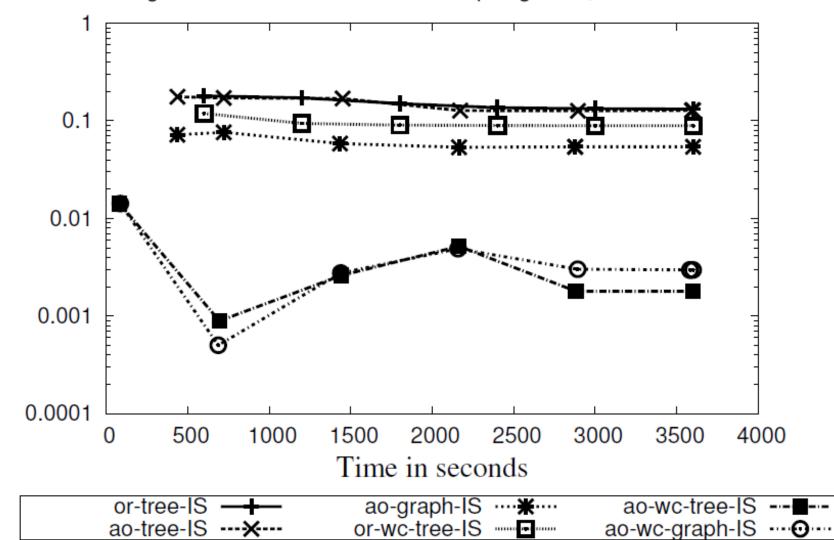
Time Bound: 1hr



### Results: Probability of Evidence Linkage instances (UAI 2008 evaluation)

Problem	$\langle n, k, E, t^*, w \rangle$	Exact	or-	ao-	ao-	or-wc-	ao-wc-	ao-wc-
			tree-IS	tree-IS	graph-IS	tree-IS	tree-IS	graph-IS
			$\Delta$	$\Delta$	$\Delta$	$\Delta$	$\Delta$	$\Delta$
pedigree18.uai	$\langle 1184, 1, 0, 26, 72 \rangle$	4.19E-79	3.17E-02	3.44E-02	3.20E-03	4.30E-02	3.49E-04	3.02E-04
pedigree19.uai	$\langle 793, 2, 0, 23, 102 \rangle$	1.59E-60	1.32E-01	1.28E-01	5.41E-02	8.92E-02	1.79E-03	2.97E-03
pedigree1.uai	(334, 2, 0, 20, 27)	7.81E-15	2.18E-03	1.90E-03	1.73E-04	3.15E-05	7.61E-06	1.13E-05
pedigree20.uai	$\langle 437, 2, 0, 25, 33 \rangle$	2.34E-30	1.52E-01	1.56E-01	2.12E-03	6.93E-02	9.17E-04	1.18E-03
pedigree23.uai	$\langle 402, 1, 0, 26, 29 \rangle$	2.00E-40	2.62E-02	2.74E-02	2.90E-02	2.82E-02	2.88E-02	2.88E-02
pedigree37.uai	(1032, 1, 0, 25, 36)	2.63E-117	2.46E-02	3.50E-03	3.24E-03	1.45E-02	3.00E-03	3.01E-03
pedigree38.uai	(724, 1, 0, 18, 45)	5.64E-55	4.08E-02	1.40E-02	1.25E-02	1.69E-02	8.91E-03	8.79E-03
pedigree39.uai	(1272, 1, 0, 29, 42)	6.32E-103	8.67E-02	5.11E-02	1.72E-03	1.89E-02	2.31E-04	2.13E-04
pedigree42.uai	$\langle 448, 2, 0, 23, 50 \rangle$	1.73E-31	4.29E-03	1.94E-03	5.06E-04	1.11E-03	3.53E-05	3.17E-05
pedigree31.uai	$\langle 1183, 2, 0, 45, 118 \rangle$		1.09E-01	1.31E-01	4.15E-02	8.34E-02	0.00E+00	2.93E-04
pedigree34.uai	$\langle 1160, 1, 0, 59, 104 \rangle$		2.12E-01	1.47E-01	8.37E-02	8.09E-02	4.83E-04	0.00E+00
pedigree13.uai	$\langle 1077, 1, 0, 51, 98 \rangle$		3.93E-01	3.93E-01	5.66E-02	9.11E-02	1.51E-04	0.00E+00
pedigree41.uai	(1062, 2, 0, 52, 95)		1.12E-01	5.06E-02	8.23E-04	5.04E-02	0.00E+00	3.15E-04
pedigree44.uai	(811, 1, 0, 29, 64)		3.16E-02	3.08E-02	2.27E-03	1.90E-02	0.00E+00	4.63E-06
pedigree51.uai	$\langle 1152, 1, 0, 51, 106 \rangle$		9.22E-02	6.39E-02	2.26E-02	4.31E-02	9.35E-05	0.00E+00
pedigree7.uai	$\langle 1068, 1, 0, 56, 90 \rangle$		7.86E-02	9.98E-02	2.31E-02	4.61E-02	4.38E-04	0.00E+00
pedigree9.uai	$\langle 1118, 2, 0, 41, 80 \rangle$		3.29E-02	3.19E-02	0.00E+00	8.25E-02	9.74E-03	1.01E-02

Time Bound: 1hr



#### Log Relative error Error vs Time for pedigree19, num-vars= 793

### Results: Solution counting Graph coloring instance

Problem	$\langle n,k,E,t^*,c\rangle$	Exact	or-	ao-	ao-	or-wc-	ao-wc-	ao-wc-
			tree-IS	tree-IS	graph-IS	tree-IS	tree-IS	graph-IS
			Δ	$\Delta$	$\Delta$	Δ	$\Delta$	$\Delta$
4-coloring1.uai	$\langle 400, 2, 0, 71, 309 \rangle$		3.82E-03	4.05E-03	4.51E-03	6.00E-03	2.35E-03	0.00E+00
4-coloring2.uai	$\langle 400, 2, 0, 95, 315 \rangle$		1.23E-02	9.54E-03	7.64E-03	3.38E-02	3.63E-02	0.00E+00
4-coloring3.uai	$\langle 800,2,0,144,617\rangle$		2.86E-03	4.58E-03	2.32E-03	2.41E-02	2.38E-02	0.00E+00
4-coloring4.uai	$\langle 800,2,0,191,620\rangle$		2.13E-02	5.06E-03	2.19E-02	1.79E-02	4.69E-03	0.00E+00
4-coloring5.uai	$\langle 1200, 2, 0, 304, 925 \rangle$		2.98E-02	2.81E-02	5.85E-02	5.70E-02	3.89E-02	0.00E+00
4-coloring6.uai	$\langle 1200, 2, 0, 338, 929 \rangle$		3.43E-02	2.72E-02	2.63E-03	3.17E-03	2.09E-03	0.00E+00

**Time Bound: 1hr** 

# Summary: AND/OR Importance sampling

- AND/OR sampling: A general scheme to exploit conditional independence in sampling
- Theoretical guarantees: lower sampling error than conventional sampling
- Variance reduction orthogonal to Rao-Blackwellised sampling.
- Better empirical performance than conventional sampling.